

CSCI 301 Lecture 26: LL(1) Parsing

You have a grammar G .

You have removed left recursion.

You have removed common prefixes.

Is it LL(1)? *Maybe.*

Running Example: Reverse Polish Notation
(postfix)

Infix: $1 + (4 - 2) * 3$

Polish (Racket) Prefix: $(+ 1 (* (- 4 2) 3))$

Reverse Polish

Postfix: $1 4 2 - 3 * +$

RPN: Look ma, no parentheses (or PEMDAS)! *100*

Evaluation with a stack:

$1 \ 4 \ 2 \ - \ 3 \ * \ +$

2	3		
4	2	6	
1	1	1	7
5	5	5	5

Ex. P+A

RPN Grammar

- no left recursion 👍
- no common prefixes 👍

Example:

Parse 12_3_+

$S \rightarrow DNT$

1 NT

1 DNT

12 NT

12 T

12_ S

12_ DNT

12_3 NT

12_3 T

12_3_ S

12_3_ PS

12_3_+ S

12_3_+

RPN expression

number

digit

operator

$S \rightarrow _S \mid PS \mid DNT \mid \epsilon$

$T \rightarrow _S \mid PS \mid \epsilon$

$N \rightarrow DN \mid \epsilon$

$D \rightarrow 0 \mid 1 \mid 2 \mid \dots \mid 9$

$P \rightarrow + \mid - \mid * \mid /$

Goal: Construct a Parse table:

		lookahead symb			
		$\downarrow$ D	P	-	\$ ← end of input
leftmost nonterminal	$S \rightarrow$	DN			
	$T \rightarrow$				
	$N \rightarrow$				
	D				
	P				

How? FIRST and FOLLOW

Let A be a variable in V .

$FIRST(A)$ is the set of terminals
any derivation from A can start
with.

$S \rightarrow$	$_S$	$ PS$	$ DNT$	$ \epsilon$
$T \rightarrow$	$_S$	$ PS$	$ \epsilon$	
$N \rightarrow$	DN	$ \epsilon$		
$D \rightarrow$	0	$ 1$	$ 2$	$ \dots 9$
$P \rightarrow$	$+$	$ -$	$ *$	$ /$

$$FIRST(A) = \{x \in \Sigma : A \Rightarrow^* x\alpha \text{ for any } \alpha \in (\Sigma \cup V)^*\}$$

Calculating FIRST:

1. If A is a terminal, $FIRST(A) = \{A\}$
2. If A is a nonterminal and $A \rightarrow Y_1 Y_2 \dots Y_k$,
add $FIRST(Y_1)$ to $FIRST(A)$ ↪ nullable? $Y_1 \Rightarrow^* \epsilon$

If $Y_1 \Rightarrow^* \epsilon$, we need to add $FIRST(Y_2)$ too

3. If $Y_1, \dots, Y_q$ are nullable, add $FIRST(Y_{q+1})$ to $FIRST(A)$

FOLLOW

What if $A \Rightarrow^* \epsilon$?

$FOLLOW(A)$: all terminals that can
come after A in a derived string

(show RPN parse table)

Recursive Descent Parsing

- Each Nonterminal is a function
- It applies rules by calling functions from RHS
- It returns the rest of the string still to be parsed

Input: 44_9_+

	<u>Parsed</u>	<u>Rest</u>
$\rightarrow S(44_9_+)$	ϵ	44_9_+
$\quad D(44_9_+)$	4	4_9_+
$\quad \rightarrow N(4_9_+)$	4	4_9_+
$\quad \quad D(4_9_+)$	44	_9_+
$\quad \quad N(_9_+)$		_9_+
$\quad \quad T(_9_+)$		