

CSCI 301 Lecture 26: LL(1) Parsing

You have a grammar G .

You have removed left recursion.

You have removed common prefixes.

Is it LL(1)? **Maybe.**

Running Example: Reverse Polish Notation

(postfix)

Infix: $1 + (4 - 2) * 3$

$\overbrace{1 \ 1} \ 4 \ 2 \ \overbrace{+} \ 3 \ * \ +$
 $1 \ \overbrace{(2 \ 3 \ *} \ +$
 $1 \ \overbrace{\begin{matrix} 6 \\ - \end{matrix}} \ +$

Polish
(prefix)

Prefix: $(+ \ 1 \ (* \ (- \ 4 \ 2) \ 3)) \ \overbrace{1 \ \begin{matrix} 6 \\ - \end{matrix}} \ +$

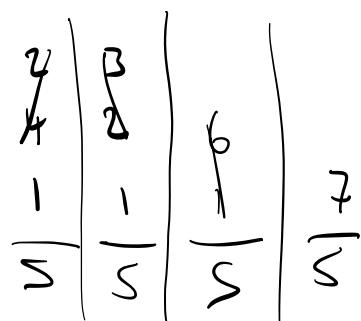
Reverse
Polish

Postfix: $1 \ 4 \ 2 \ - \ 3 \ * \ +$

RPN: Look ma, no parentheses (or PEMDAS)! **100**

Evaluation with a stack:

$1 \ 4 \ 2 \ - \ 3 \ * \ +$



Ex. Pta

RPN Grammar

- No left recursion
- No common prefixes

Example:

Parse $12 - 3 +$

$S \rightarrow DNT$

$| NT$
 $| DNT$
 $| 12 NT$
 $| 12 T$
 $| 12 - S$
 $| 12 - DNT$
 $| 12 - 3 NT$
 $| 12 - 3 T$
 $| 12 - 3 , S$
 $| 12 - 3 , PS$
 $| 12 - 3 - + S$
 $| 12 - 3 - +$

RPN expression

number
 digit
 operator

$S \rightarrow _S PS DNT \epsilon$
$T \rightarrow _S PS \epsilon$
$N \rightarrow DN \epsilon$
$D \rightarrow 0 1 2 \dots 9$
$P \rightarrow + - * /$

Goal: Construct a Parse table:

left most non-terminal	lookahead symbol				end of input
	D	P	-	\$	
$S \rightarrow$	DN				
$T \rightarrow$					
$N \rightarrow$					
D					
P					

How? FIRST and FOLLOW

Let A be a variable in V.

FIRST(A) is the set of terminals
any derivation from A can start
with.

$S \rightarrow _S PS DN \epsilon$
$T \rightarrow _S PS \epsilon$
$N \rightarrow DN \epsilon$
$D \rightarrow 0 1 2 \dots 9$
$P \rightarrow + - * /$

$$\text{FIRST}(A) = \{x \in \Sigma : A \Rightarrow^* x\alpha \text{ for any } \alpha \in (\Sigma \cup V)^*\}$$

Calculating FIRST:

1. If A is a terminal, $\text{FIRST}(A) = \{A\}$

2. If A is a nonterminal and $A \Rightarrow Y_1 Y_2 \dots Y_k$,
add $\text{FIRST}(Y_i)$ to $\text{FIRST}(A)$ $\text{nullable? } Y_i \stackrel{?}{\Rightarrow} \epsilon$

If $Y_1 \Rightarrow^* \epsilon$, we need to add $\text{FIRST}(Y_2)$ too

3. If $Y_1 \dots Y_q$ are nullable, add $\text{FIRST}(Y_{q+1})$ to $\text{FIRST}(A)$

FOLLOW

What if $A \Rightarrow^* \epsilon$?

$\text{FOLLOW}(A)$: all terminals that can
come after A in a derived string

(Show RPN parse table)

Recursive Descent Parsing

- Each Nonterminal is a function
- It applies rules by calling functions from RHS
- It returns the rest of the string still to be parsed

Input: 44_9_+

$\rightarrow S(44_9_+)$

$D(44_9_+)$

$\rightarrow N(\underline{4_9_+})$

$D(4_9_+)$

$N(-9_+)$

$T(-9_+)$

Parsed Rest

ϵ 44_9_+

4 4_9_+

4 4_9_+

44 $_9_+$

-9_+