

# CSCI 301 - Lecture 6: Logic 3

## Logical Equivalence:

Two statements are logically equivalent if their truth tables match exactly. We write this  $P \equiv Q$

Example:

$P \Leftrightarrow Q$  is true when  $P$  and  $Q$  are both true or both false

| P | Q | $P \Leftrightarrow Q$ | $(P \wedge Q) \vee (\neg P \wedge \neg Q)$ |
|---|---|-----------------------|--|
| T | T | T                     | T  |
| T | F | F                     | F  |
| F | T | F                     | F  |
| F | F | T                     | T  |

$(P \wedge Q) \vee (\neg P \wedge \neg Q)$

Example: does  $\wedge$  distribute over  $\vee$ ?

$$A \wedge (B \vee C) \stackrel{?}{\equiv} (A \wedge B) \vee (A \wedge C)$$

| A | B | C | $A \wedge (B \vee C)$ | $(A \wedge B) \vee (A \wedge C)$ |
|---|---|---|-----------------------|----------------------------------|
| T | T | T | T                     | T                                |
| T | T | F | T                     | T                                |
| T | F | T | T                     | T                                |
| T | F | F | F                     | F                                |
| F | T | T | F                     | F                                |
| F | T | F | F                     | F                                |
| F | F | T | F                     | F                                |
| F | F | F | F                     | F                                |

Do Ex. A

| P | Q | $P \Rightarrow Q$ | $(P \wedge \neg Q)$ | $\neg(P \wedge \neg Q)$ |
|---|---|-------------------|---------------------|-------------------------|
| T | T | T                 |                     | T                       |
| T | F | F                 |                     | F                       |
| F | T | T                 |                     | T                       |
| F | F | T                 |                     | T                       |

$\neg(P \wedge \neg Q)$

$$P \Rightarrow Q$$

# Negating Statements

$$P \rightarrow \neg P, \text{ done!}$$

But there are some tricks.

Ex. #45 De Morgan's Laws: Negating  $\wedge$  and  $\vee$

$$\neg(P \wedge Q) \equiv \neg P \vee \neg Q$$

$$\neg(P \vee Q) \equiv \neg P \wedge \neg Q$$

Ex. #6: Negating a conditional.  $\neg(P \Rightarrow Q) \equiv \underline{P \wedge \neg Q}$

$$\neg(P \Rightarrow Q) \neq (Q \Rightarrow P)$$

Quantifiers?

$$\neg(\forall x, P(x)) \equiv \underline{\exists x, \neg P(x)}$$

$$\neg(\exists x, P(x)) \equiv \underline{\forall x, \neg P(x)}$$

Do Ex. B

1.  $x$  is positive, but  $y$  is not positive

2. Every even integer greater than 2 is the sum of two primes.

$$\forall x \in \mathbb{Z}, \left[ \underbrace{E(x) \wedge x > 2}_P \Rightarrow \underbrace{[x = y + z \wedge y \text{ is prime} \wedge z \text{ is prime}]}_Q \right]$$
$$\exists x \in \mathbb{Z}, \left[ E(x) \wedge x > 2 \wedge \left( x \neq y + z \vee y \text{ is not prime} \vee z \text{ is not prime} \right) \right]$$

3. At least one of the integers  $a$  and  $b$  is odd,

4.  $2a$  is even iff  $a$  is an integer

5. There exists a real number  $y$  for which  $x < y$  for any real number  $x$ .

# Logical Inference

Knowing some things what else can you conclude?  
(infer)  
(deduce)

Duh:

$$\frac{P \wedge Q}{P}$$

$$\frac{P}{Q}$$
$$\frac{P \wedge Q}{P \wedge Q}$$

$$\frac{P}{P \vee Q}$$

Somewhat less duh:

Modus Ponens:

$$\frac{P \Rightarrow Q}{P}$$

Modus Tollens:

$$\frac{P \Rightarrow Q}{\neg Q}$$

Elimination:

$$\frac{P \vee Q}{\neg P}$$

Def An integer  $n$  is \_\_\_\_\_ if  $n = 2a$  for some integer  $a$ .

Def \_\_\_\_\_; if  $n = 2a + 1$  for some integer  $a$ .

Def Suppose  $a$  and  $b$  are integers. We say that  $a$  \_\_\_\_\_  $b$ ,  
written \_\_\_\_\_, if  $b = ac$  for some integer  $c$ . In this case,  
we also say  $a$  is a \_\_\_\_\_ of  $b$ , and  $b$  is a \_\_\_\_\_ of  $a$ .

Def Two integers have the \_\_\_\_\_ if they are both even  
or both odd. Otherwise, they have \_\_\_\_\_.