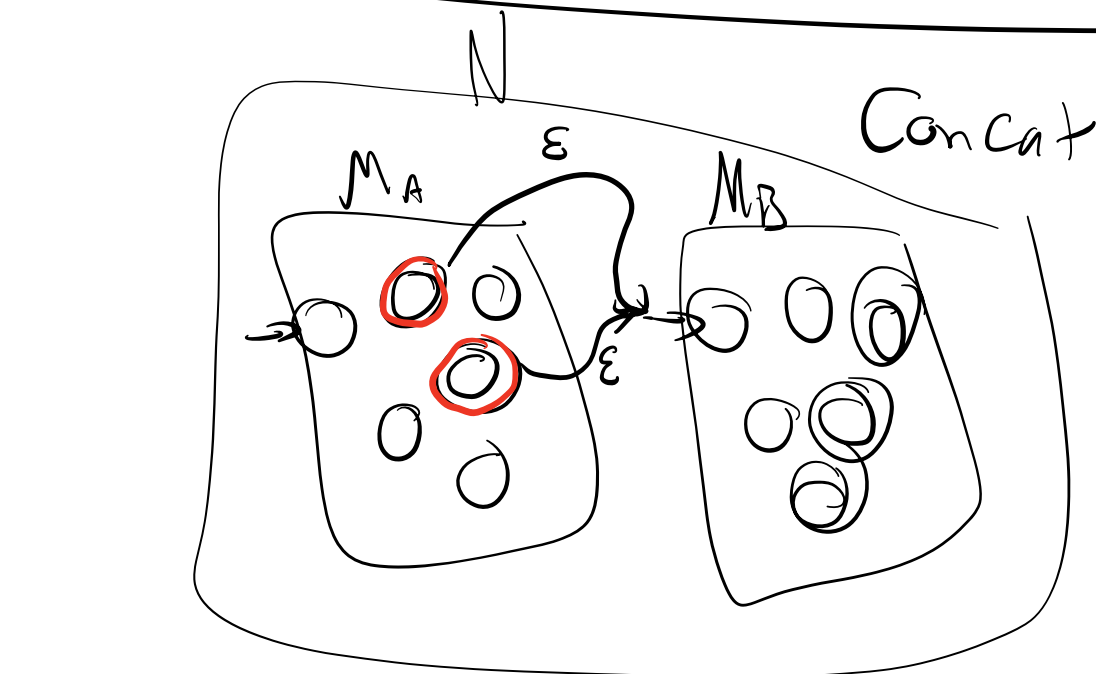
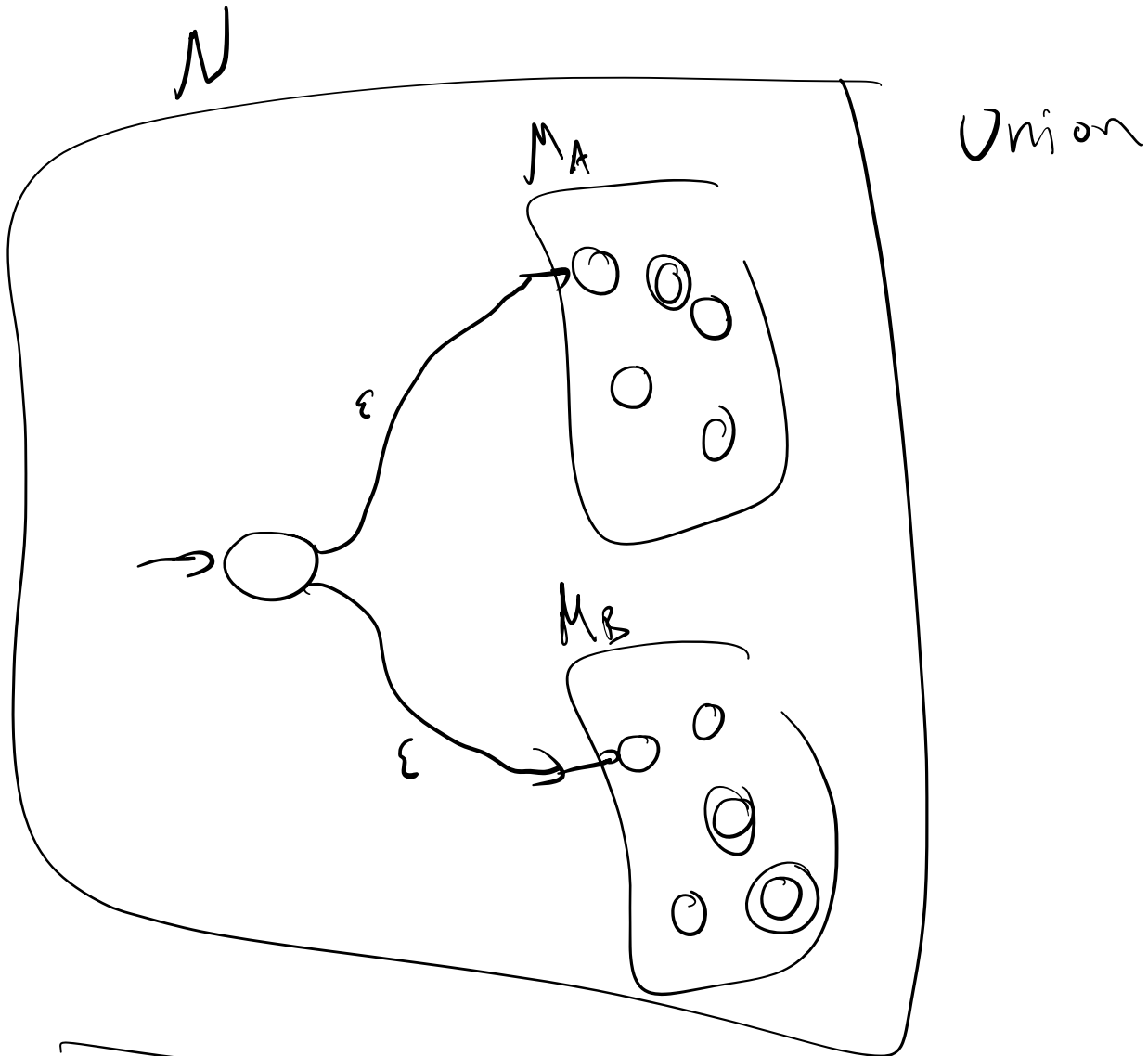
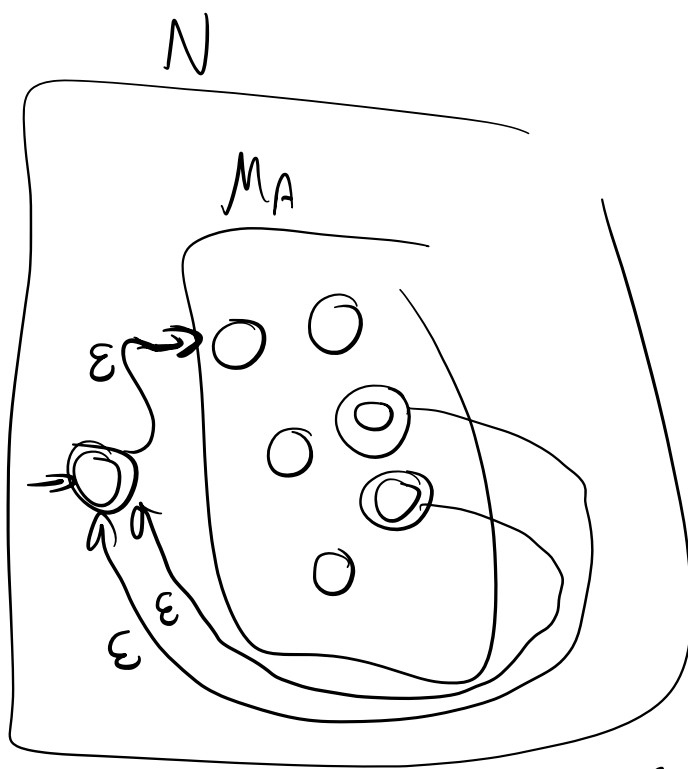


CSCI 301 - Lecture 26 - Regular Closure Proof

Regular Expressions

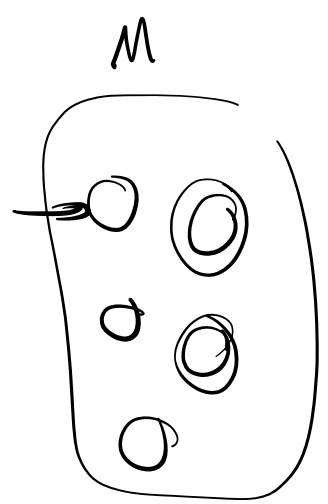




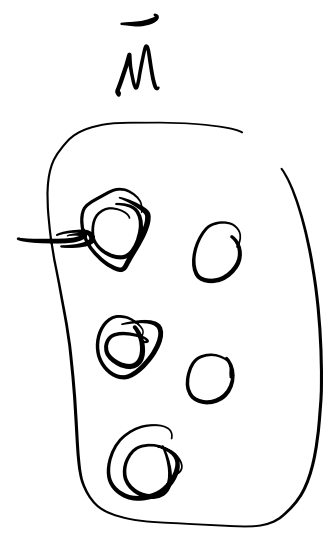
*

$$\rightarrow A = \{11, 00\}$$

$$\rightarrow A^* \subseteq \{\epsilon, 11, 00, 1111, 0000, 1100, 0011, \dots\}$$



\bar{A}



$$A \cap B = \overline{\bar{A} \cup \bar{B}}$$

Regular Expressions: A meta-language for regular languages.

Definition: A regular expression over alphabet Σ is defined inductively (recursively) as follows:

Regular Expression	Language described
• Σ is a R.E.	$\{\Sigma\}$
• \emptyset is a R.E.	\emptyset
• a is a R.E. for all $a \in \Sigma$	$\{a\}$
• If R_1, R_2 are R.E.s,	
- $R_1 \cup R_2$ is a R.E.	$R_1 \cup R_2$
- $R_1 R_2$ is a R.E.	$R_1 R_2$
• If R is a R.E., R^* is a R.E.	R^*

$$\Sigma = \{0, 1\}$$

$$(0 \cup 1) 0 1^*$$

$$\{00, 10, 001, 101, 0011, 1011, \dots\}$$

$$L = \{w : w \text{ has } 1011 \text{ as a substring}\}$$

$$\text{RE: } (0 \cup 1)^* 1011 (0 \cup 1)^*$$

R.E	L
1. ab	$\{ab\}$
2. $(a \cup b \cup \epsilon)$	$\{a, b, \epsilon\}$
4. $\underbrace{a(a \cup b)} \cup \underbrace{b(b \cup a)}$	$\{aa, ab, bb, ba\}$
6. $(aa)^*$	$\{\epsilon, aa, aaaa, aaaaaa, \dots\}$
8. $\underbrace{(a^* b^*)^*}_{\underline{a}ab\underline{a}}$	Σ^*

B.2 $\Sigma = \{a, b\}$
 $\{w : w \text{ has at least 3 a's}\}$

$(a \cup b)^* a (a \cup b)^* a (a \cup b)^* a (a \cup b)^*$

$\Sigma = \{a, b\} \{ |w| \text{ is odd} \}$

$\Sigma = \{0, 1\}$ { w : every odd position is 1 }