

CSCI 241

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Merge Sort: Runtime Analysis

Goals

Know how to derive the worst-case runtime of mergesort.

Mergesort: Runtime

A strategy for analyzing recursive methods:

1. Count work done in a call
excluding recursive calls.
2. Multiply by overall number of calls made

```
def fact(n):  
    if n <= 1:  
        return n  
    return n * fact(n-1)
```

1. $O(1)$ work per call
2. Called once per value in $1..n+1$ for a total of $O(n)$ work

Mergesort: Runtime

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2. Multiply by overall number of calls made

```
/** sort A[start..end] */
mergeSort(A, start, end):
    if (end-start < 2):
        return
    mid = (end+start)/2
    mergeSort(A, start, mid)
    mergeSort(A, mid, end)
    merge(A, start, mid, end)
```

Mergesort: Runtime

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```
/** sort A[start..end] */  
mergeSort(A, start, end):
```

```
    O(1)  if (end-start < 2):  
          return
```

```
    O(1)  mid = (end+start)/2
```

```
(excluded) mergeSort(A, start, mid)  
            mergeSort(A, mid, end)
```

```
    O(??) merge(A, start, mid, end)
```

Merge: Runtime

Let $n = \text{end} - \text{start}$

$O(\text{end} - \text{start})$

$O(n)$
 $O(1)$

$O(n)$

```
merge(A, start, mid, end):
```

```
  B = deep copy of A
```

```
  initialize i, j, and k
```

```
  while neither half is empty
```

```
    [ copy the smaller  
      "front" element into A
```

```
    [ copy any remaining  
      left half elements
```

```
    [ copy any remaining  
      right half elements
```

Mergesort: Runtime

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```
/** sort A[start..end] */
```

```
mergeSort(A, start, end):
```

```
    O(1)  if (end-start < 2):  
          return
```

```
    O(1)  mid = (end+start)/2
```

```
(excluded) mergeSort(A, start, mid)
```

```
mergeSort(A, mid, end)
```

```
O(end - start) merge(A, start, mid, end)
```

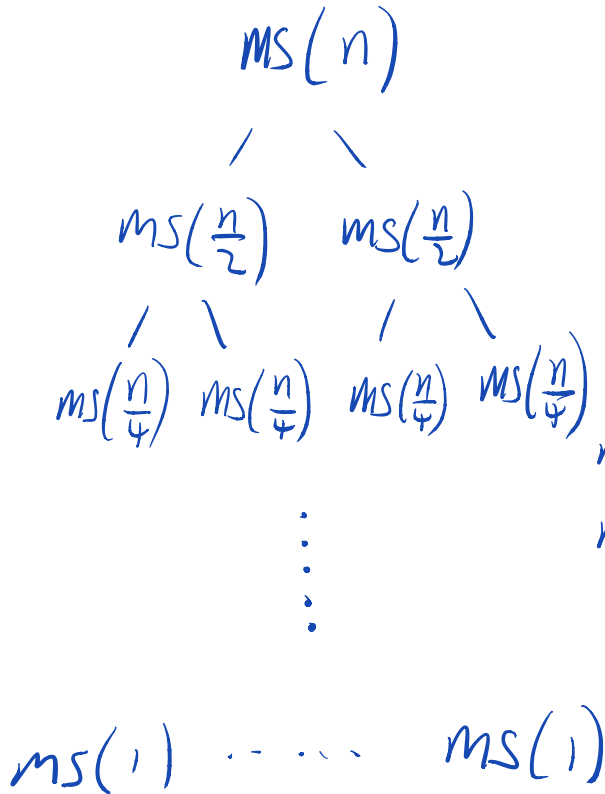
Mergesort: Runtime

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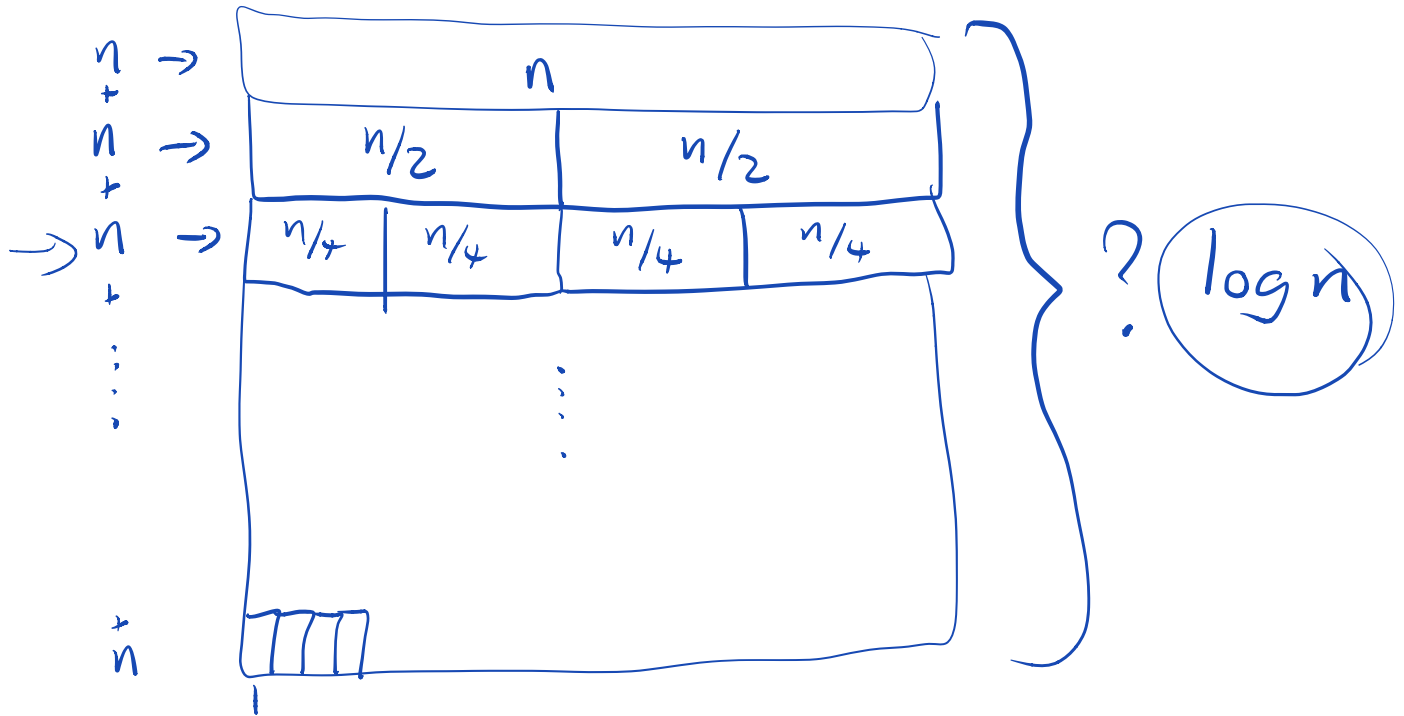
Problem: sometimes work depends on n , which varies from call to call.

Mergesort: Runtime



```
/** sort A[start..end] */
mergeSort(A, start, end):
  if (end-start < 2):
    return
  mid = (end+start)/2
  mergeSort(A, start, mid)
  mergeSort(A, mid, end)
  merge(A, start, mid, end)
```

$O(1)$ → mergeSort(A, start, mid)
 $n/2$ → mergeSort(A, mid, end)
 $O(n)$ merge(A, start, mid, end)



Overall work: $O(n \cdot \log n)$