

# CSCI 241

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Dijkstra's Algorithm:  
Detailed Pseudocode  
The Path Itself

# Goals

Know how to augment Dijkstra's algorithm to keep backpointers in order to reconstruct the sequence of nodes in a shortest path.

# Dijkstra's Shortest Paths: Pseudocode

```
S = { }; F = {v}; v.d = 0;
while (F ≠ { }) {
    f = node in F with min d value;
    Remove f from F, add it to S;
    for each neighbor w of f {
        if (w not in S or F) {
            w.d = f.d + weight(f, w);
            add w to F;
        } else if (f.d+weight(f,w) < w.d) {
            w.d = f.d+weight(f,w);
        }
    }
}
```

# Dijkstra's Shortest Paths: Pseudocode

Initialize Settled to empty  
Initialize Frontier to the start node

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            w.d = f.d+weight(f,w);
        }
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}
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Initialize Settled to empty
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While the frontier isn't empty:
    move node f with smallest d
    from F to S
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      w.d = f.d + weight(f, w);
    }
  }
}
```

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Initialize Settled to empty
Initialize Frontier to the start node

While the frontier isn't empty:
  move node f with smallest d
  from F to S

For each neighbor w of f:
  if we've never seen w before:
    set its path length
    add it to frontier
  else if path to w via f is shorter:
    update w's shortest path length
```

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# What if we want to know the shortest path?

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  for each neighbor w of f {
    if (w not in S or F) {
      w.d = f.d + weight(f, w);
      add w to F;
    } else if (f.d + weight(f, w) < w.d) {
      w.d = f.d + weight(f, w);
    }
  }
}
```

At termination: for each reachable node  $n$ ,  $n.d$  stores the **length** of the shortest path from  $v$  to  $n$ .

We didn't keep track of **how** to get from  $v$  to  $n$ !



# What if we want to know the shortest path?

```
S = { }; F = {v}; v.d = 0; v.bp = null;
while (F ≠ {}) {
  f = node in F with min d value;
  Remove f from F, add it to S;
  for each neighbor w of f {
    if (w not in S or F) {
      w.d = f.d + weight(f, w);
      w.bp = f;
      add w to F;
    } else if (f.d + weight(f, w) < w.d) {
      w.d = f.d + weight(f, w);
      w.bp = f
    }
  }
}
```

Each node could store the full path, but that would be expensive to keep updated.

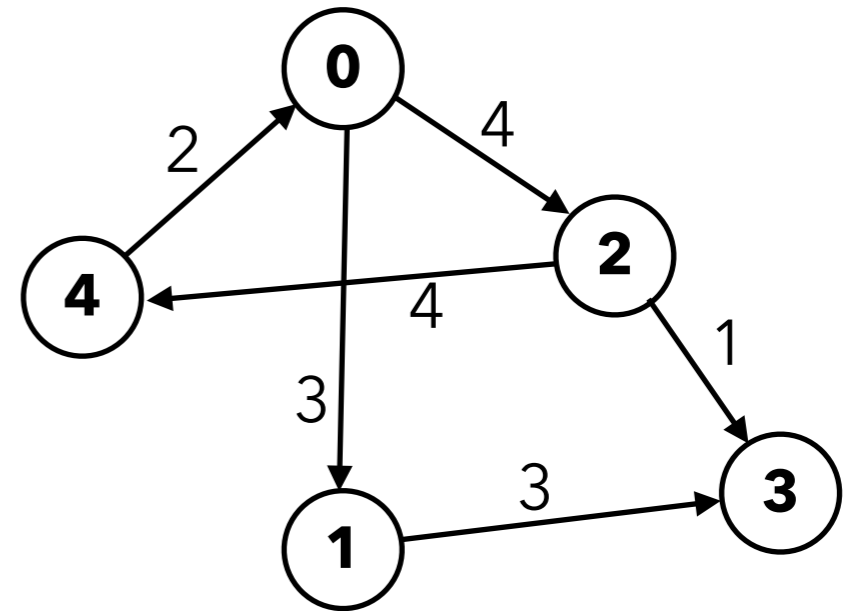
Strategy: maintain a **backpointer** at each node pointing to the previous node in the shortest path.

# What if we want to know the shortest path? Example

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            w.bp = f
        }
    }
}

```



shortest-paths ( 4 )

Node	d	bp
0	2	4
1	5	0
2	6	0
3	7	2
4	0	null

# What if we want to know the shortest path? Example

$S = \{ \}; F = \{v\}; v.d = 0; v.bp = \text{null};$

**while** ( $F \neq \{ \}$ ) {

$f = \text{node in } F \text{ with min } d \text{ value};$

  Remove  $f$  from  $F$ , add it to  $S$ ;

**for** each neighbor  $w$  of  $f$  {

**if** ( $w$  not in  $S$  or  $F$ ) {

$w.d = f.d + \text{weight}(f, w);$

$w.bp = f;$

      add  $w$  to  $F$ ;

    } **else if** ( $f.d + \text{weight}(f, w) < w.d$ ) {

$w.d = f.d + \text{weight}(f, w);$

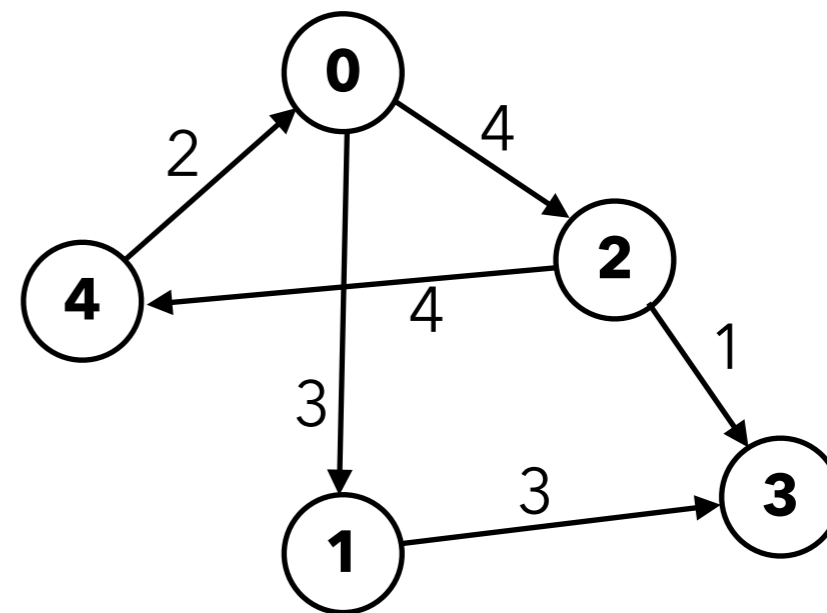
$w.bp = f$

  }

    Shortest path to Node 3:

}

}



shortest-paths ( 4 )

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2	6	0
3	7	2
4	0	null

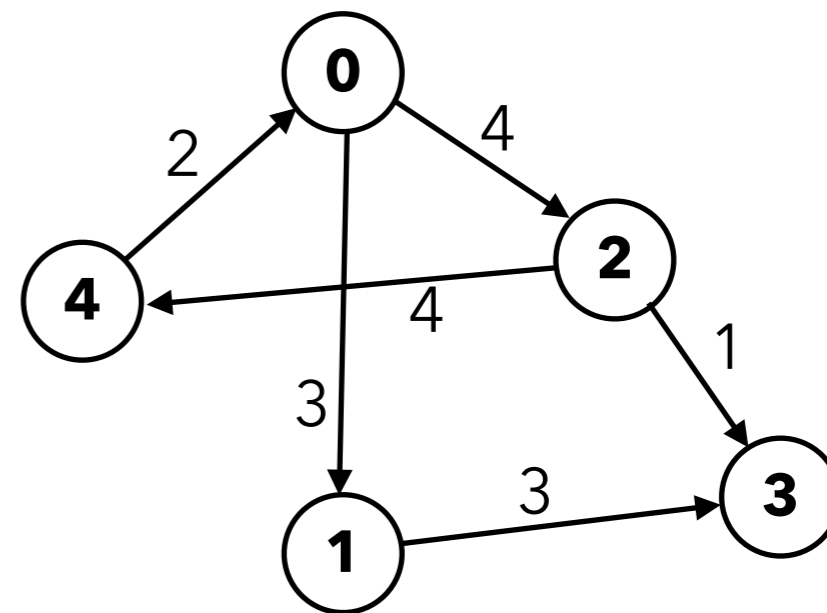
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Shortest path to Node 3: **3**



shortest-paths ( 4 )

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3	7	2
4	0	null

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Remove  $f$  from  $F$ , add it to  $S$ ;

**for** each neighbor  $w$  of  $f$  {

**if** ( $w$  not in  $S$  or  $F$ ) {

$w.d = f.d + \text{weight}(f, w);$

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add  $w$  to  $F$ ;

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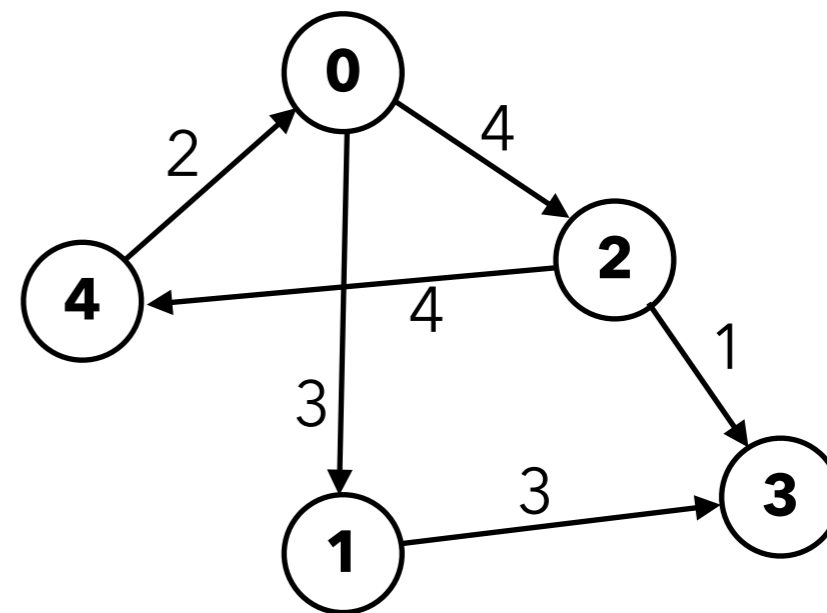
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}

Shortest path to Node 3:

**2 3**



shortest-paths ( 4 )

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0	2	4
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2	6	0
3	7	2
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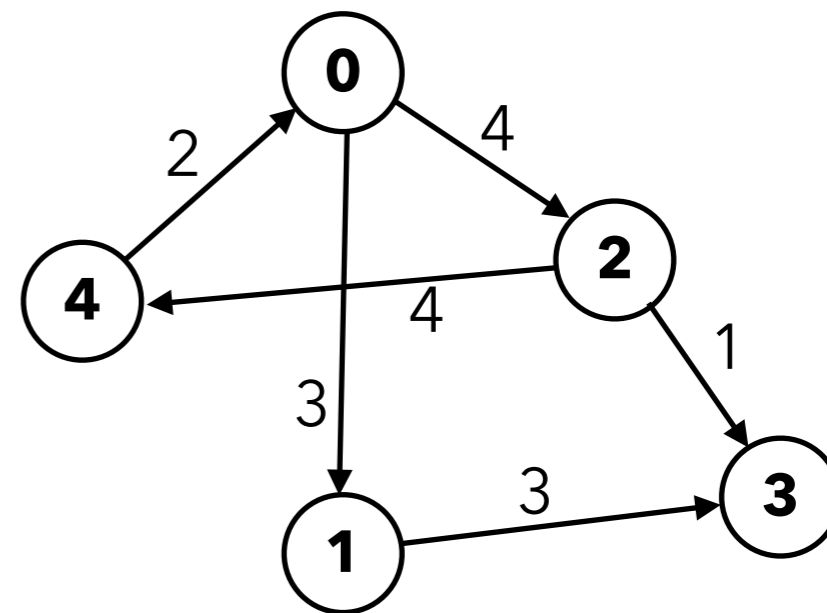
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Shortest path to Node 3:

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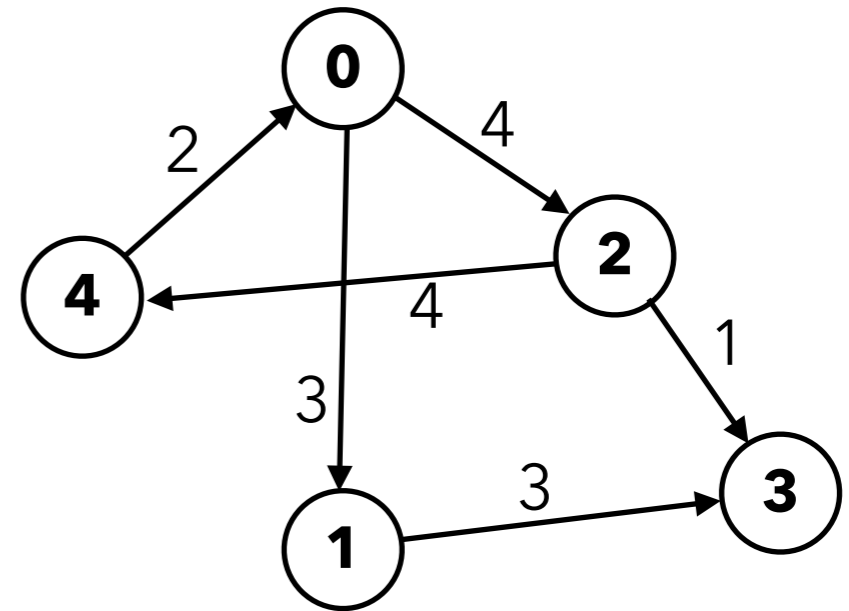
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Shortest path to Node 3:

**4 0 2 3**



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