

CSCI 241

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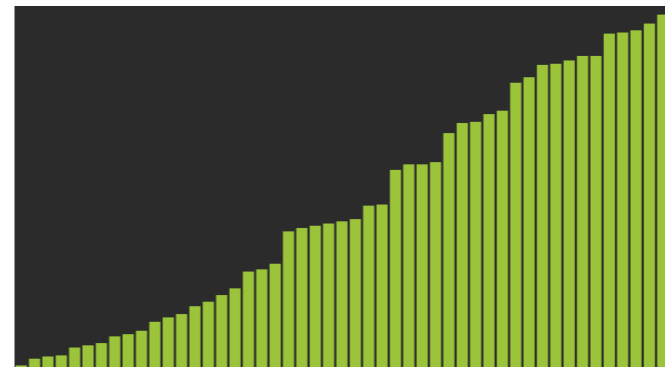
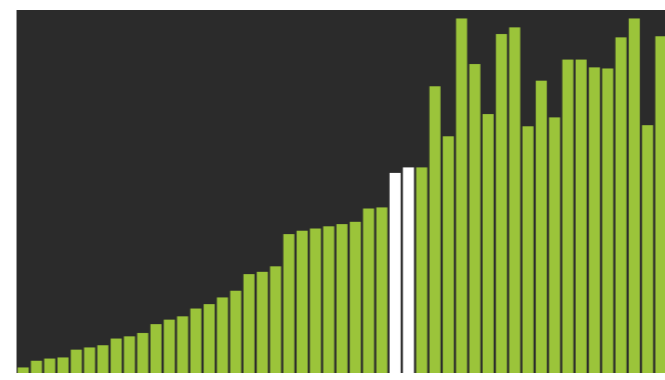
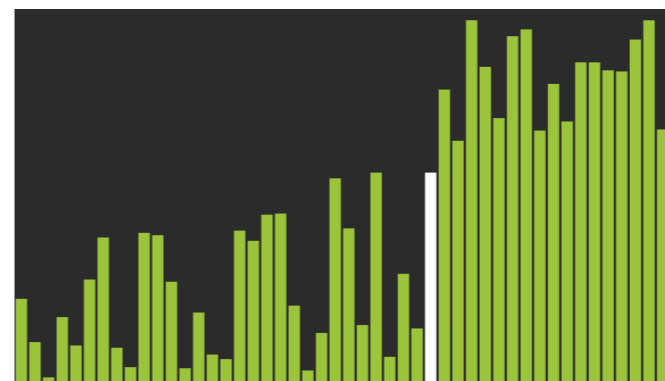
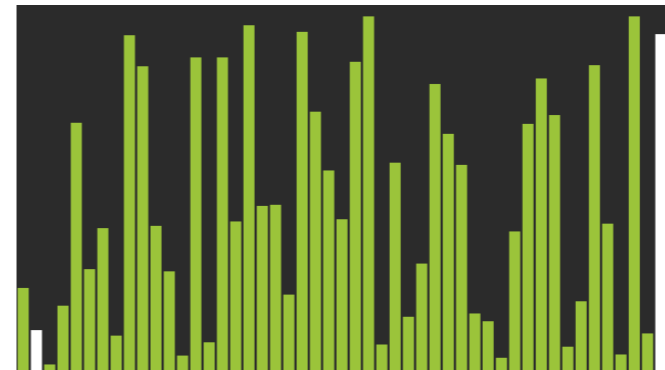
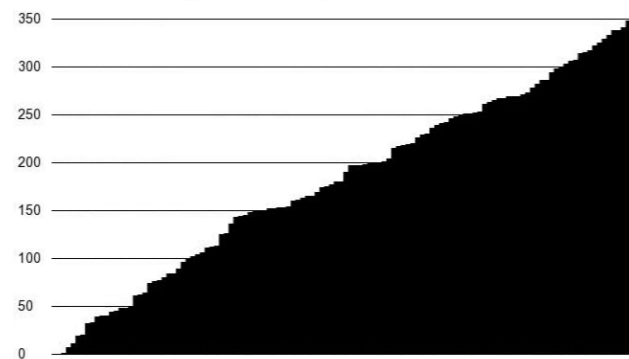
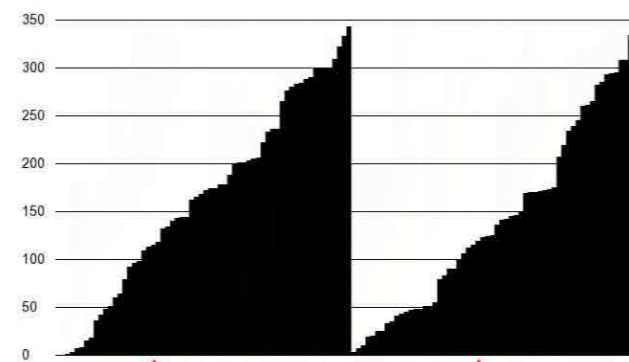
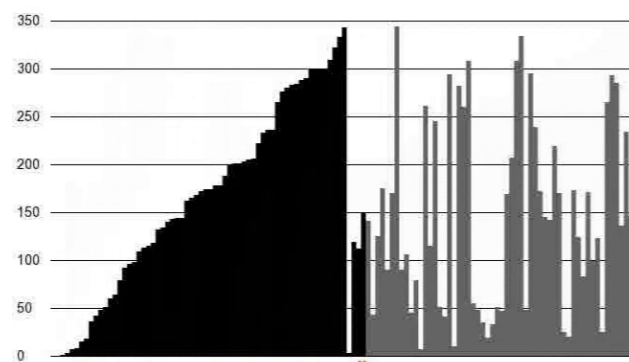
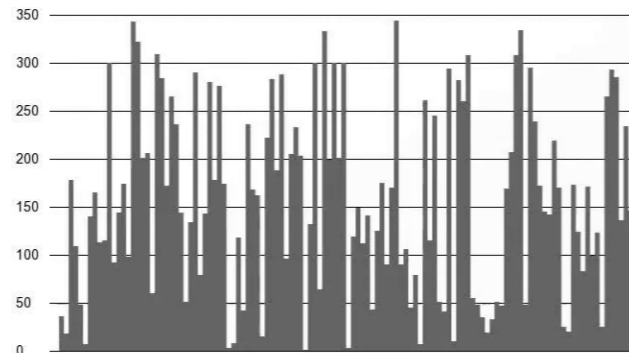
Quick Sort: Runtime

Goals

Understand the best-case and worst-case runtime analysis of quicksort.

Know the average-case runtime of quicksort.

Merge vs Quick



"real work"
done here

"real work"
done here

Quicksort: Runtime

```
/** quicksort A[st..end]*/  
quickSort(A, st, end):  
    if (small):  
        return
```

$O(1)$

$O(??)$

```
mid = partition(A, st, end)
```

(excluded)

```
quickSort(A, st, mid)  
quickSort(A, mid+1, end)
```

```
# (nothing to do!)
```

Partition: Runtime

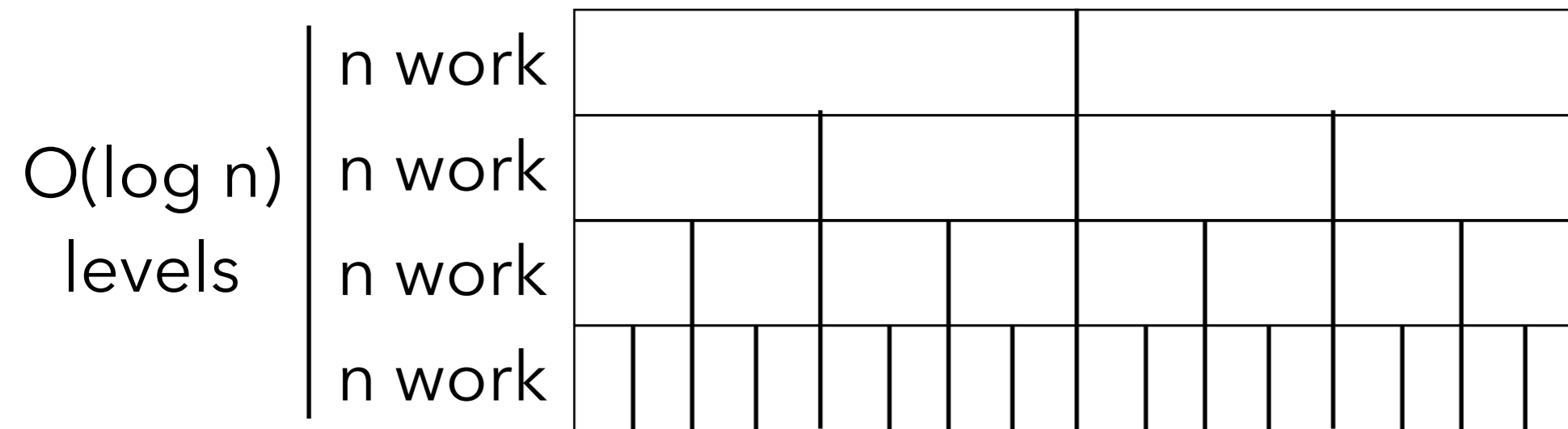
```
partition(A, start, end)
  initialize i, j
O(1)  choose pivot
      swap pivot to A[0]
      while [?] section != []
        # process A[i]:
        if <= p:
n * O(1)  move to <= p section
        else:
          move to > p section
```

Total: $O(n)$, where $n = \text{end} - \text{start}$.

Runtime: Best case

Best case:

- pivot is the median of the array
- partition splits the array exactly in half
- same analysis as merge sort

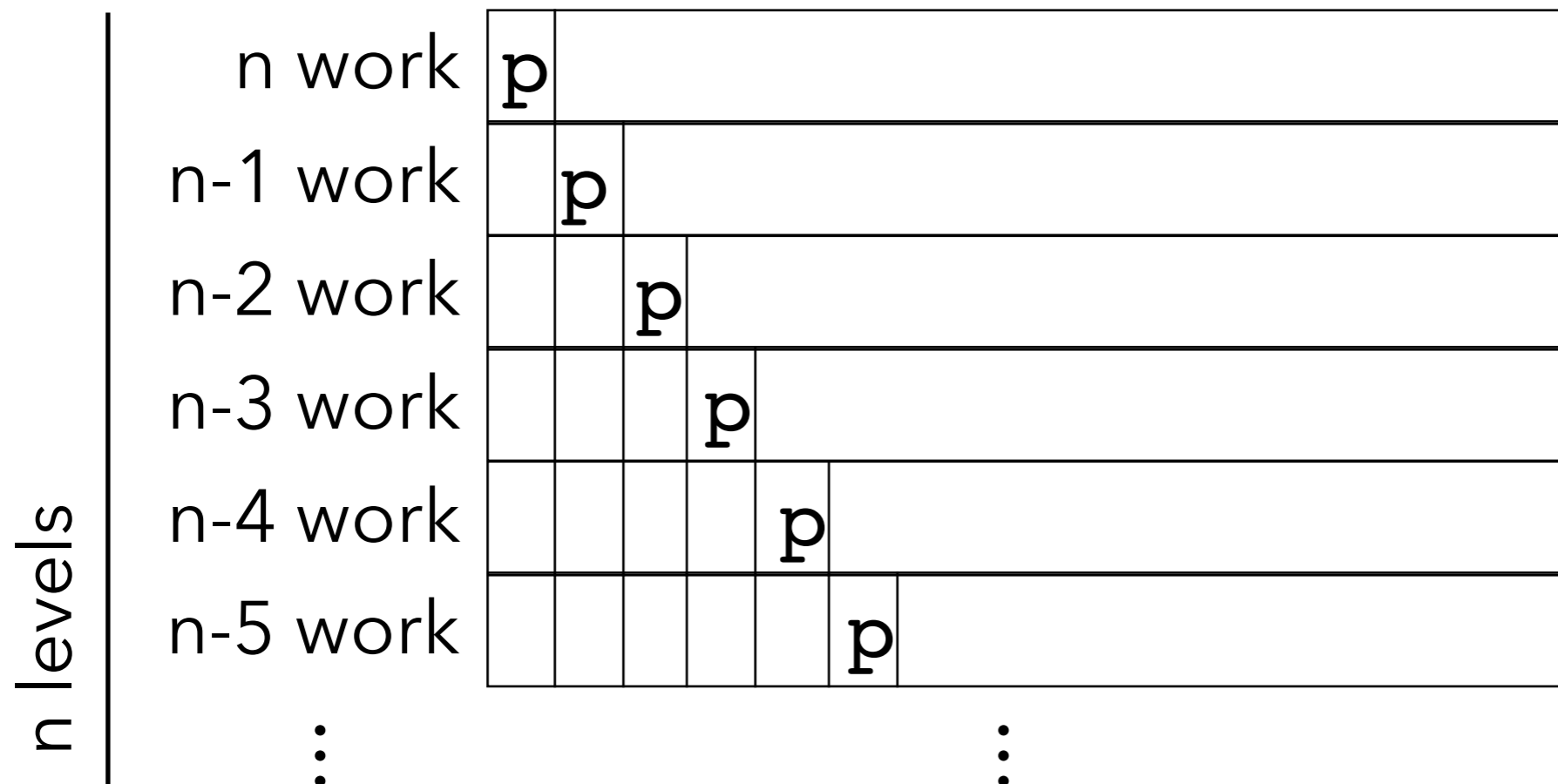


Best-case runtime: $O(n \log n)$

Runtime: Worst case

Worst case:

- pivot is the minimum or maximum of the array
- partition splits the array into 1 and $n-1$.

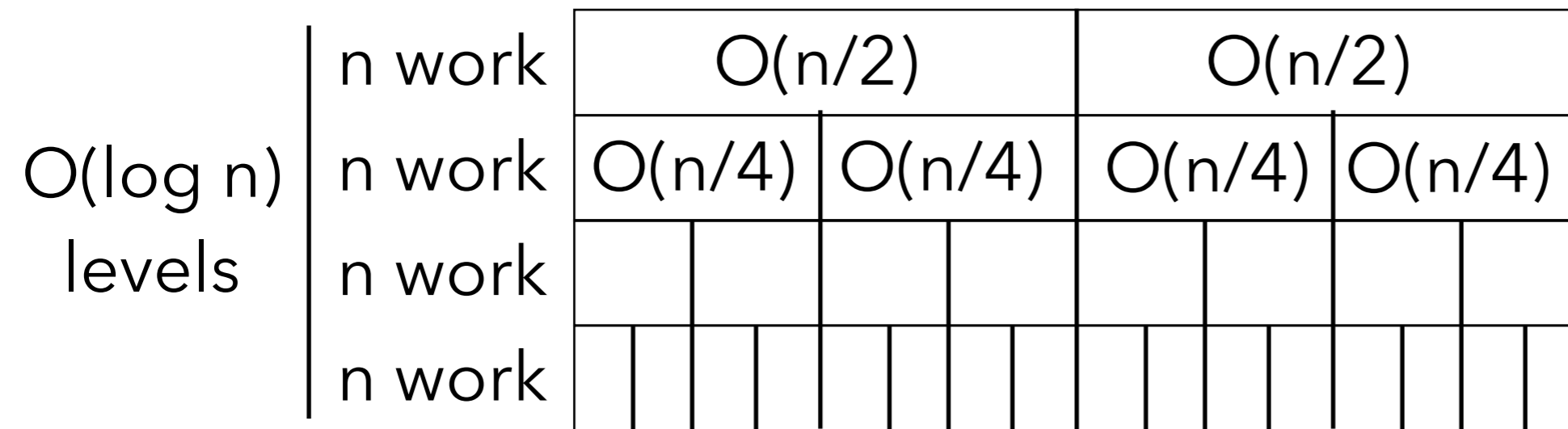


Worst-case runtime: $O(n^2)$

Runtime: Average case

Average case:

- more like best case than worst case (this is rare!)
- full analysis is out of scope, but you should know this result



Average-case runtime: $O(n \log n)$