

CSCI 241

Lecture 11 Binary Search Trees

please get logged into Socrative now! socrative.com room name: CSCI241

Announcements

Goals

- Know how to perform (and code) three tree traversals: preorder, in-order, and post-order.
- Know the definition and uses of a binary search tree.
- Be prepared to implement, and know the runtime of, the following BST operations:
 - searching
 - inserting
 - deleting
- Know what a balanced BST is and why we want it.

Tree Terminology

M is the **root** of this tree

N is the **left child** of P

S is the **right child** of P

P is the parent of N

G is the **root** of the **left subtree** of M

B, H, J, N, S are leaves

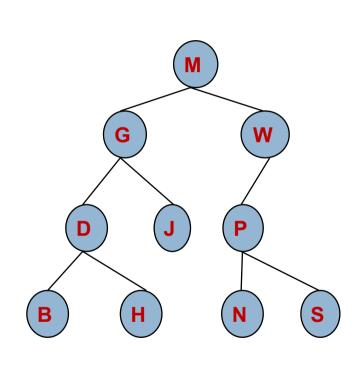
M and G are ancestors of D

P, N, S are **descendants** of W

J is at depth 2

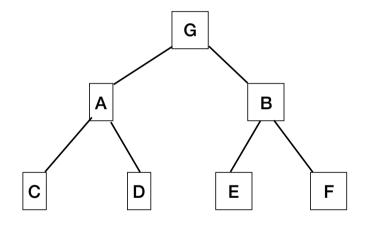
The subtree rooted at W has height 2

A collection of several trees is called a forest.





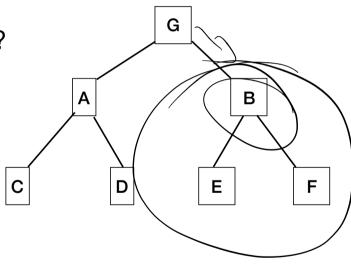
ABCD (name the node!):





ABCD (name the node!):

What's the root of G's right subtree?

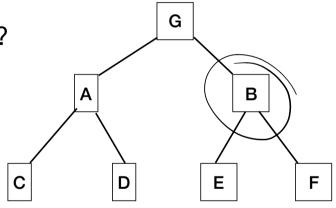




ABCD (name the node!):

What's the root of G's right subtree?

What's an ancestor of F?



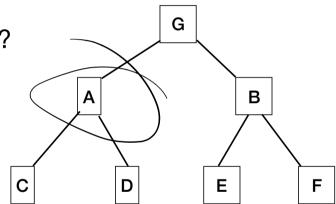


ABCD (name the node!):

What's the root of G's right subtree?

What's an ancestor of F?

What's C's parent?





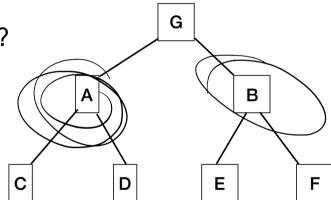
ABCD (name the node!):

What's the root of G's right subtree?

What's an ancestor of F?

What's C's parent?

What's a node at depth 1?





ABCD (name the node!):

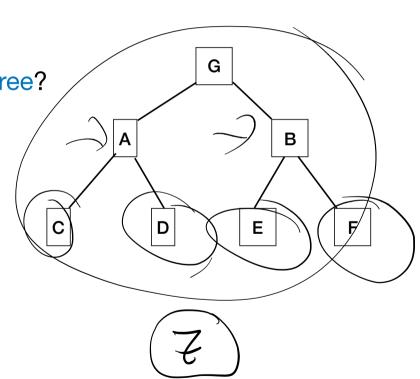
What's the root of G's right subtree?

What's an ancestor of F?

What's C's parent?

What's a node at depth 1?

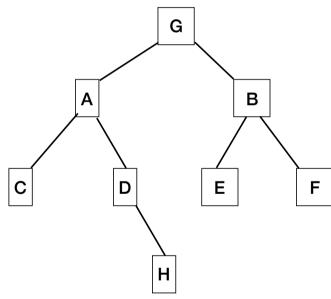
What's a node at the root of a subtree of height 0?





ABCD:

What's the height of the tree rooted at G?





ABCD:

What's the height of the tree rooted

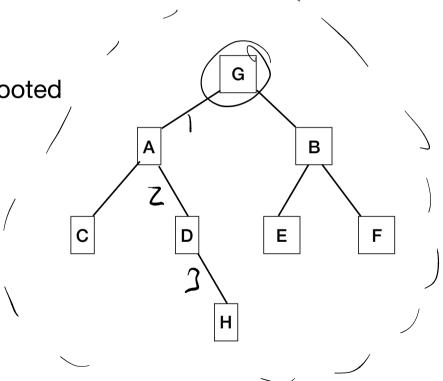
at G?

A. 1

B. 2

C. 3

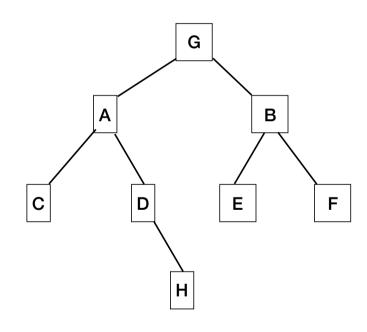
1). 4





ABCD:

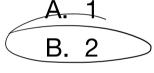
What's the depth of node D?





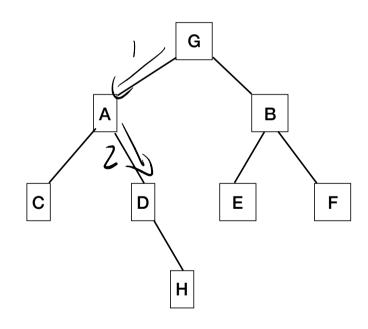
ABCD:

What's the depth of node D?



C. 3

D. 4



Print (or otherwise process) every node in a tree:

- A binary tree is
 - Empty, or
 - Three things:
 - value
 - a left binary tree
 - a right binary tree

Print (or otherwise process) every node in a tree:

A binary tree is

Print all nodes in a binary tree:

boolean printTree(Tree t):

Empty, or_

value

(base case - nothing to print)

if t == null:

Three things:

> (print this node's value)

System.out.println(t.value)

a left binary tree

(recursive call - print left subtree)

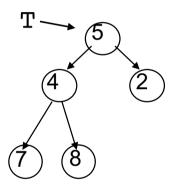
printTree(t.left)

a right binary tree

(recursive call - print left subtree)

printTree(t.right)

Print (or otherwise process) every node in a tree:



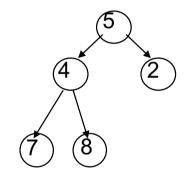
```
Print all nodes in a binary tree:
boolean printTree(Tree t):
    (base case - nothing to print)
    if t == null:
        return
```

```
(print this node's value)
System.out.println(t.value)

(recursive call - print left subtree)
printTree(t.left)
(recursive call - print left subtree)
printTree(t.right)
```

Practice Exercise

- Write the values printed by a:
 - pre-order
 - in-order
 - post-order



traversal of this (or any other) binary tree.

"Walking" over the whole tree is called a tree traversal This is done often enough that there are standard names. Previous example was a pre-order traversal:

- 1. Process root
- 2. Process left subtree
- 3. Process right subtree

Other common traversals:

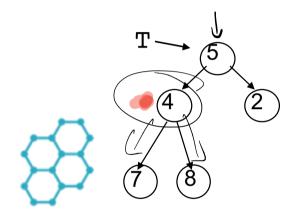
in-order traversal:

- 3. Process right subtree 3. Process root

post-order traversal:

- Process left subtree ←
 Process left subtree
- 2. Process root 2. Process right subtree

Print (or otherwise process) every node in a tree:



Print all nodes in a binary tree:

boolean preorder(Tree t):

(base case - nothing to print)

if t == null:

return

ABCD: T is a reference to the node with value 5. What is printed

by the call preorder (T)?

A. 54278

B. 74852

C. 78425

(print this node's value)
System.out.println(t.value)

(recursive call - print left subtree)

printTree(t.left)

(recursive call - print left subtree)

printTree(t.right)

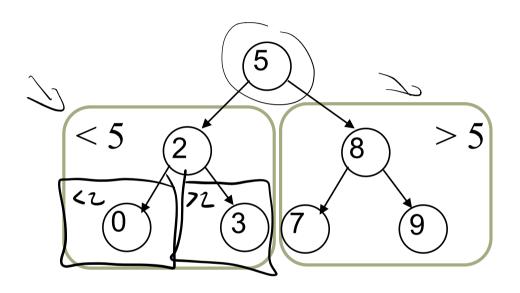
Binary Tree

```
public class Tree {
  int value;
  Tree parent;
  Tree left;
  Tree right;
}
```

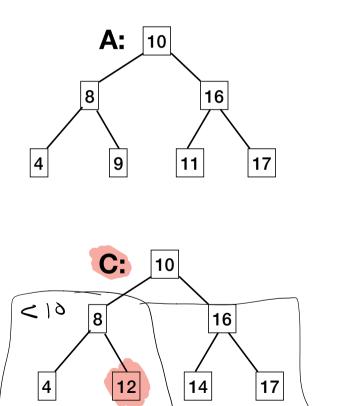
```
/** BST: a binary tree, in which:
 * -all values in left are < value
 * -all values in right are > value
 * -left and right are BSTs */
public class BST {
 int value;
 BST parent;
 BST left; < Value
 BST right; > value
```

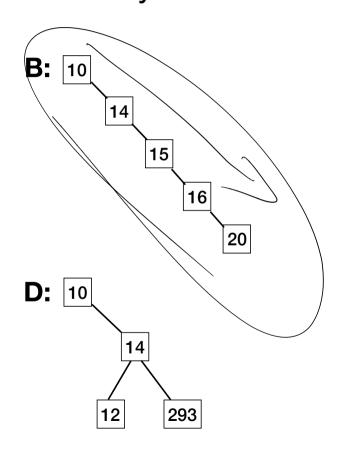
```
/** BST: a binary tree, in which:
 * -all values in left are < value
 * -all values in right are > value
 * -left and right are BSTs */
public class BST {
  int value;
  BST parent;
  BST left;
  BST right;
```

```
/** BST: a binary tree, in which:
 * -all values in left are(<)value
 * -all values in right are > value
 * -left and right are BST/s/
public class BST {
  int value;
  BST parent;
                  consequence: no duplicates!
  BST left;
  BST right;
```



ABCD: Which of these is **not** a binary search tree?





Traversing a BST

pre-order traversal:

- Process root
- 2. Process left subtree
- 3. Process right subtree

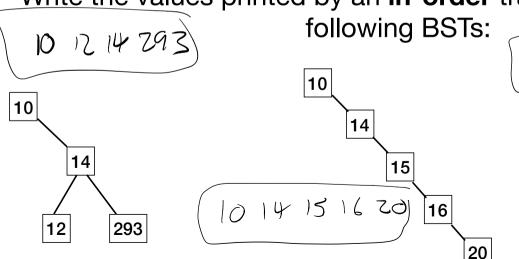
in-order traversal:

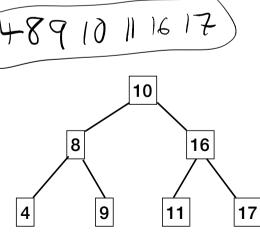
- 1. Process left subtree
- 2. Process root
- 3. Process right subtree

post-order traversal:

- 1. Process left subtree
- 2. Process right subtree
- 3. Process root

Write the values printed by an **in-order** traversal of each of the following BSTs:





(not Search!)

Searching a Binary Tree

A binary tree is

(not BST!) Find v in a binary tree:

• Empty, or

(base case - not found!)

• Three things:

value

(base case - is this v?)

• a left binary tree

(recursive call - is v in left?)

• a right binary tree

(recursive call - is v in right?)

(not Search!)

Searching a Binary Tree

A binary tree is

• Empty, or

Three things:

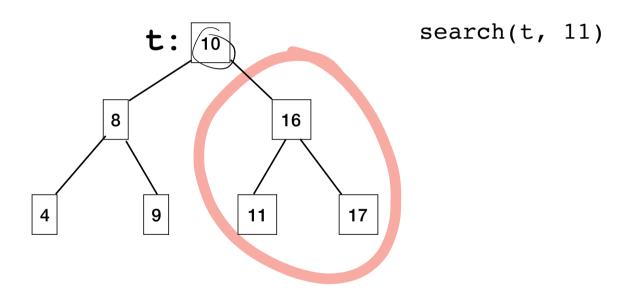
value

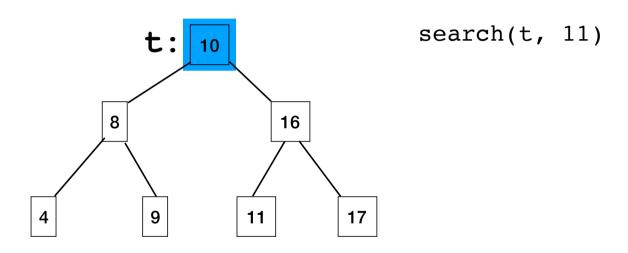
- a left binary tree
- a right binary tree

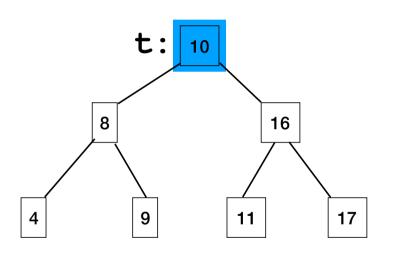
```
(not BST!)
Find v in a binary tree:
boolean findVal(Tree t, int v):
```

```
(base case - not found!)
if t == null:
    return false
```

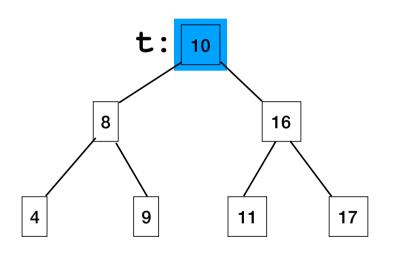
```
(base case - is this v?)
if t.value == v: return true
```



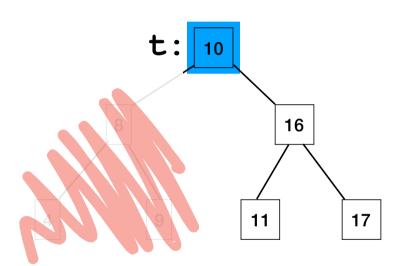




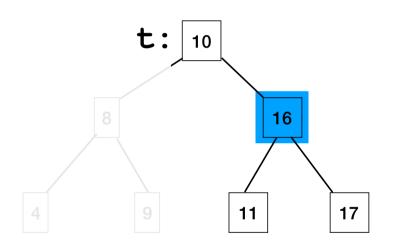
search(t, 11) 11 > 10



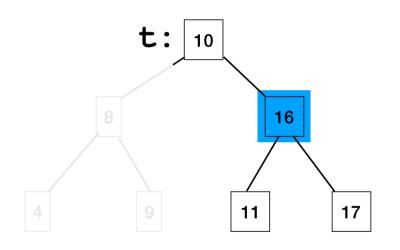
```
search(t, 11)
11 > 10
search(right, 11)
```



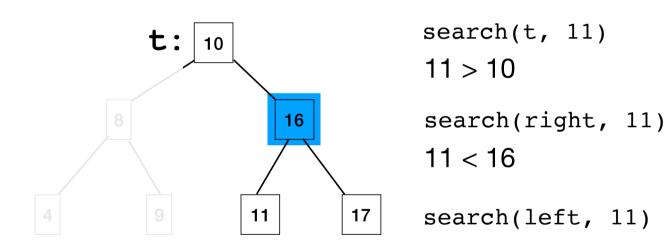
```
search(t, 11)
11 > 10
search(right, 11)
```

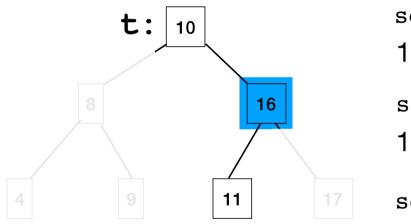


```
search(t, 11)
11 > 10
search(right, 11)
```

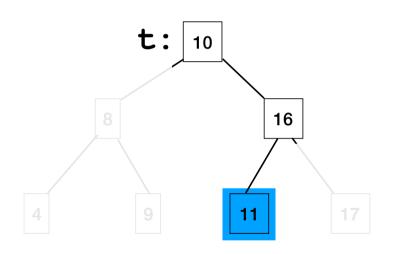


```
search(t, 11)
11 > 10
search(right, 11)
11 < 16</pre>
```

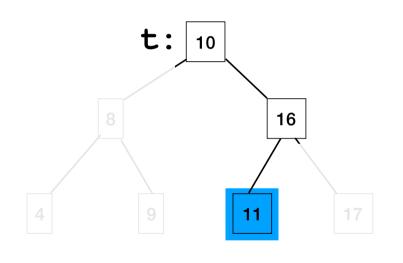




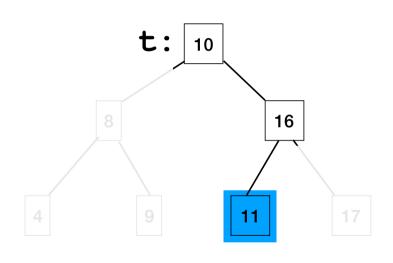
```
search(t, 11)
11 > 10
search(right, 11)
11 < 16
search(left, 11)</pre>
```



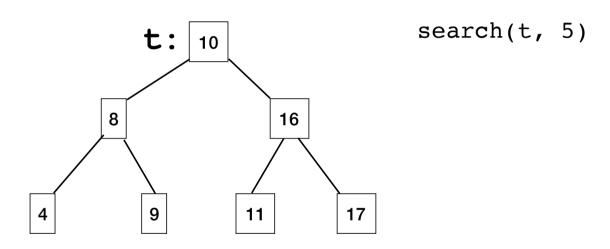
```
search(t, 11)
11 > 10
search(right, 11)
11 < 16
search(left, 11)</pre>
```

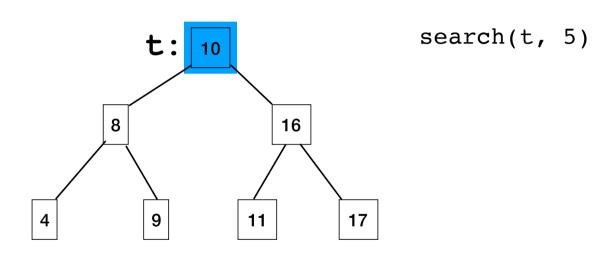


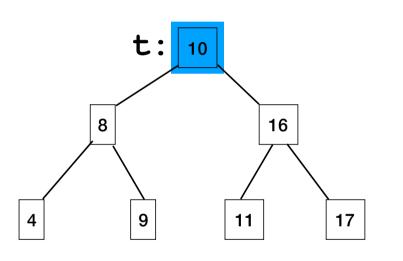
```
search(t, 11)
11 > 10
search(right, 11)
11 < 16
search(left, 11)
11 == 11</pre>
```



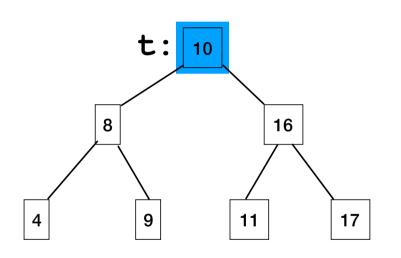
```
search(t, 11)
11 > 10
search(right, 11)
11 < 16
search(left, 11)
11 == 11
found it! return.</pre>
```



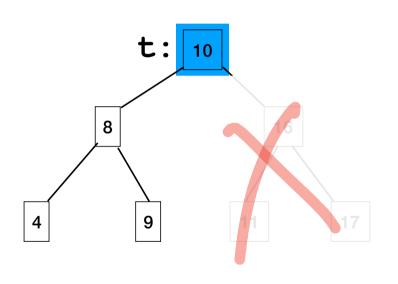




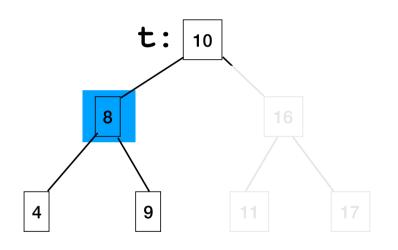
search(t, 5) 5 < 10



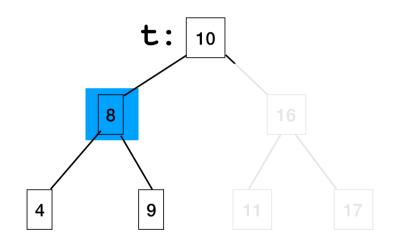
```
search(t, 5)
5 < 10
search(left, 5)</pre>
```



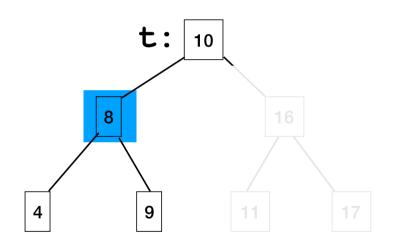
```
search(t, 5)
5 < 10
search(left, 5)</pre>
```



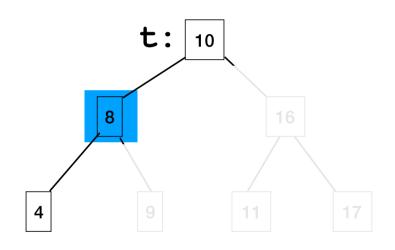
```
search(t, 5)
5 < 10
search(left, 5)</pre>
```



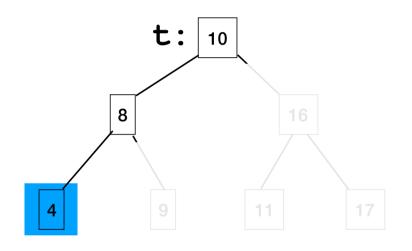
```
search(t, 5)
5 < 10
search(left, 5)
5 < 8</pre>
```



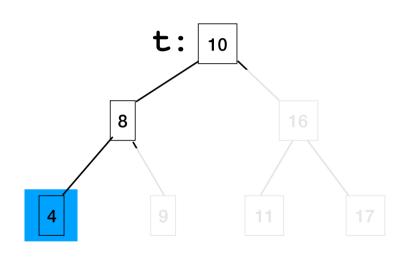
```
search(t, 5)
5 < 10
search(left, 5)
5 < 8
search(left, 5)</pre>
```



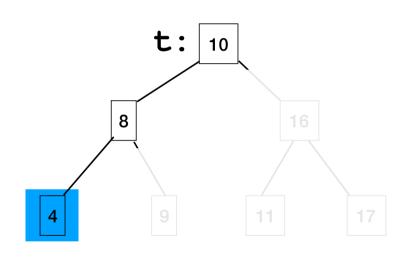
```
search(t, 5)
5 < 10
search(left, 5)
5 < 8
search(left, 5)</pre>
```



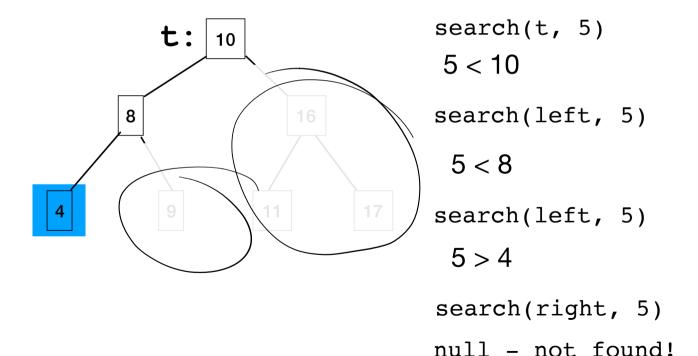
```
search(t, 5)
5 < 10
search(left, 5)
5 < 8
search(left, 5)</pre>
```



```
search(t, 5)
5 < 10
search(left, 5)
5 < 8
search(left, 5)
5 > 4
```



```
search(t, 5)
5 < 10
search(left, 5)
 5 < 8
search(left, 5)
 5 > 4
search(right, 5)
```



Searching: BT vs BST

```
/** Searches the binary tree
                                /** Searches the binary *search*
 * rooted at n for value v,
                                 * tree rooted at n for value v,
 * returning true iff it is
                                 * returning true iff it is in
 * in the tree. */
                                 * the tree. */
boolean srchBT(n, v) {
                                public srchBST(n, v) {
  if (n == null) {
                                  if (n == null) {
    return false;
                                    return false;
  if (n.v == v) 
                                  if (n.v == v) 
    return true;
                                    return true;
  return srchBT(n.left, v)
                                  if (v < n.v) 
         srchBT(n.right, v);
                                    return srchBST(n.left, v);
                                  } else {
                                    return srchBST(n.right, v)
```

```
boolean search (BST)t, int v):
  if t == null:
                                           10
    return false
  if t.value == v:
    return true
  if v < t.value:</pre>
 return search(t.left)
  else:
                                              11
                                                      17
 return search(t.right)
```

```
boolean search(BST t, int v):
   if t == null:
      return false
   if t.value == v:
      return true
   if v < t.value:
      return search(t.left)
   else:
      return search(t.right)</pre>
```

If h is the tree's **height**, search can visit at most h+1 nodes!

Runtime of search is **O(h)**.

```
boolean search(BST t, int v):
    if t == null:
        return false
    if t.value == v:
        return true
    if v < t.value:
        return search(t.left)
    else:
        return search(t.right)</pre>
```

If h is the tree's height, search can visit at most h+1 nodes!

Runtime of search is **O(h)**.

That's great, but how does h relate to n, the number of nodes?

nodes. What's The Suppose me have n Min height? Max height?

Consider h = 2:

depth

0

1

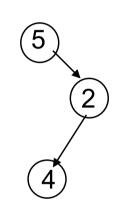
1

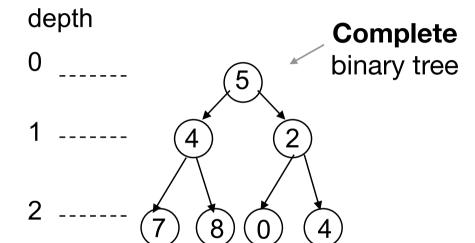
2

1

2

Consider h = 2:

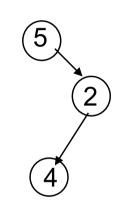


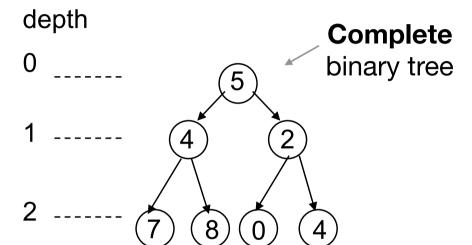


Fewest possible:

n = h+1n is O(h)

Consider h = 2:

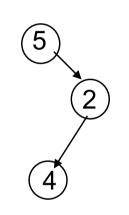




Fewest possible:

- n = h+1
- n is O(h)
- h is O(n)

Consider h = 2:

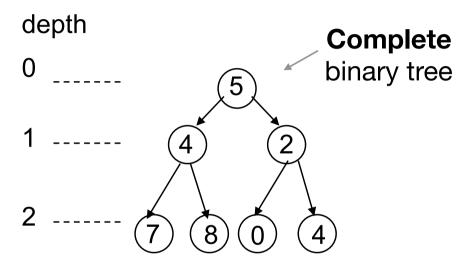


Fewest possible:

$$n = h+1$$

n is O(h)

h is O(n)



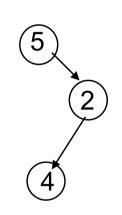
Most possible:

At depth d: 2^d nodes possible.

At all depths: $2^0 + 2^1 + ... + 2^h$

$$= 2^{h+1} - 1$$

Consider h = 2:

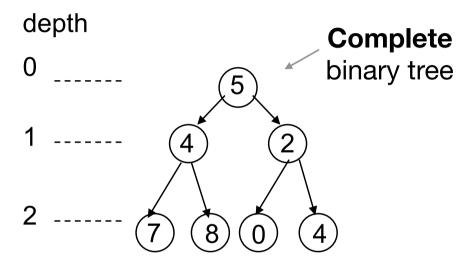


Fewest possible:

$$n = h+1$$

n is O(h)

h is O(n)



Most possible:

At depth d: 2^d nodes possible.

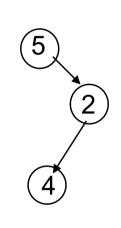
At all depths: $2^0 + 2^1 + ... + 2^h$

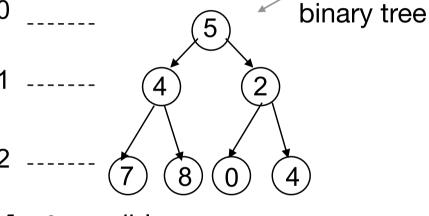
$$= 2^{h+1} - 1$$

 $n = 2^{h+1} - 1$

depth

Consider h = 2:





Complete

Fewest possible: n = h+1

h is O(n)

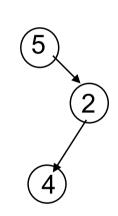
Most possible:

At depth d: 2^d nodes possible.

At all depths: $2^0 + 2^1 + ... + 2^h$ $= 2^{h+1} - 1$

 $n = 2^{h+1} - 1$ n is $O(2^h)$



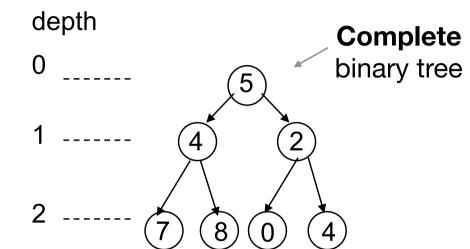


Fewest possible:

$$n = h+1$$

n is O(h)

h is O(n)



Most possible:

At depth d: 2^d nodes possible.

At all depths: $2^0 + 2^1 + ... + 2^h$ = $2^{h+1} - 1$

 $n = 2^{h+1} - 1$

n is $O(2^h)$

h is O(log n)

```
boolean search(BST t, int v):
  if t == null:
                          10
    return false
  if t.value == v:
                              14
                                                     10
    return true
                                 15
  if t.value < v:</pre>
                                                          16
    return search(t.left)
                                    16
  else:
                                                  9
                                       20
                                                        11
    return search(t.right)
```

Runtime of search is O(h). Worst: O(n) Best: O(log n)

```
We want our trees to
boolean search(BST t, int v):
                                     look more like this
  if t == null:
                          10
    return false
  if t.value == v:
                              14
                                                      10
    return true
                                 15
  if t.value < v:</pre>
                                                           16
    return search(t.left)
                                    16
  else:
                                                   9
                                        20
                                                         11
    return search(t.right)
```

Runtime of search is O(h). Worst: O(n) Best: O(log n)