

CSCI 241

Lecture 7

~~Quicksort~~

Stability; Non-Comparison Sorts

Radix Sort

Announcements

- Quiz 1 grades and review video out soon

Goals:

- Know what it means for a sorting algorithm to be **stable**
- Understand the distinction between **comparison** and **non-comparison** sorts.
- Be prepared to implement **radix sort**.
- Know the definition of an **in-place** sorting algorithm.

Stability

Objects can be sorted on **keys** - **different** objects may have the same value.

A **stable** sort maintains the order of distinct elements with the same key.

- Example: sort the following list on the **tens** place only:

Sorted stably:

unstably:

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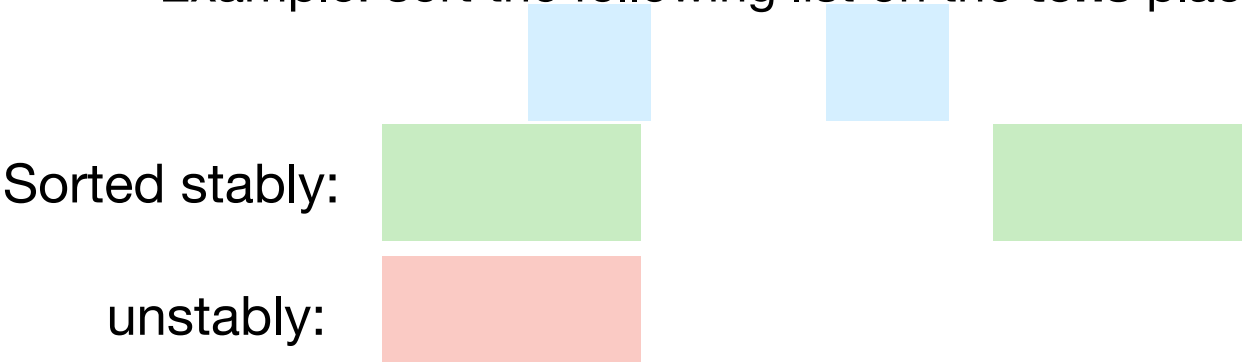
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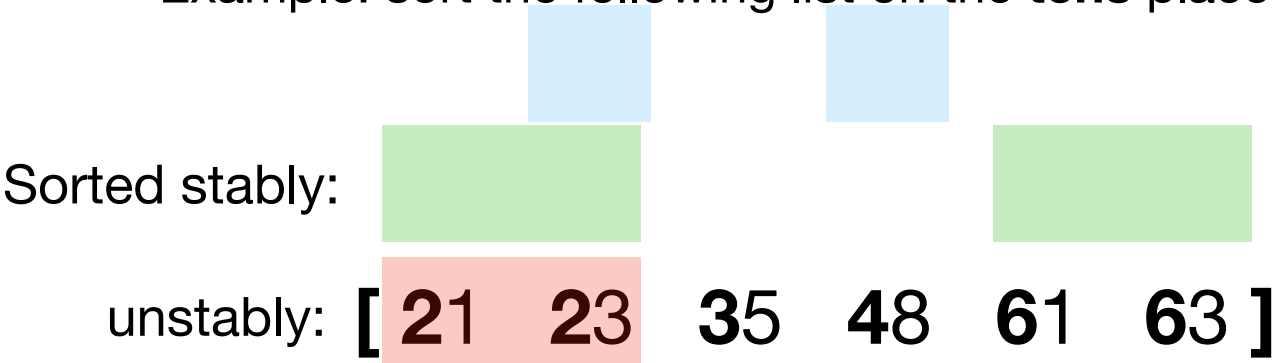


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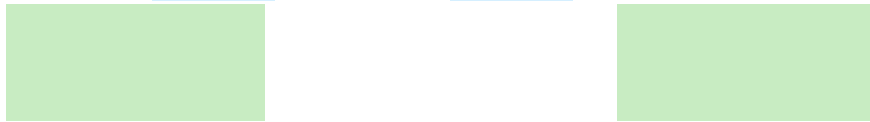
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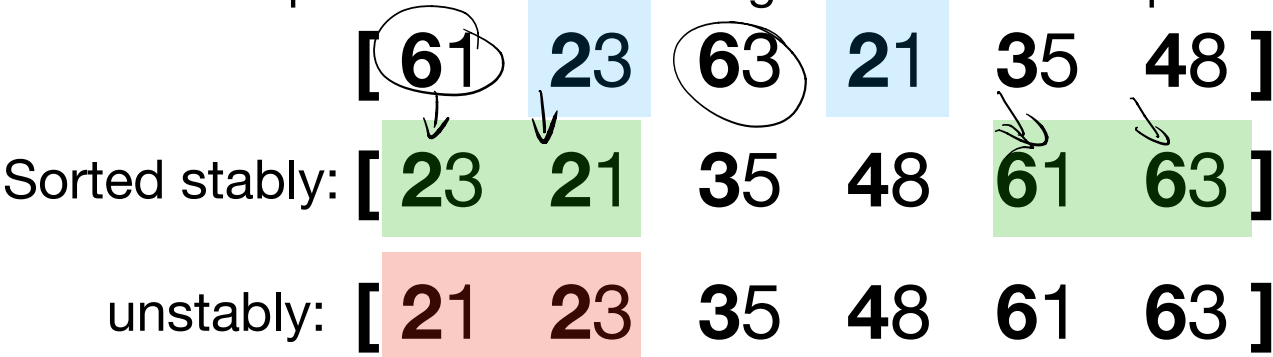
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Comparison sorts operate by comparing pairs of elements.

Examples: all four sorts we've seen so far!

...is there any other way to do it?

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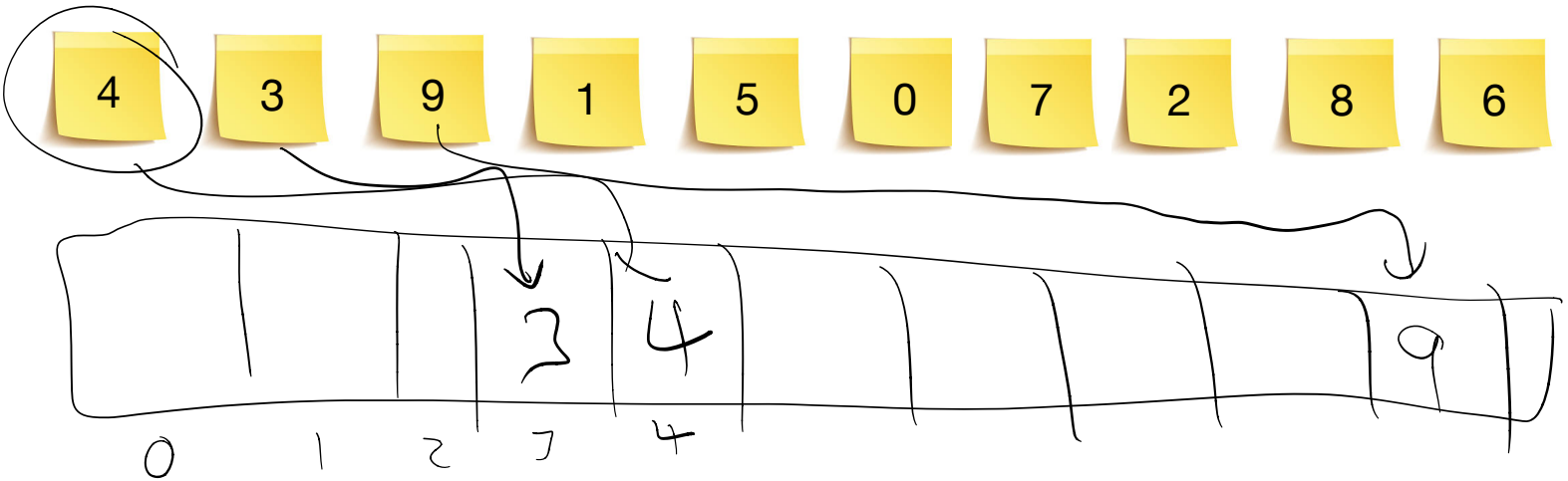
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**How do you sort without
comparing elements?**

How do you sort things without comparing them?

Suppose I gave you 10 sticky notes with the digits 0 through 9.
What algorithm would you use to sort them?



How do you sort things without comparing them?

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How many times did you need to look at each sticky note?

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How many times did you need to look at each sticky note?

What if there are duplicates?

LSD Radix Sort


```
/** least significant digit radix sort A */  
LSDRadixSort(A):
```

→ `max_digits = max # digits in any element of A`

```
for d in 0..max_digits:
```

```
    stably sort A on the dth least significant  
    digit
```

```
// A is now sorted(!)
```

- 
1. ones place, then
 2. tens place, then
 3. hundreds place,
and so on
- ↓

Do you believe me?

```
/** least significant digit radix sort A */  
LSDRadixSort(A):  
  2= max_digits = max # digits in any element of A  
  for d in 0..max_digits:  
    ↷ stably sort A on the dth least significant  
      digit
```

```
// A is now sorted(!)    [45, 26, 42, 32] 07
```

↓	↓	↓	↓
42	32	45	26
26	32	42	45

Do you believe me?

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    digit  
  
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```

Still don't believe me? <https://visualgo.net/en/sorting>

LSD Radix Sort using queue buckets

Pseudocode from visualgo.net:

LSDRadixSort(A):

 create 10 buckets (queues) for each digit (0 to 9)

 for each digit (least- to most-significant):

 for each element in A:

 move element into its bucket based on digit

 for each bucket, starting from smallest digit

 while bucket is non-empty

 restore element to list

LSD Radix Sort using queue buckets

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```
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```
    for each element in A:
```

```
      move element into its bucket based on digit
```

```
    for each bucket, starting from smallest digit
```

```
      while bucket is non-empty
```

```
        restore element to list
```

LSD Intuition: sort on most-significant digit **last**; if tied, yield to the next most significant digit, and so on.

Only works because **stability** preserves orderings from less significant digits (previously sorted).

Exercise: Radix sort this

[7, 19, 21, 11, 14, 54, 1, 8]

Hint: [07, 19, 21, 11, 14, 54, 01, 08]

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create 10 buckets (queues) for each digit (0 to 9)

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Exercise: Radix sort this

[07, 19, 61, 11, 14, 54, 01, 08]



Buckets
on 1's place:

0	1	2	3	4	5	6	7	8	9
	01								
	11			54					
	61			14			07	08	19

Sorted on
1's place:

61 11⁰¹ 14 54 07 08 19

Buckets
on 10's place:

0	1	2	3	4	5	6	7	8	9
08	19								
07	14								
01	11			54	61				

Sorted on
10's place:

01 07 08 11 14 19 54 61

LSD Radix Sort using counting sort

```
/** least significant digit radix sort A */  
LSDRadixSort(A):  
max_digits = max # digits in any element of A  
for d in 0..max_digits:  
    counting sort A on the dth least  
    significant digit  
  
// A is now sorted(!)
```

Counting Sort

Formalizes what you did with the 0-9 sticky notes:

- Handles duplicates
- Stable sort
- Less memory overhead than queue buckets

Intuition:

<http://www.cs.miami.edu/home/burt/learning/Csc517.091/workbook/countingsort.html>

Pseudocode in CLRS (reproduced on the next slide).

Counting Sort - from CLRS

Notes:

- k is the base or radix (10 in our examples)
- B is filled with the sorted values from A .
- C maintains counts for each bucket.
- The final loop **must** go back-to-front to guarantee stability.

COUNTING-SORT(A, B, k)

```
1  let  $C[0..k]$  be a new array
2  for  $i = 0$  to  $k$ 
3       $C[i] = 0$ 
4  for  $j = 1$  to  $A.length$ 
5       $C[A[j]] = C[A[j]] + 1$ 
6  //  $C[i]$  now contains the number of elements equal to  $i$ .
7  for  $i = 1$  to  $k$ 
8       $C[i] = C[i] + C[i - 1]$ 
9  //  $C[i]$  now contains the number of elements less than or equal to  $i$ .
10 for  $j = A.length$  downto 1
11      $B[C[A[j]]] = A[j]$ 
12      $C[A[j]] = C[A[j]] - 1$ 
```

in-place sorts

One more property of sorting algorithms aside from runtime.

A sorting algorithm is considered **in-place** if:

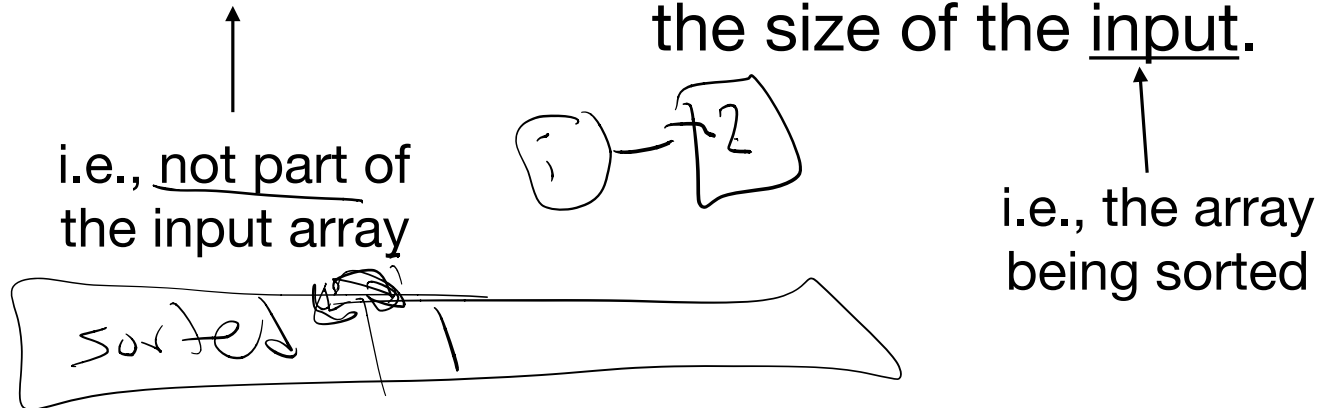
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	in-place?
insertion	Y
selection	Y
merge	?
quick	?
radix	?