Α	0110	J	110000000	s	1001
в	00010	ĸ	101000	Т	1000
С	01111	L	10101	U	0000
D	0010	м	00011	v	1100001
E	111	N	010	w	001101
F	101001	0	01110	х	1100000101
G	11001	Р	. 110001	Y	110000011
н	00111	Q	1100000100	z	001100
I	1101	R	1011		

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## **CSCI 241**

Lecture N Huffman Coding

# **Course Evaluations**

 Please spend the few minutes to provide feedback!

### More sources of fun/practice

- Easier: <u>https://codingbat.com/java</u>
  - good for recursion
- Easy -> Hard <u>https://adventofcode.com/</u>
  - see how far into December you can get before giving up
  - A few years ago I got to ~Dec 21
- Easy -> Nuts <a href="https://projecteuler.net/">https://projecteuler.net/</a>
  - first 50 are pretty manageable
- <u>https://leetcode.com/</u>
  - I have no personal experience, but I know many use it



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What gives?

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### Image Compression: JPEG



## Video Compression: H.265



**HEVC Encoder Blocks** 

## Video Compression: H.265



**HEVC Encoder Blocks** 

### Image Compression: JPEG



#### PROCEEDINGS OF THE I.R.E.

### A Method for the Construction of Minimum-Redundancy Codes\* DAVID A. HUFFMAN<sup>+</sup>, ASSOCIATE, IRE

Summary—An optimum method of coding an ensemble of messages consisting of a finite number of members is developed. A minimum-redundancy code is one constructed in such a way that the average number of coding digits per message is minimized.

#### INTRODUCTION

NE IMPORTANT METHOD of transmitting messages is to transmit in their place sequences of symbols. If there are more messages which might be sent than there are kinds of symbols available, then some of the messages must use more than one symbol. If it is assumed that each symbol requires the same time for transmission, then the time for transmission (length) of a message is directly proportional to the number of symbols associated with it. In this paper, the symbol or sequence of symbols associated with a given message will be called the "message code." The entire number of messages which might be transmitted will be will be defined here as an ensemble code which, for a message ensemble consisting of a finite number of members, N, and for a given number of coding digits, D, yields the lowest possible average message length. In order to avoid the use of the lengthy term "minimum-redundancy," this term will be replaced here by "optimum." It will be understood then that, in this paper, "optimum code" means "minimum-redundancy code."

The following basic restrictions will be imposed on an ensemble code:

- (a) No two messages will consist of identical arrangements of coding digits.
- (b) The message codes will be constructed in such a way that no additional indication is necessary to specify where a message code begins and ends once the starting point of a sequence of messages is known.

# This is a coding tree.

Encodes a **map**ping from bit strings to words: 0 means go left 1 means go right

0: a

111: d

101: b



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Key intuition: put common words near the root.

	a	b	С	d	е	f
Frequency (in thousands)	45	13	12	16	9	5

HUFFMAN(C)

1 n = |C|

$$2 \quad Q = C$$

4

3 for i = 1 to n - 1

allocate a new node z

5 
$$z.left = x = EXTRACT-MIN(Q)$$

6 
$$z.right = y = EXTRACT-MIN(Q)$$

7 
$$z.freq = x.freq + y.freq$$

- 8 INSERT(Q, z)
- 9 return EXTRACT-MIN(Q) // return the root of the tree

	a	b	С	d	е	f
Frequency (in thousands)	45	13	12	16	9	5
Variable-length codeword	0	101	100	111	1101	1100













Smallest two: 12, 13



Smallest two: 12, 13 b:13 **d**:16 **c**:12 a:45 0 **f**:5 **e**:9 Smallest two: 14, 16 **d**:16 a:45 0 0 I **f**:5 b:13 **e**:9 **c**:12

0

**f**:5

**e**:9

Smallest two: 5, 9

Smallest two: 12, 13 b:13 **d**:16 **c**:12 a:45 0 **f**:5 **e**:9 Smallest two: 14, 16 **d**:16 **a**:45 0 0 **f**:5 b:13 **c**:12 **e**:9 30 **a**:45 0 b:13 **d**:16 **c**:12

**f**:5

**e**:9

Smallest two: 5, 9









#### **Huffman Tree:**



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**f**:5

**e**:9



# Coding Huffman Coding