

CSCI 241

Lecture 8:
Abstract Data Types
Introduction to Trees

Announcements

- A1: Look into the future: read the rubric!
- Submitting late (using slip days or otherwise) requires sending me email **after** you submit.

Goals:

- Understand the motivation for trees:
 - To model **tree-structured data**.
 - To implement **abstract data types**.
- Understand the definition of a tree.
- Know the basic terminology associated with trees:
 - Root, child, parent, leaf, height, depth, subtree, descendent, ancestor
- Be able to write a tree class and simple recursive methods such as size, height, and traversals.

Last Time:

Big-Deal CS Concept #1: Runtime

Big-Deal CS Concept #2: Interface vs Implementation and Abstract Data Types

What the operations do



An abstract data type specifies only **interface**,
not **implementation**

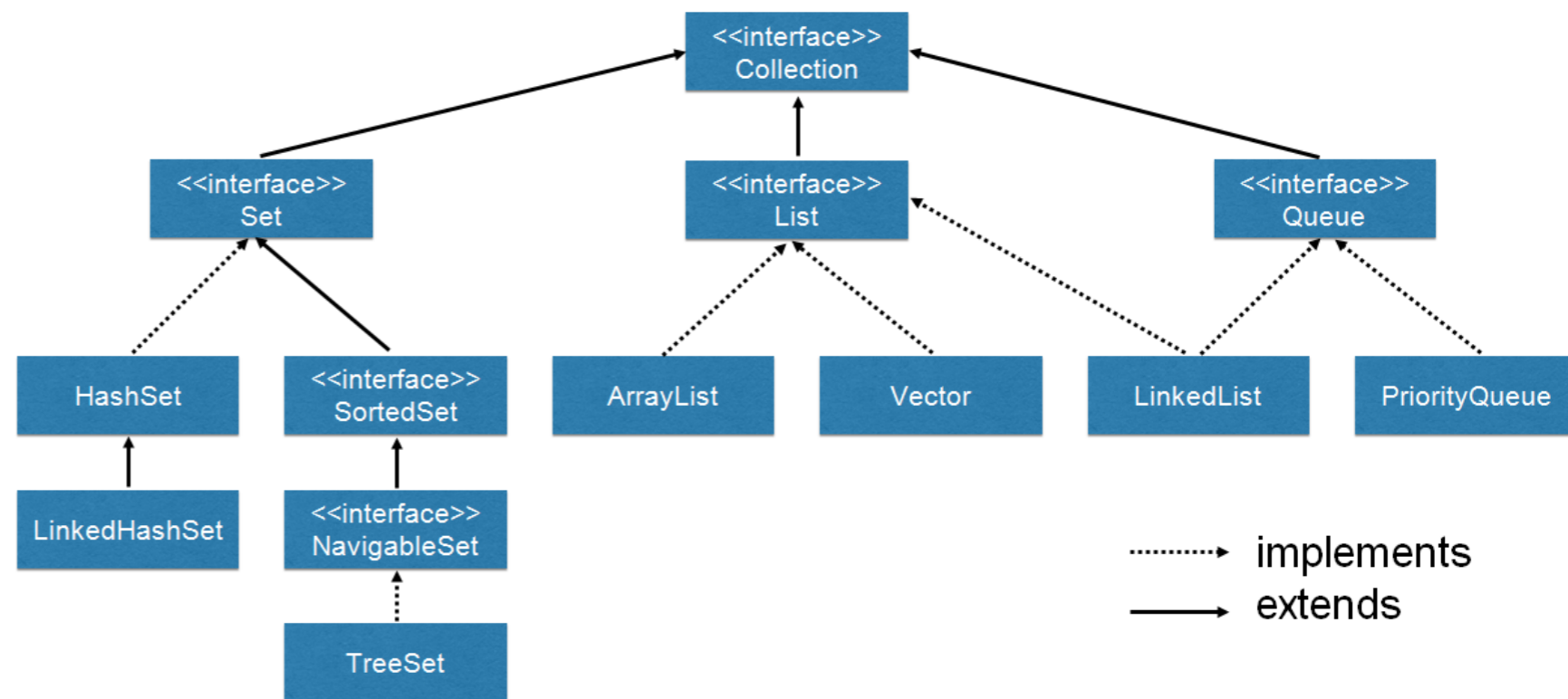


How they are accomplished

Abstract Data Types: Examples

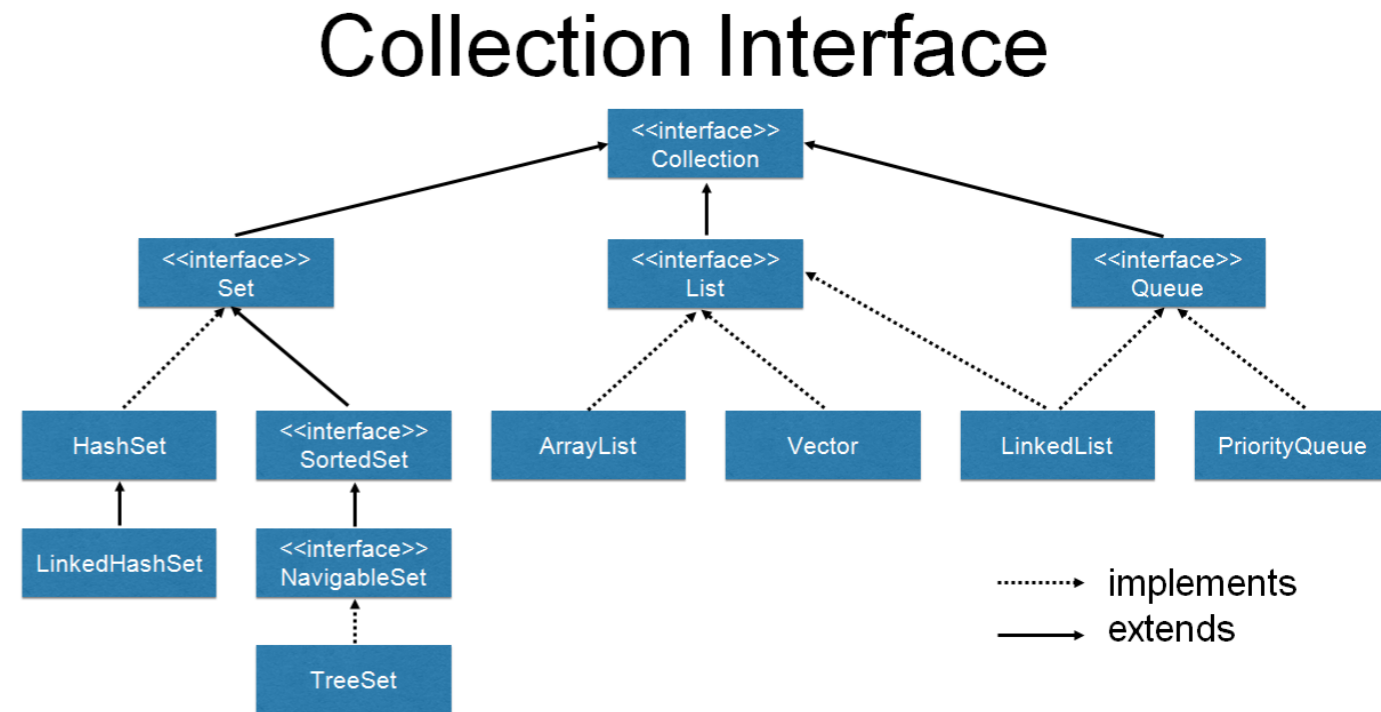
- List, Queue, Stack
- Set
- Tree
- Priority Queue
- Map
- Graph

Collection Interface

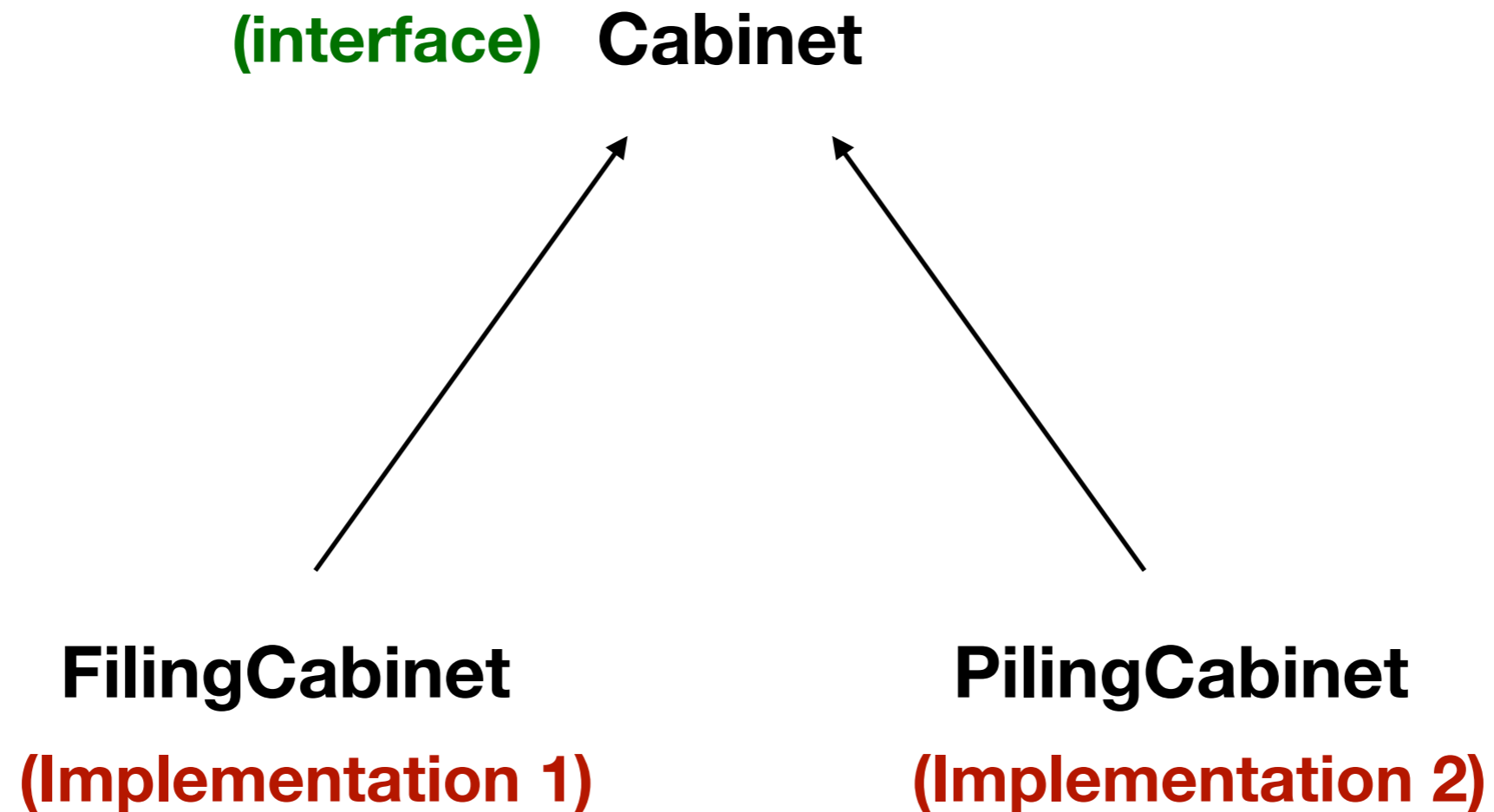


Abstract Data Types: Examples

- List, Queue, Stack (145)
- Set (Weeks 4,5,7)
- Tree (Weeks 4-6; A2)
- Priority Queue (Week 6; A3)
- Map (Week 7; A3)
- Graph (Weeks 8-9; A4)



Interface vs Implementation: Example



Interface vs Implementation: Example

Interface

Cabinet:

(short for “if and only if”)

- Contains(item) - returns true **iff** item is in the cabinet
- Add(item) - adds item to the cabinet
- Remove(item) - removes item from the cabinet if it exists

FilingCabinet implements Cabinet:

Contains(item):

`look up drawer by first letter range`

`find folder by first letter`

`search folder for item`

`return true if item is found, false otherwise`

Implementation

Comparing Implementations

class FilingCabinet:

- Contains(item):

look up drawer by first letter range

find folder by first letter

search folder for item

return true if item is found, false otherwise

class PilingCabinet:

- Contains(item):

for each drawer:

exhaustively search drawer

if found, return true

return false

Comparing Implementations

class FilingCabinet:

- Add(item):

look up drawer by first letter range

find folder by first letter

insert item into folder

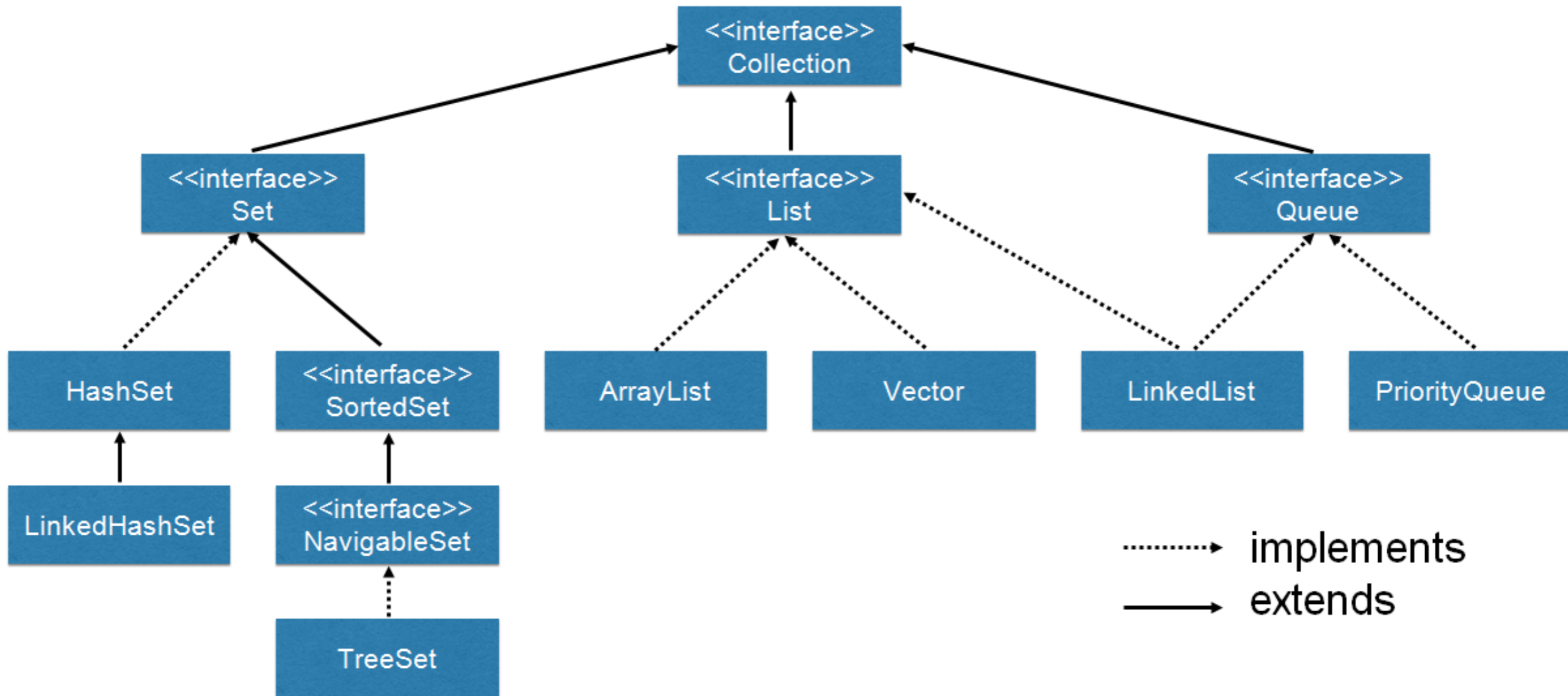
class PilingCabinet:

- Add(item):

open random drawer

insert item into drawer

Collection Interface



Is an array an ADT?

ADTs and Runtime:

Why we care

Runtime comparison of **List** implementations:

| Class: | ArrayList | LinkedList |
|------------------|-----------|---------------|
| Backing storage: | array | chained nodes |
| add(i, val) | $O(n)$ | $O(n)$ |
| add(0, val) | $O(n)$ | $O(1)$ |
| add(n, val) | $O(1)$ | $O(1)$ |
| get(i) | $O(1)$ | $O(n)$ |
| get(0) | $O(1)$ | $O(1)$ |
| get(n) | $O(1)$ | $O(1)$ |

Assume: i = arbitrary index. n = last index + 1.

Linked List

```
public class ListNode {  
    int value;  
    ListNode next;  
}
```

Linked List

```
public class List {  
    int value;  
    List next;  
}
```

The node *is the list*.

Next points to the **tail** of the list (also a list!)

Binary Tree

```
public class Tree {  
    int value;  
    Tree left;  
    Tree right;  
}
```

The node *is the tree*.

`left` points to the **left child** of the tree (also a tree!)

`right` points to the **right child** of the tree (also a tree!)

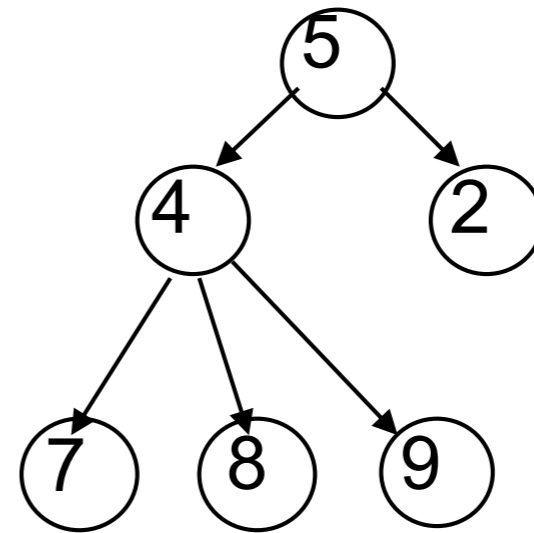
Tree - Definition

Tree: like a linked list, but:

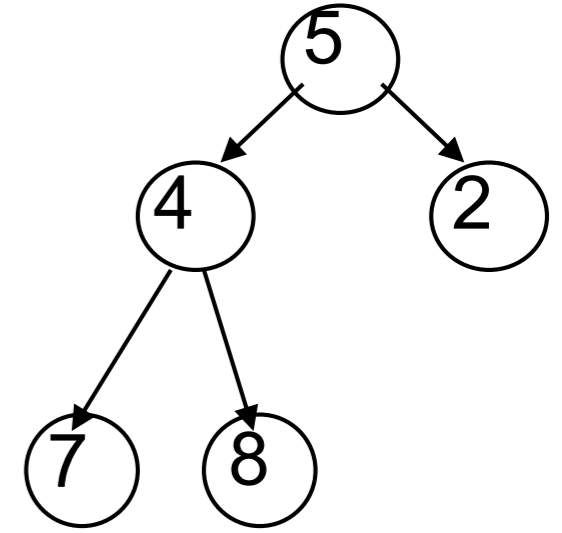
- Each node may have zero or more *successors* (**children**)
- Each node has exactly one *predecessor* (parent) except the *root*, which has none
- All nodes are reachable from *root*

Binary tree: A tree, but:

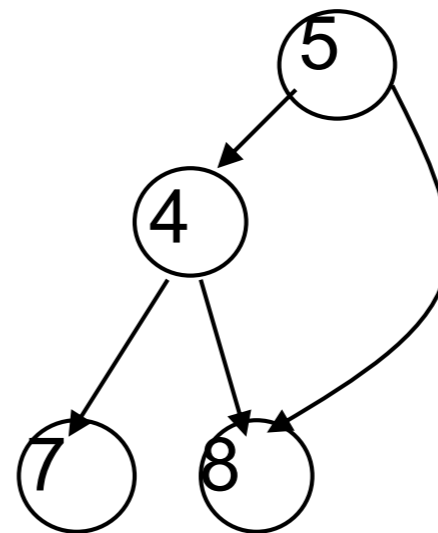
- Each can have at most **two** children (left child, right child)



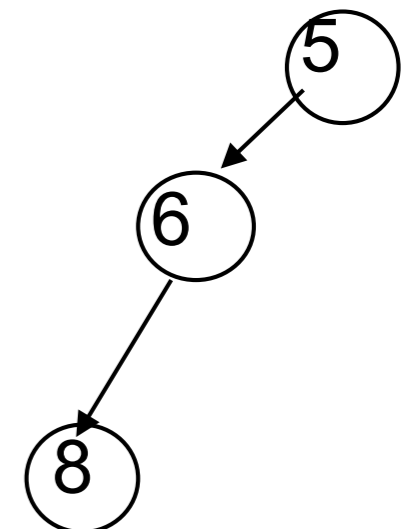
General tree



Binary tree



Not a tree



List-like tree

Tree Terminology

M is the **root** of this tree

G is the **root** of the **left subtree** of M

B, H, J, N, S are **leaves** (have no children)

N is the **left child** of P

S is the **right child** of P

P is the **parent** of N

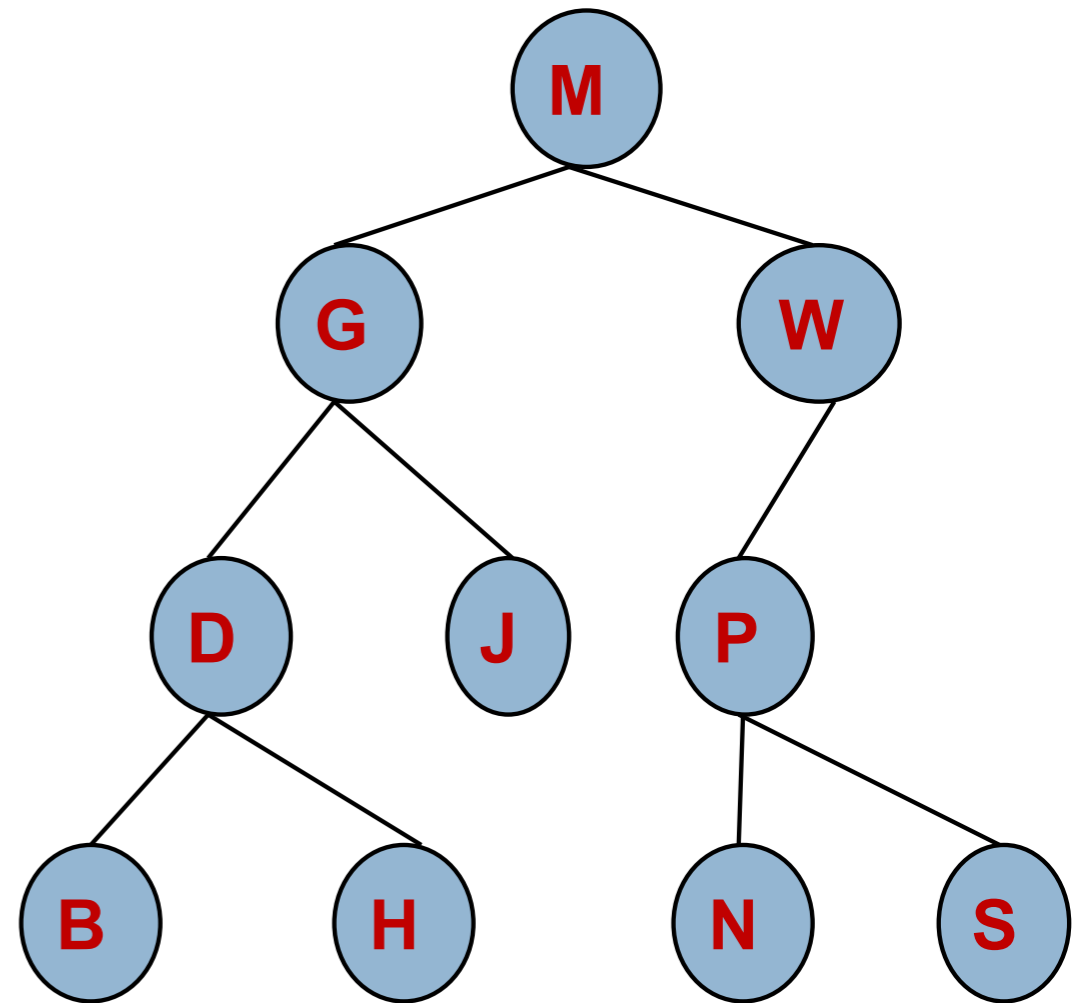
M and G are **ancestors** of D

P, N, S are **descendants** of W

J is at **depth** 2 (length of path from root)

The subtree rooted at W has **height** (length of longest path to a leaf) of 2

A collection of several trees is called a _____?



```
public class BinaryTreeNode {  
    private int value;  
    private BinaryTreeNode parent; (null if no left child)  
    private BinaryTreeNode left; // left subtree  
    private BinaryTreeNode right; // right subtree  
                                     (null if no right child)  
}
```

```
public class GeneralTreeNode {  
    private int value;  
    private GeneralTreeNode parent;  
    private List<GeneralTreeNode> children;  
}
```

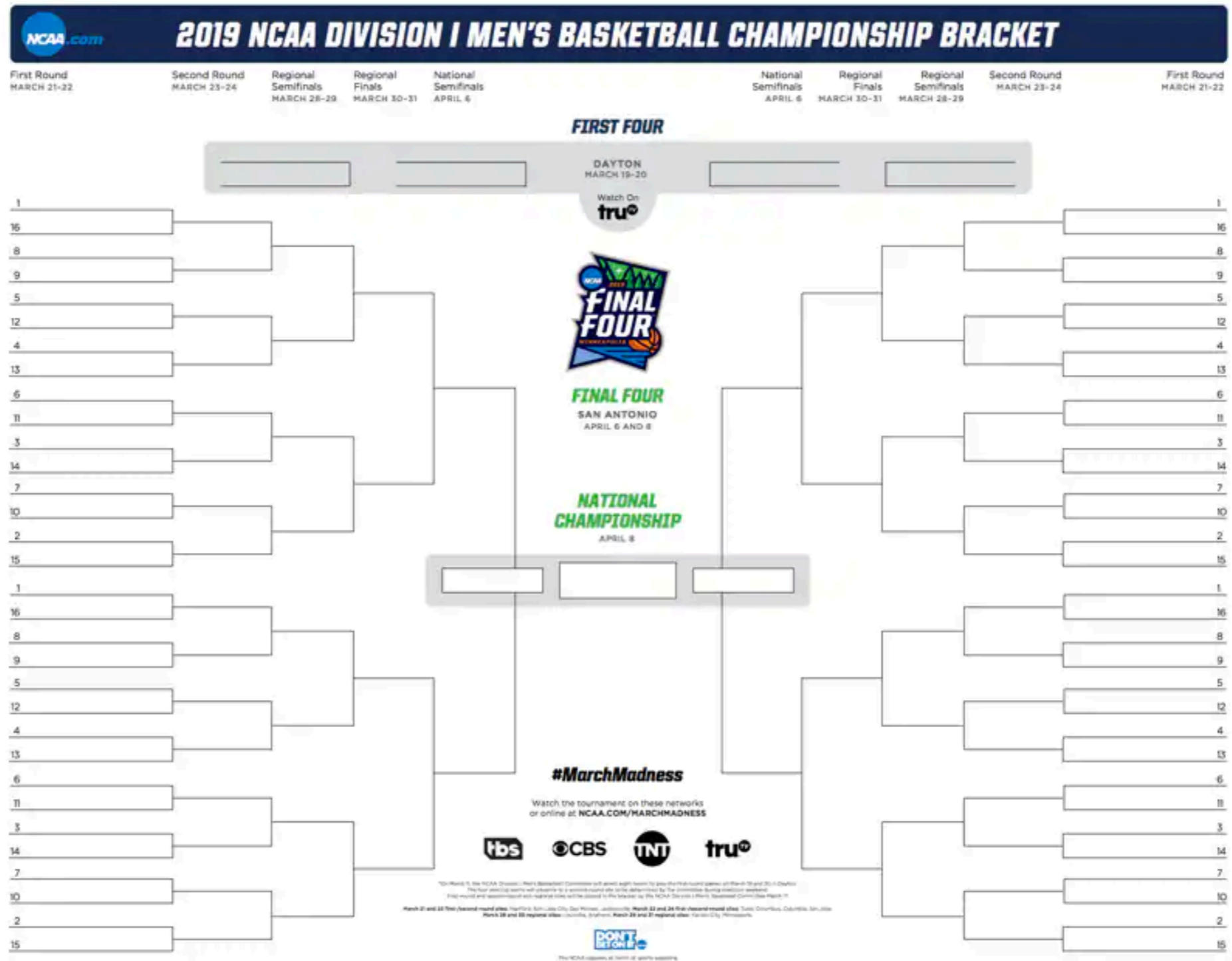
Why do we need these?

Why do we need these?

to represent **hierarchical structure**.

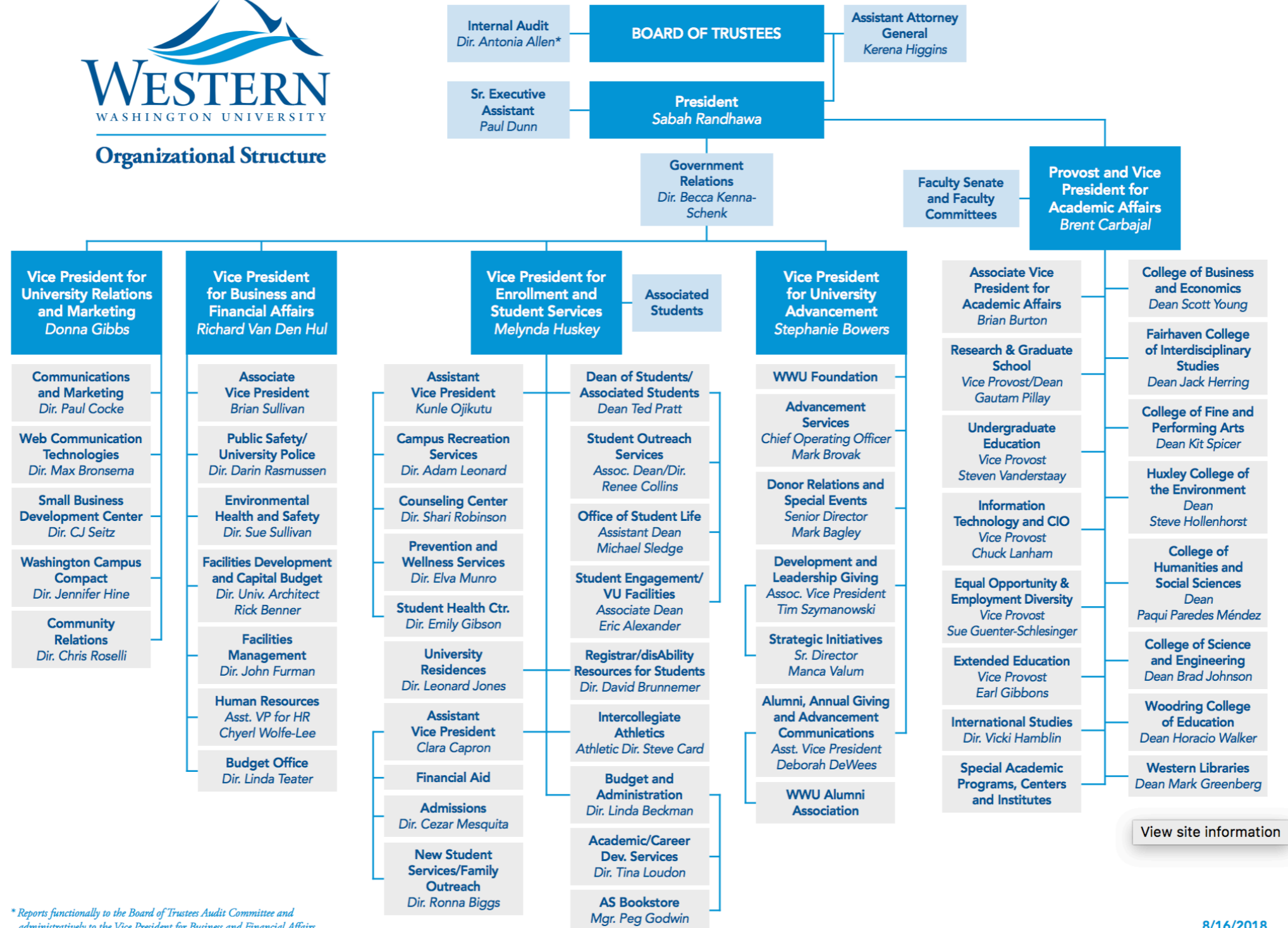
Why do we need these?

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Why do we need these?

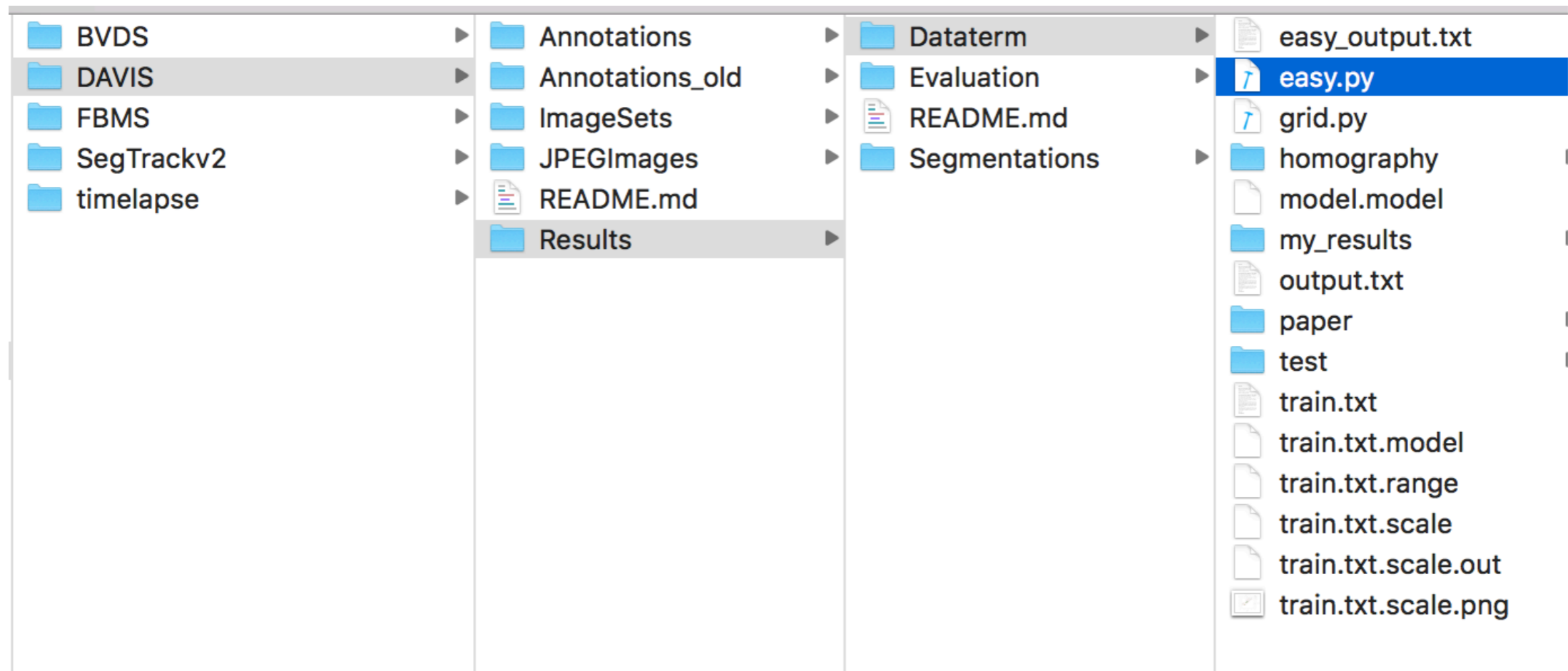
to represent hierarchical structure.



* Reports functionally to the Board of Trustees Audit Committee and administratively to the Vice President for Business and Financial Affairs

Why do we need these?

to represent **hierarchical structure**.



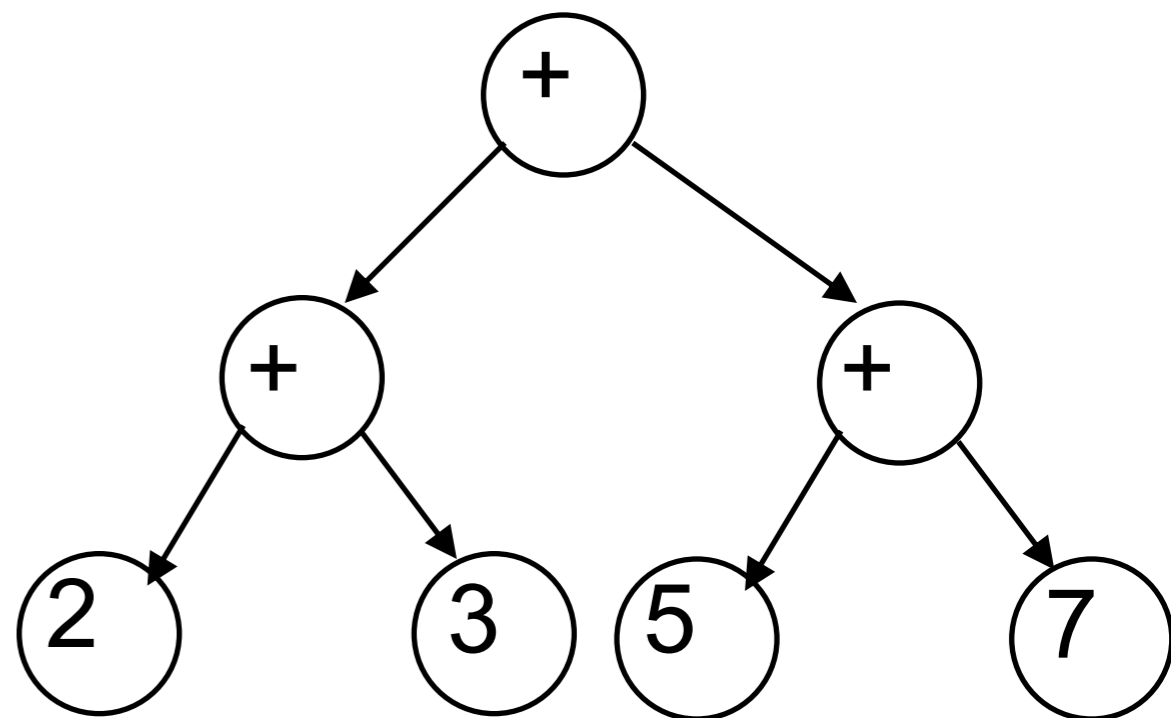
Why do we need these?

to represent **hierarchical structure**.

Syntax Trees:

- In textual representation, **parentheses** show hierarchical structure
- In tree representation, hierarchy is explicit in the tree's **structure**

$((2+3) + (5+7))$



Also used for **natural languages** and **programming languages!**

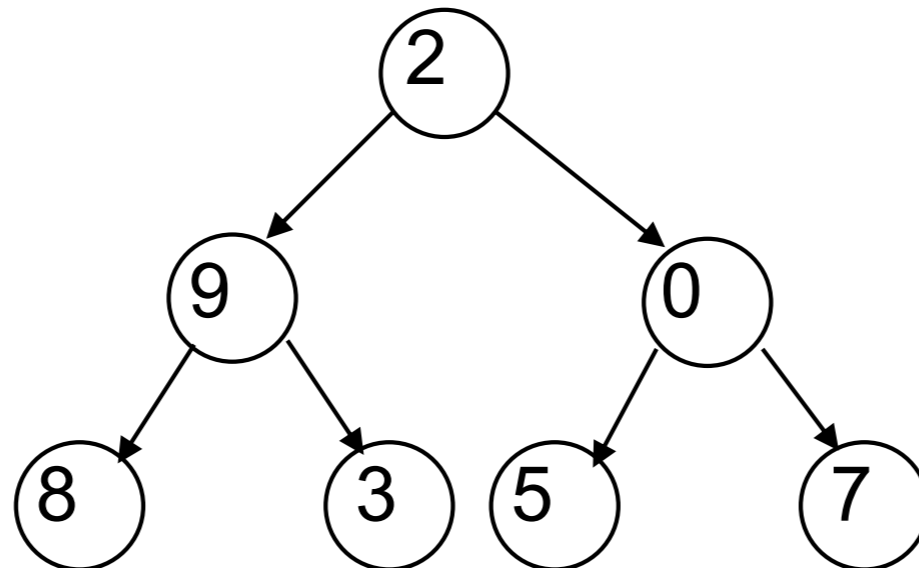
Why do we need these?

to implement various ADTs **efficiently**.

TreeSet, **TreeMap**

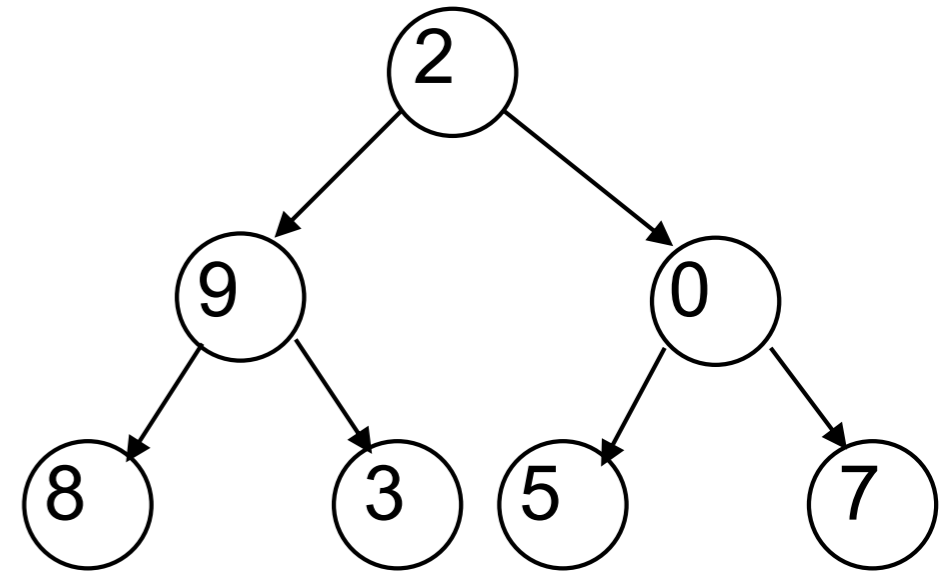
Height of a balanced binary tree is $O(\log n)$

Consequence: Many operations (find, insert, ...) can be done in **$O(\log n)$** in carefully-designed trees.



Thinking about trees recursively

- **A binary tree is**
 - Empty, or
 - Three things:
 - value
 - a left **binary tree**
 - a right **binary tree**



Thinking about trees recursively

- **A binary tree is**

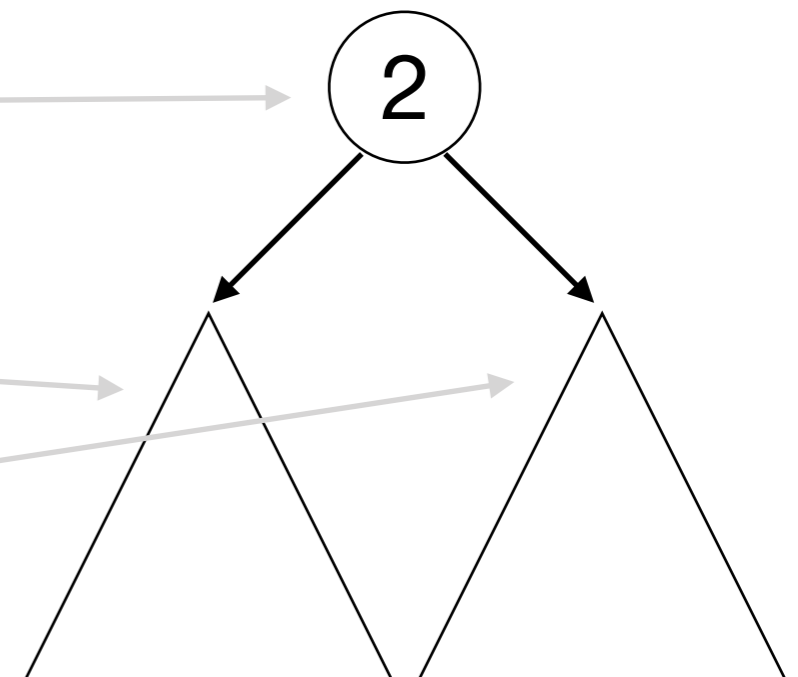
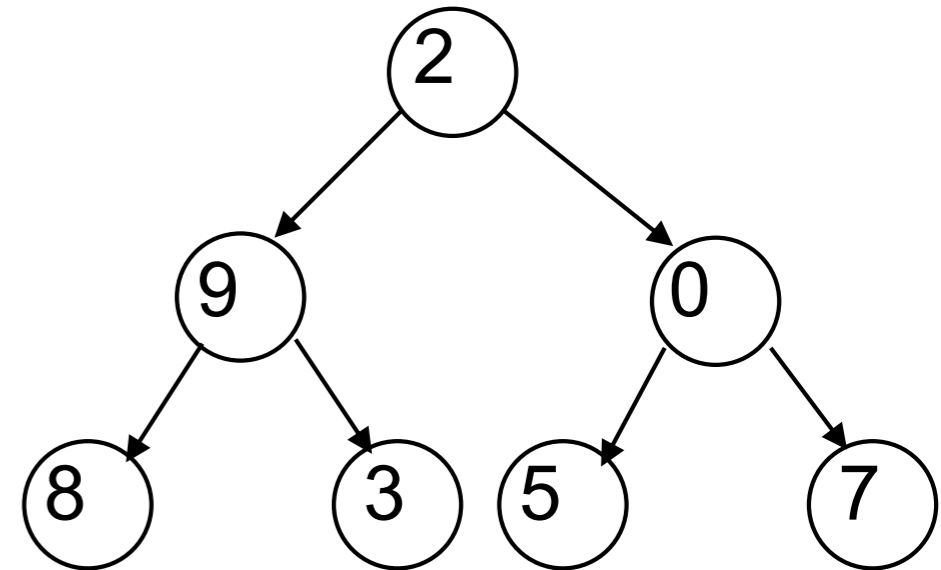
- Empty, or

- Three things:

- value

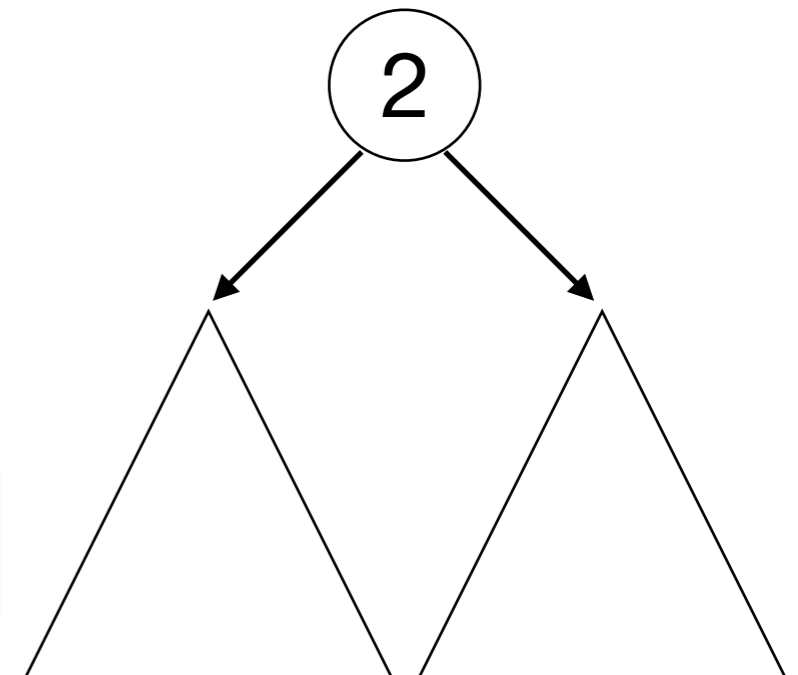
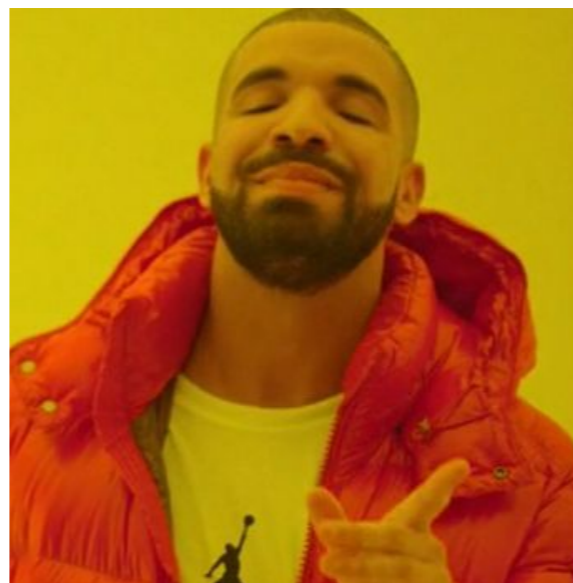
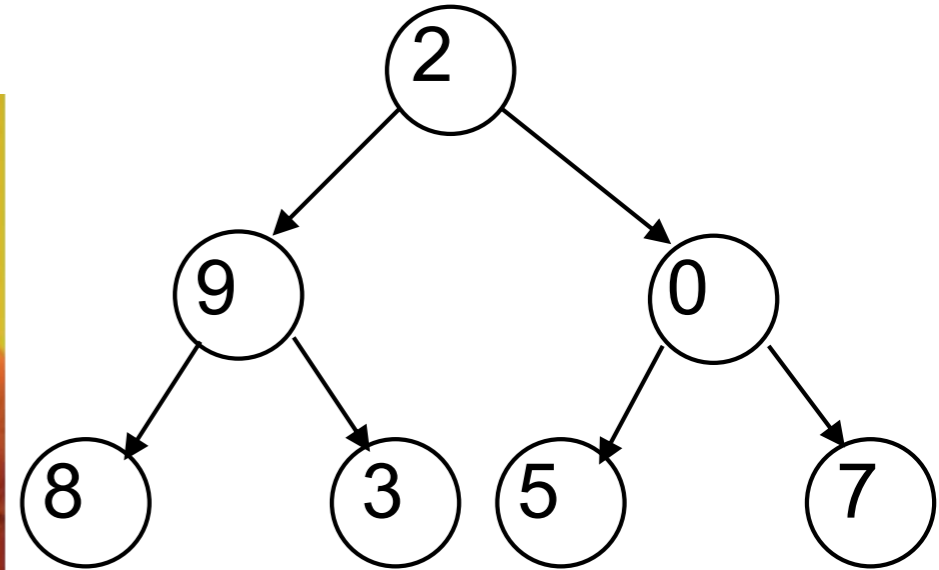
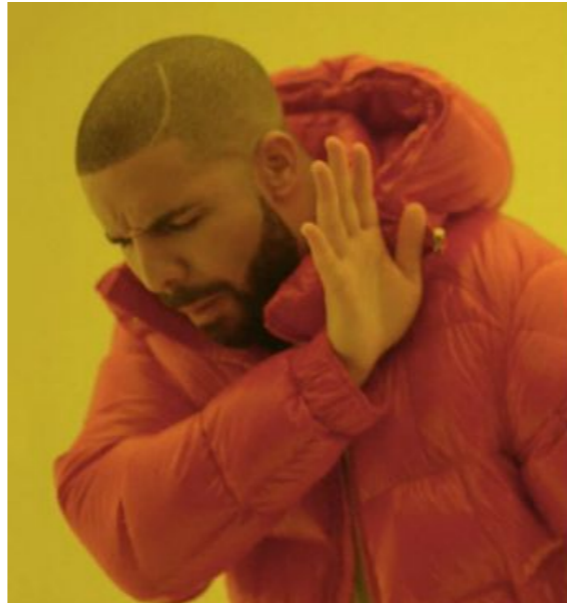
- a left **binary tree**

- a right **binary tree**



Thinking about trees recursively

- A **binary tree** is
 - Empty, or
 - Three things:
 - value
 - a left **binary tree**
 - a right **binary tree**



Operations on trees

often follow naturally from the definition of a tree:

- **A binary tree is** Find v in a binary tree:
 - Empty, or (base case - not found!)
 - Three things:
 - value (base case - is this v ?)
 - a left **binary tree** (recursive call - is v in left?)
 - a right **binary tree** (recursive call - is v in right?)

Operations on trees

often follow naturally from the definition of a tree:

- **A binary tree is**

- Empty, or

- Three things:

- value

- a left **binary tree**

- a right **binary tree**

Find v in a binary tree:

```
boolean findVal(Tree t, int v):
```

(base case - not found!)

```
if t == null:
```

```
    return false
```

(base case - is this v ?)

```
if t.value == v: return true
```

(recursive call - is v in left?)

```
return findVal(t.left)
```

```
    || findVal(t.right)
```

(recursive call - is v in right?)

Tree Traversals

Print (or otherwise process) every node in a tree:

- **A binary tree is**

- Empty, or

- Three things:

- value

- a left **binary tree**

- a right **binary tree**

Print all nodes in a binary tree:

```
boolean printTree(Tree t):
```

(base case - nothing to print)

```
if t == null:
```

```
    return
```

(print this node's value)

```
System.out.println(t.value)
```

(recursive call - print left subtree)

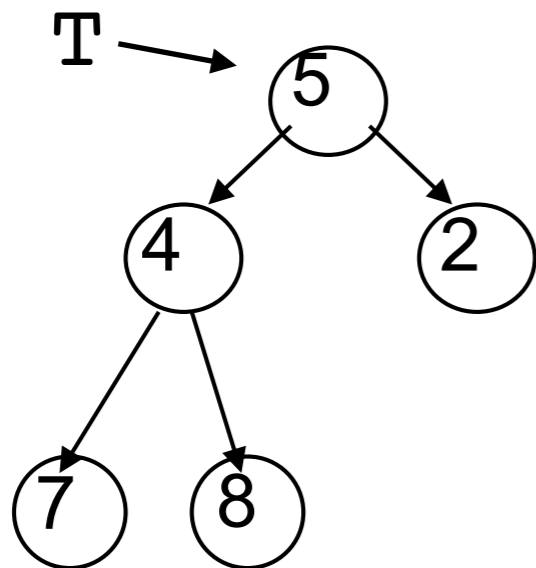
```
printTree(t.left)
```

(recursive call - print right subtree)

```
printTree(t.right)
```

Tree Traversals

Print (or otherwise process) every node in a tree:



Print all nodes in a binary tree:

```
boolean printTree(Tree t):
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(base case - nothing to print)

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if t == null:
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System.out.println(t.value)
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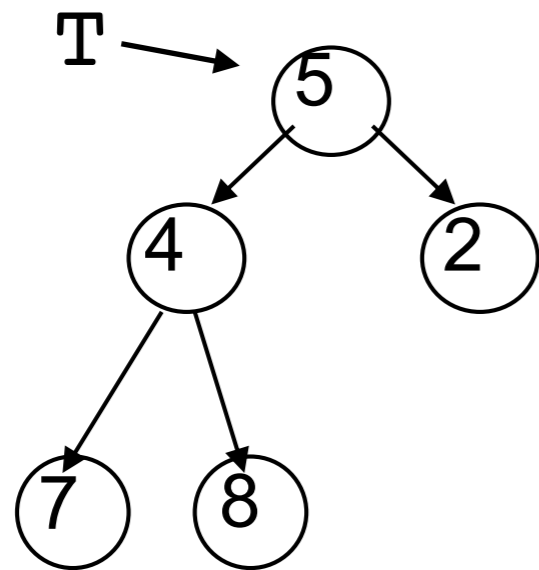
```
printTree(t.left)
```

(recursive call - print right subtree)

```
printTree(t.right)
```

Tree Traversals

Print (or otherwise process) every node in a tree:



ABCD: T is a reference to the node with value 5. What is printed by the call `printTree(T)`?

- A. 5 4 2 7 8
- B. 7 4 8 5 2
- C. 7 8 4 2 5
- D. 5 4 7 8 2

Print all nodes in a binary tree:

```
boolean printTree(Tree t):
```

(base case - nothing to print)

```
if t == null:
```

```
    return
```

(print this node's value)

```
System.out.println(t.value)
```

(recursive call - print left subtree)

```
printTree(t.left)
```

(recursive call - print right subtree)

```
printTree(t.right)
```

Tree Traversals

“Walking” over the whole tree is called a **tree traversal**. This is done often enough that there are standard names. Previous example was a **pre-order traversal**:

1. **Process root**
2. Process left subtree
3. Process right subtree

Other common traversals:

in-order traversal:

1. Process left subtree
2. **Process root**
3. Process right subtree

post-order traversal:

1. Process left subtree
2. Process right subtree
3. **Process root**

Why do we need these?

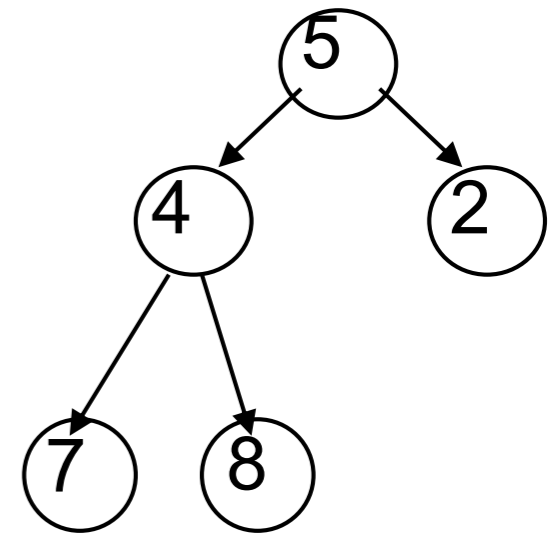
to represent **hierarchical structure**.

Quadrees in graphics and simulation:

<https://www.youtube.com/watch?v=fuexOsLOf10>

Practice Exercise

- Write the values printed by a:
 - pre-order
 - in-order
 - post-order



traversal of this tree.

Terminology - Self-Quiz

root

subtree

leaf

child

parent

ancestor

descendant

depth

height

