



CSCI 241

Lecture 11

Balanced Binary Search Trees

Announcements

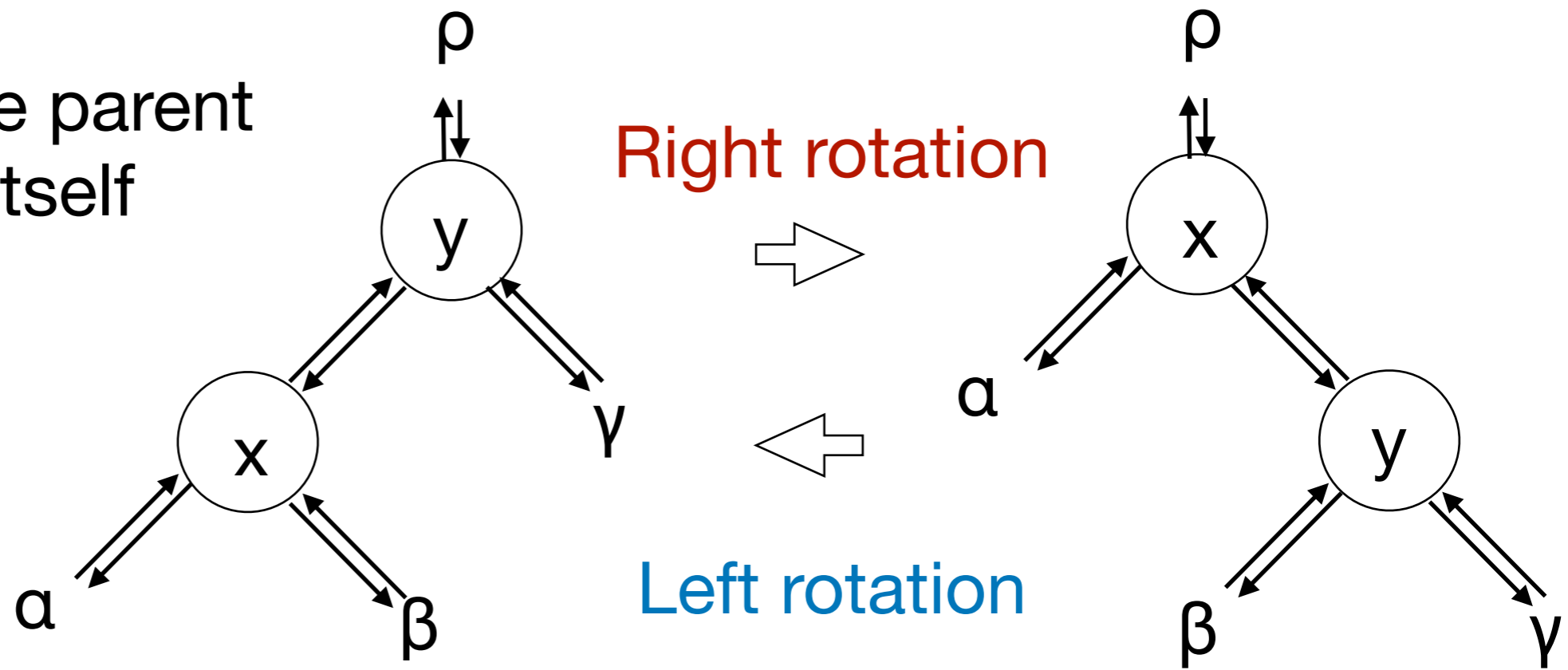
Goals

- Be prepared to implement **rotations** in BSTs
- Be prepared to implement AVL rebalancing.

Tree Rotations

Steps in left rotation (move y up to x 's position):

1. Transfer β
2. Transfer the parent
3. Transfer x itself



$x.R$ gets $y.L$

$y.L.p$ gets x

$y.p$ gets $x.p$

$p.[L/R]$ gets y

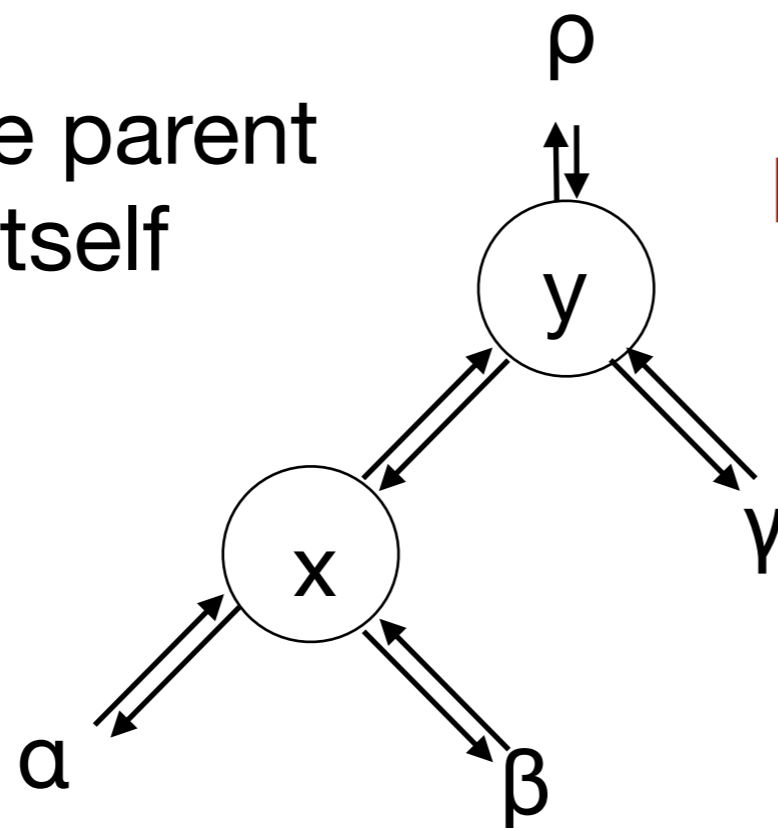
$y.L$ gets x

$x.p$ gets y

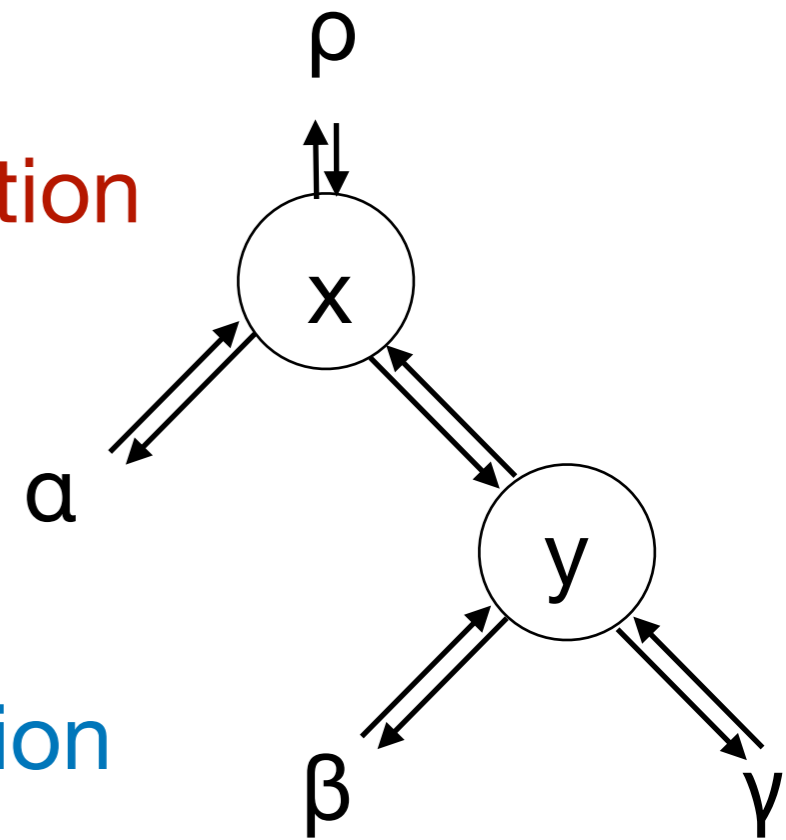
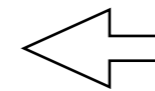
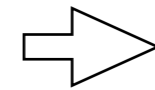
Warm-up

Steps in left rotation (move y up to x 's position):

1. Transfer β
2. Transfer the parent
3. Transfer x itself



Right rotation



Left rotation

$x.R$ gets $y.L$

$y.L.p$ gets x

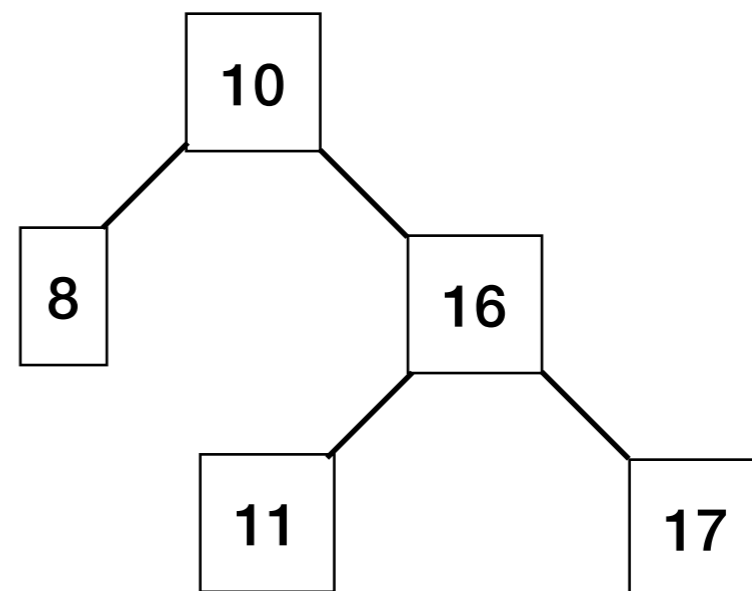
$y.p$ gets $x.p$

$p.[L/R]$ gets y

$y.L$ gets x

$x.p$ gets y

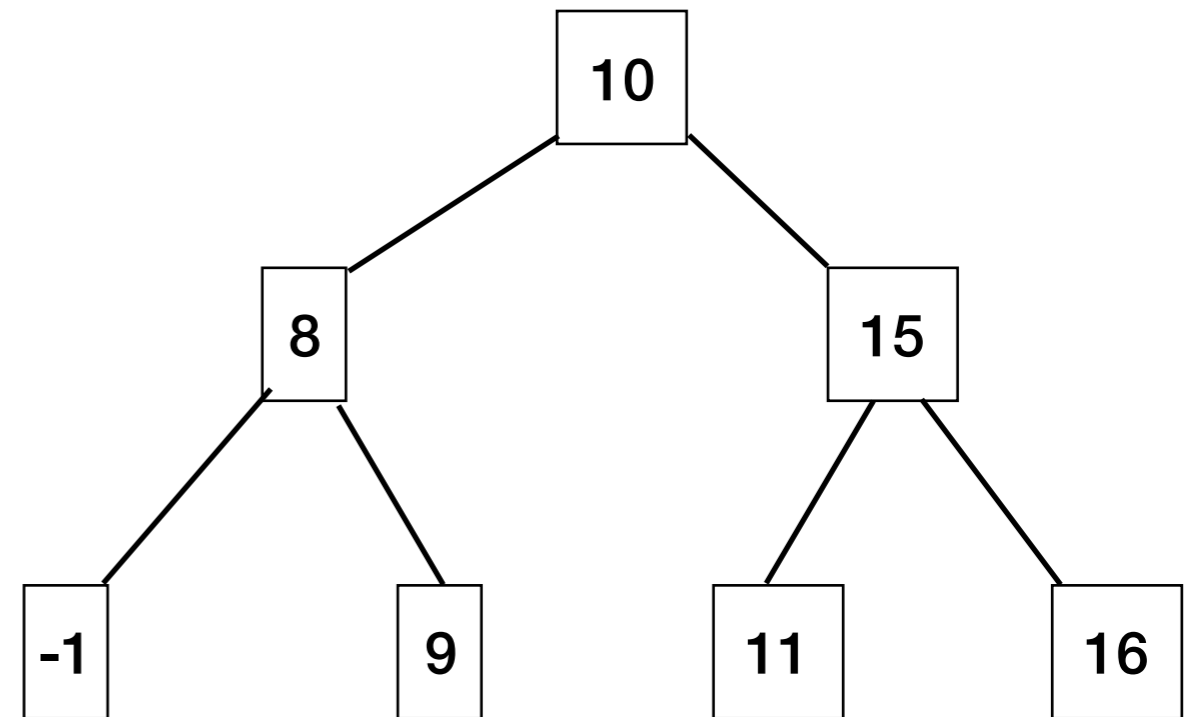
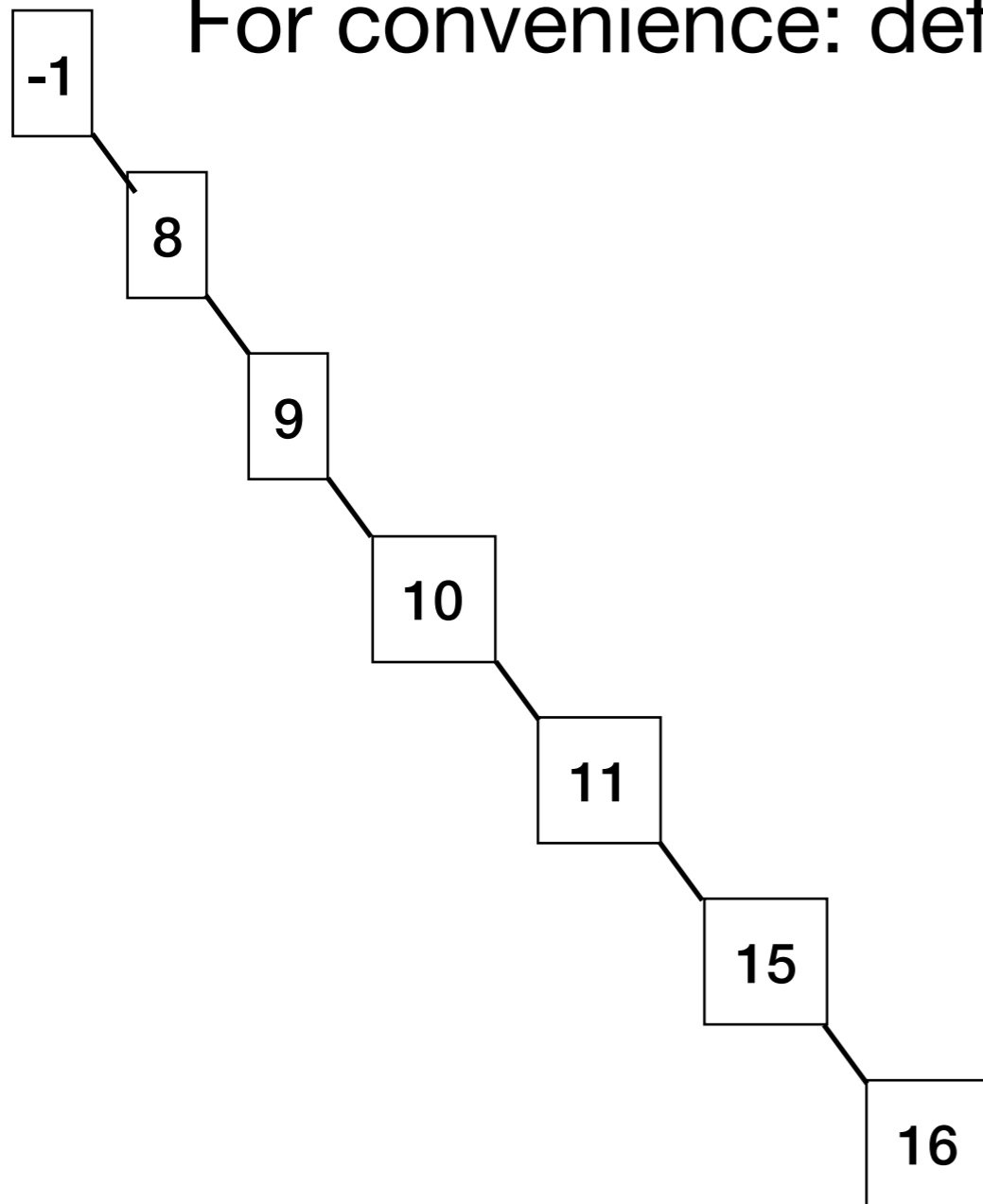
Perform a **left** rotation
on the root of this tree:



Can we improve balance?

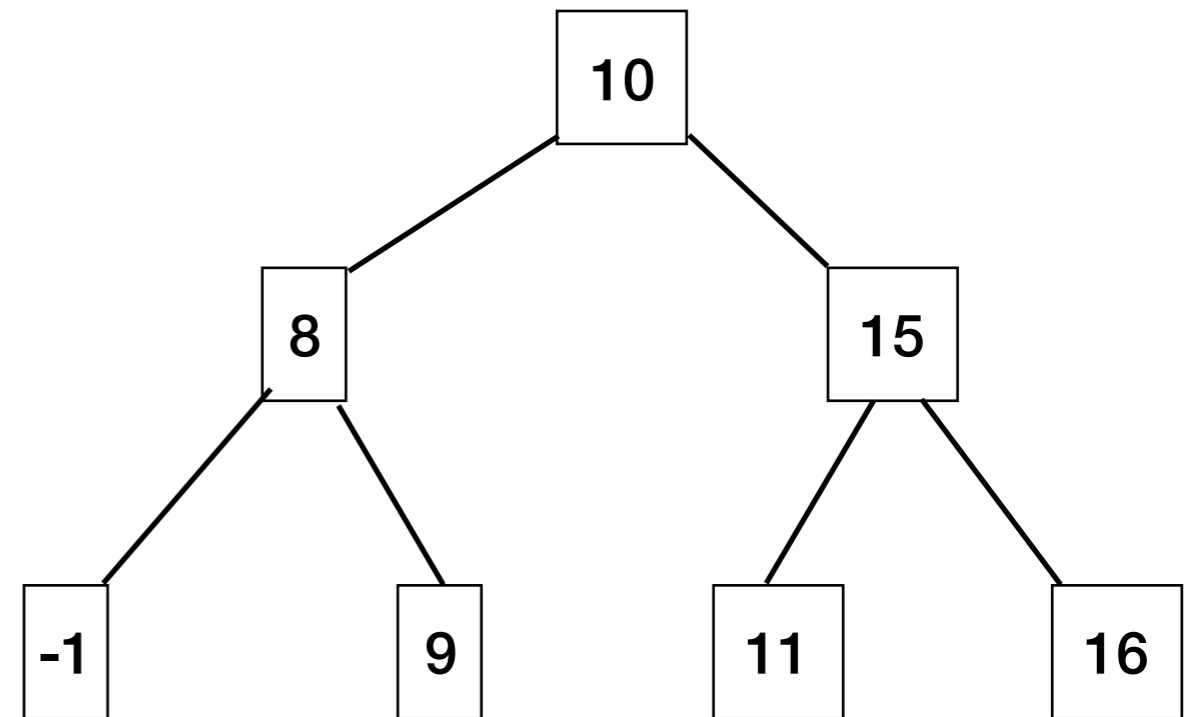
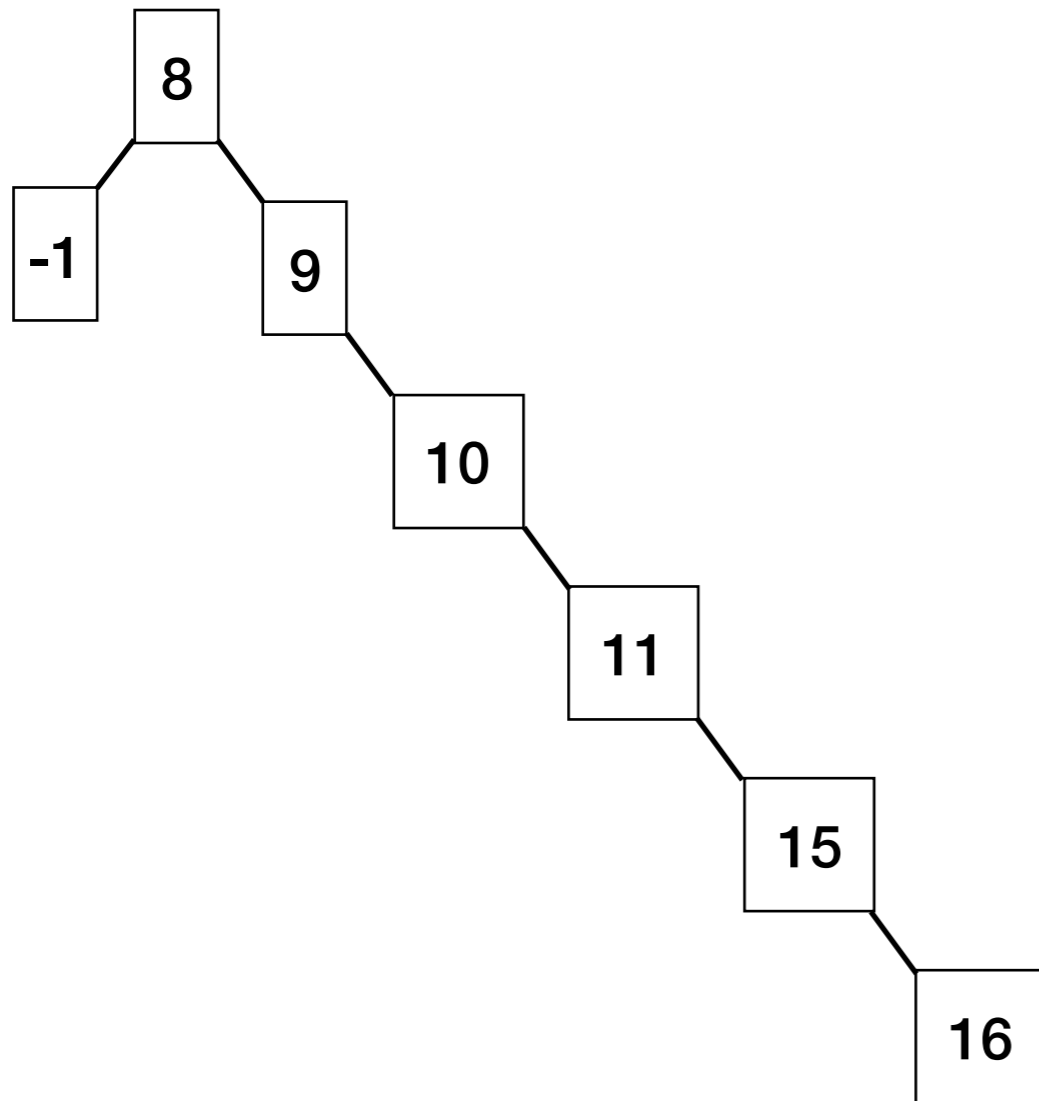
Balance Factor(n) = height(n.right) - height(left)

For convenience: define **height(null)** = -1



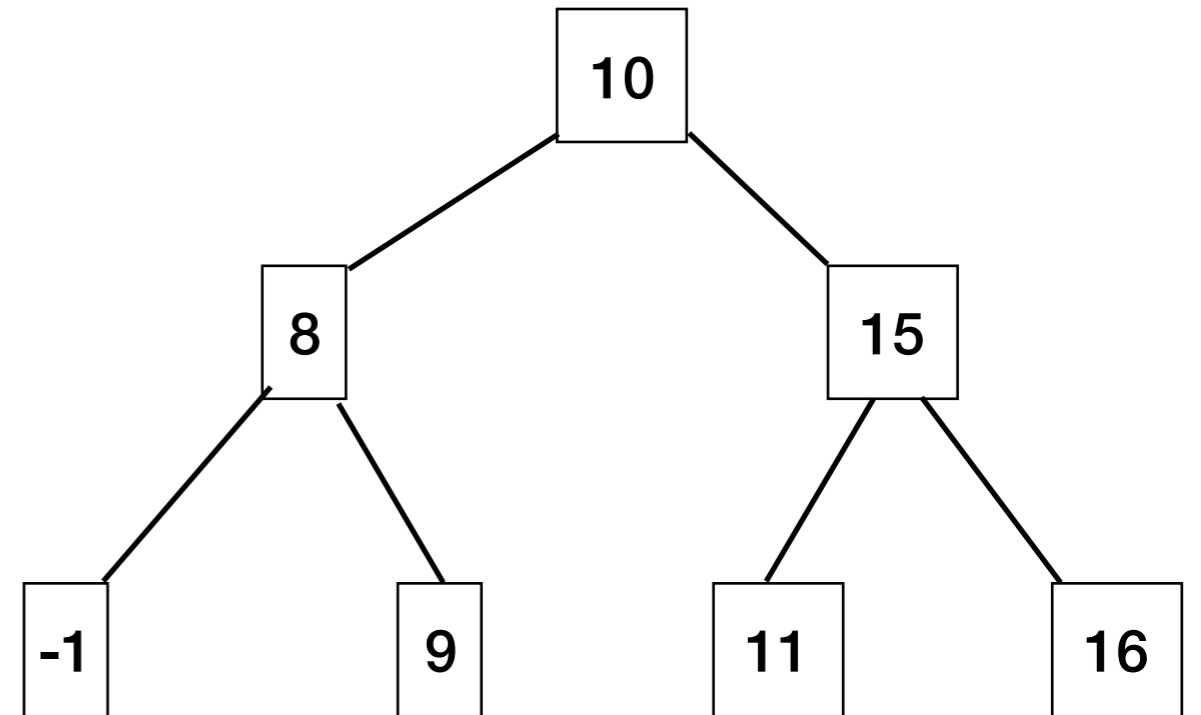
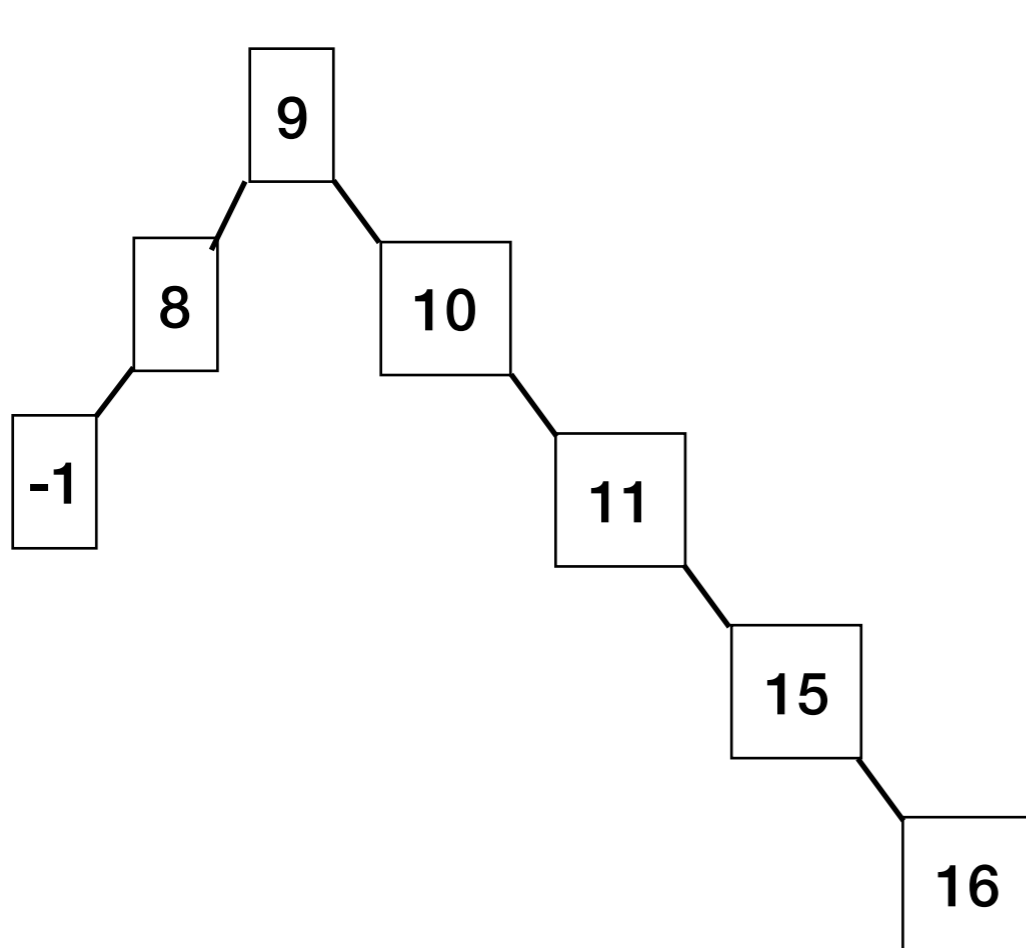
Can we improve balance?

Balance Factor(n) = height(n.right) - height(left)



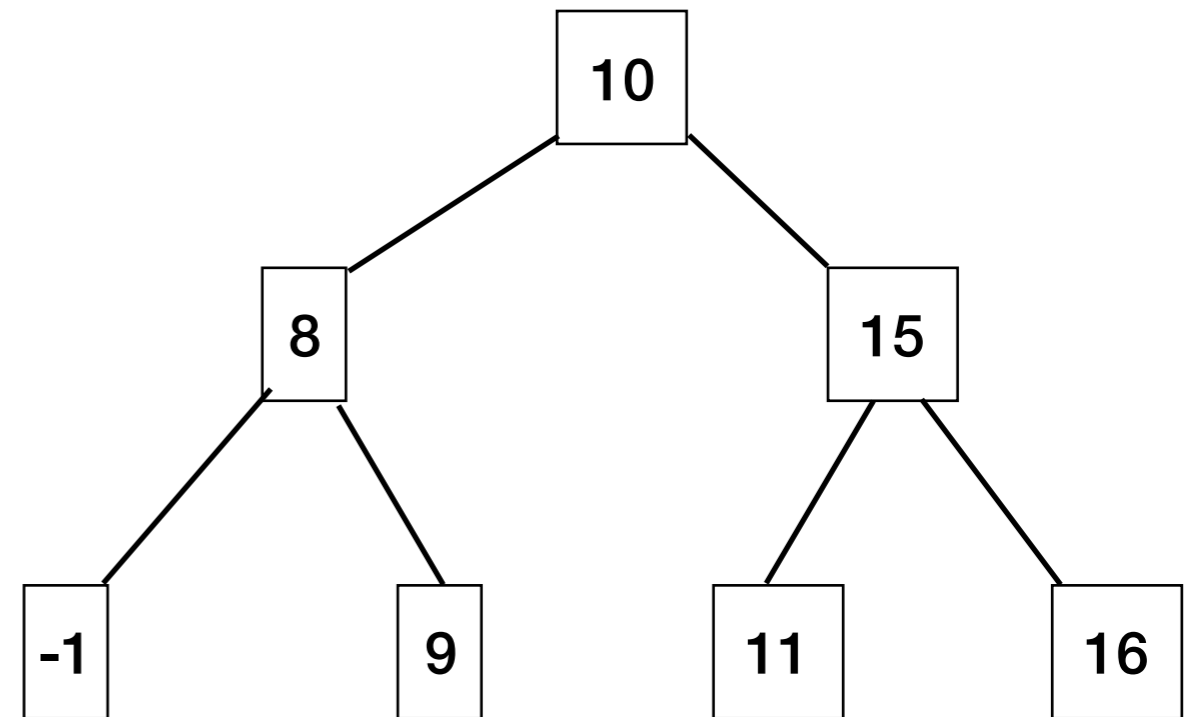
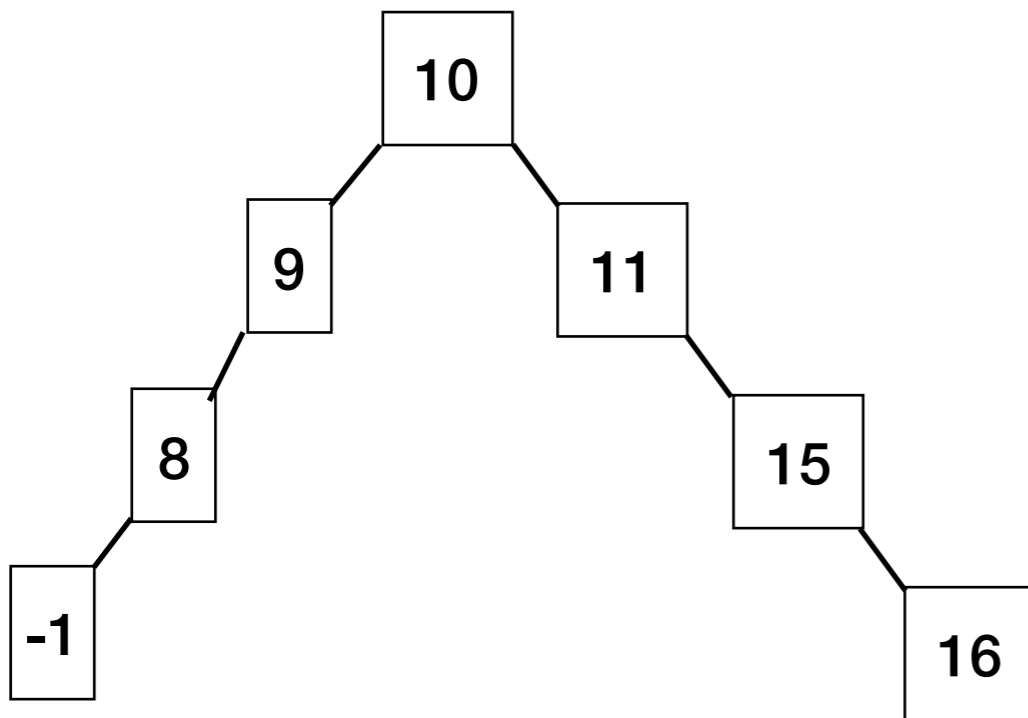
Can we improve balance?

Balance Factor(n) = height(n.right) - height(left)



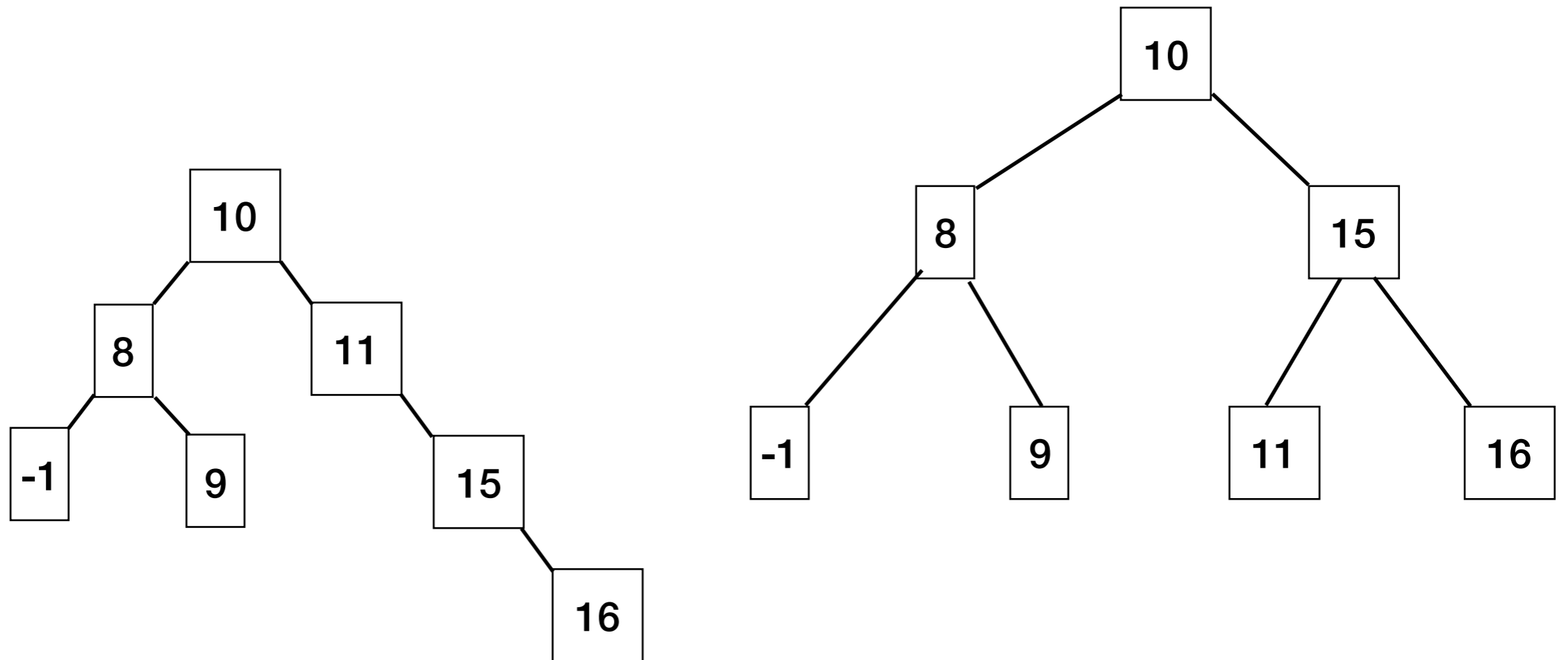
Can we improve balance?

Balance Factor(n) = height(n.right) - height(left)



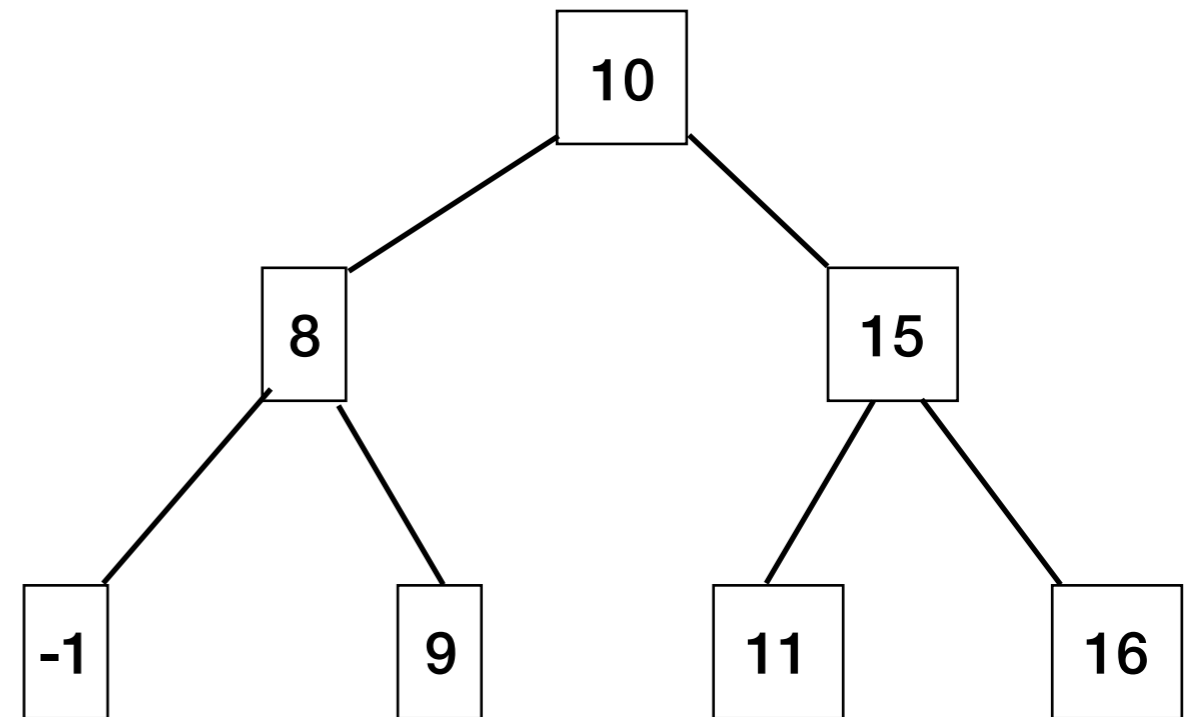
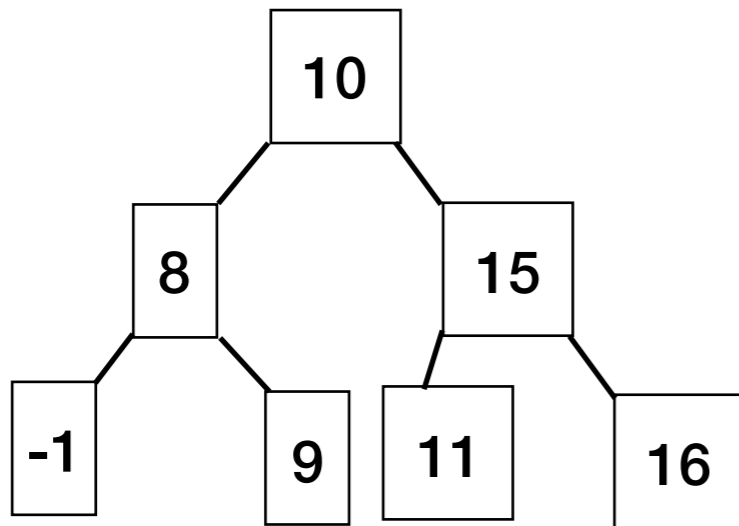
Can we improve balance?

Balance Factor(n) = height(n.right) - height(left)



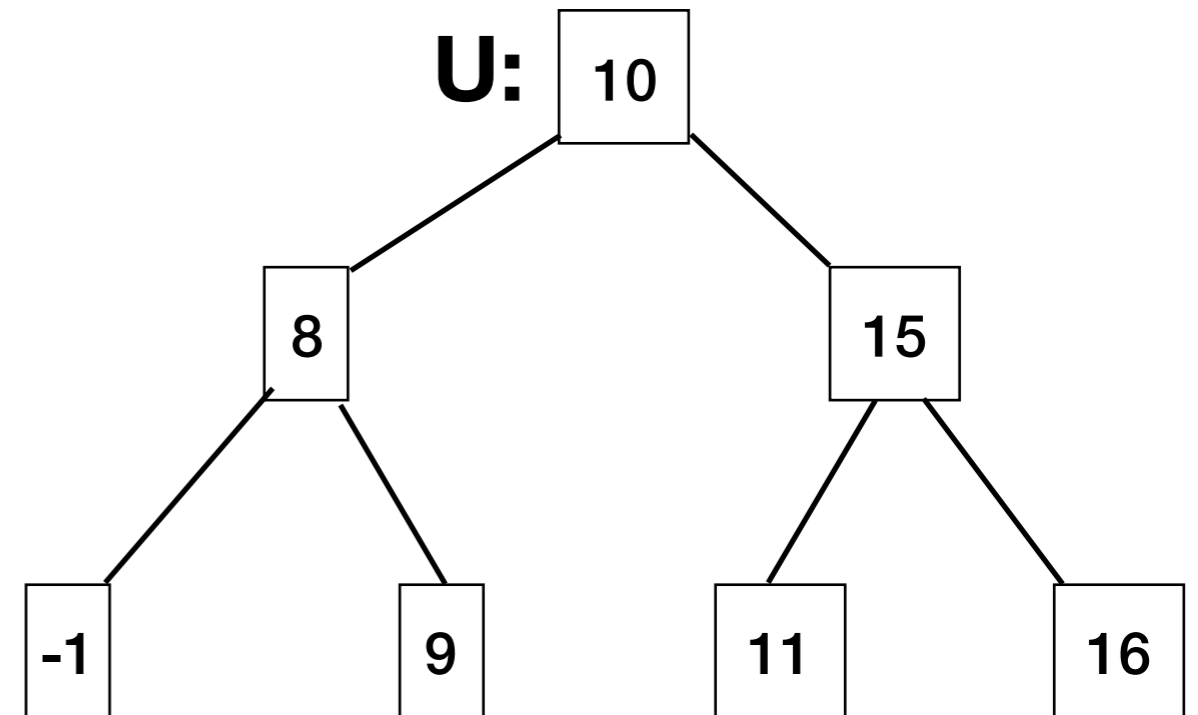
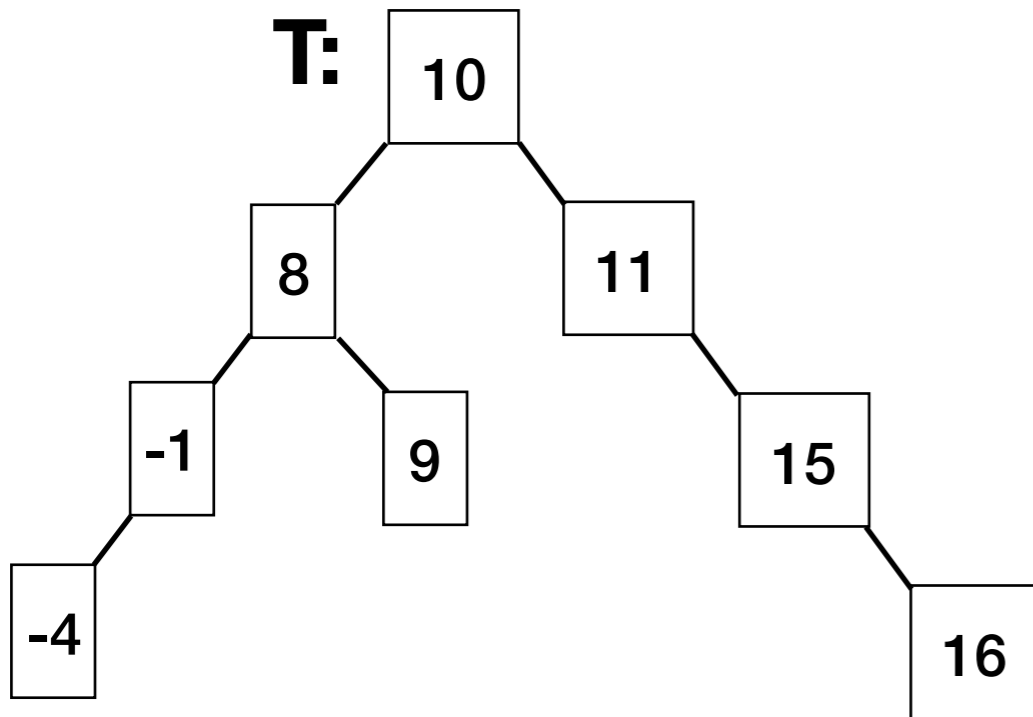
Can we improve balance?

Balance Factor(n) = height(n.right) - height(left)



Balance Factor

Balance Factor $b(n)$ = height(n.right) - height(left)



ABCD: What's the largest *absolute* balance factor of any node in each tree?

	T	U
A	0	0
B	2	1
C	2	0
D	1	1

AVL Trees

Balance Factor $b(n)$ = height(n.right) - height(left)

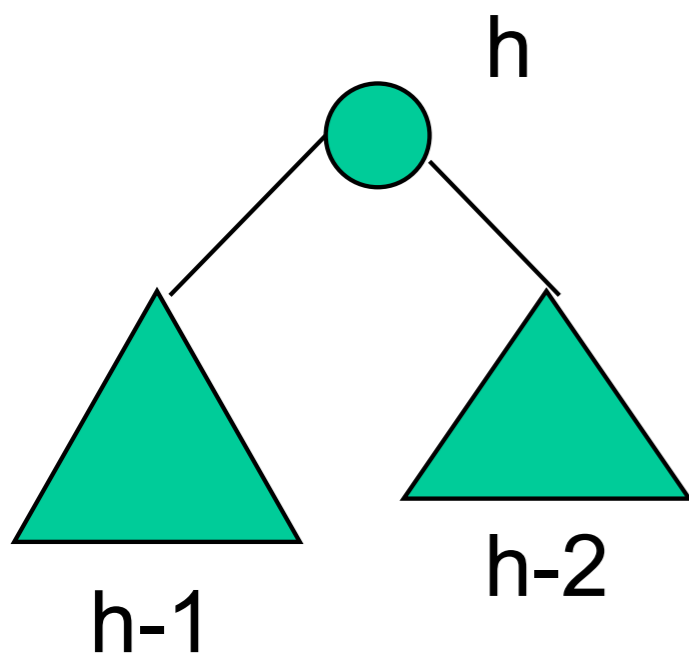
- Devised by **Adelson-Velsky** and **Landis**
- An AVL tree is a Binary Search Tree in which the following property holds:

AVL property: $-1 \leq b(n) \leq 1$ for all nodes n .

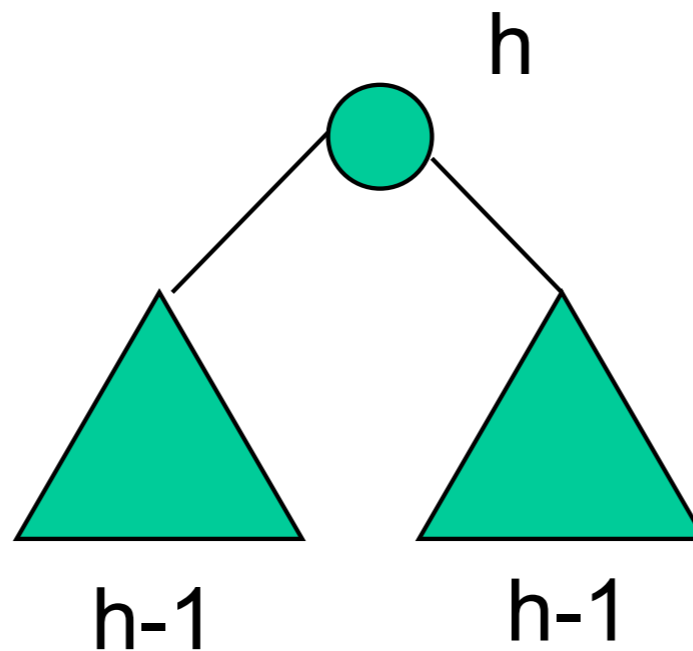
Balance Factor in AVL Trees

AVL property: $-1 \leq b(n) \leq 1$ for all nodes n .

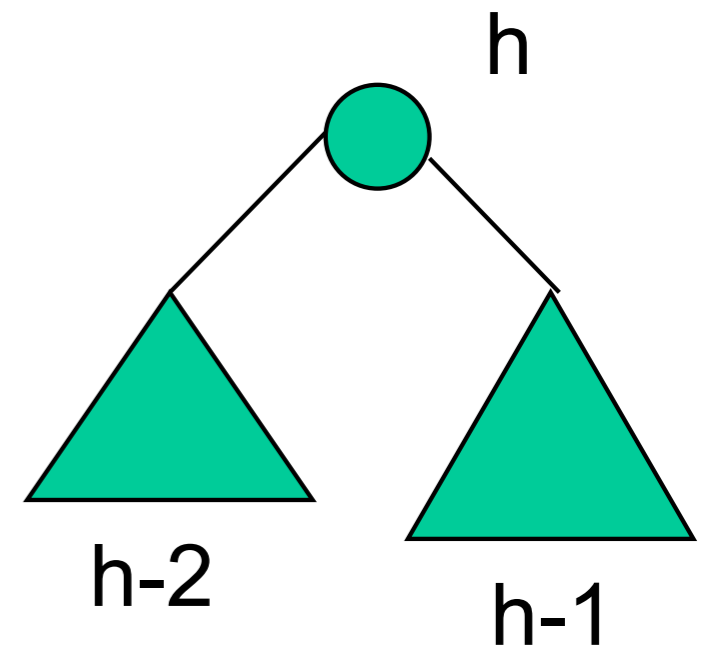
Every subtree in an AVL tree looks like one of these three trees:



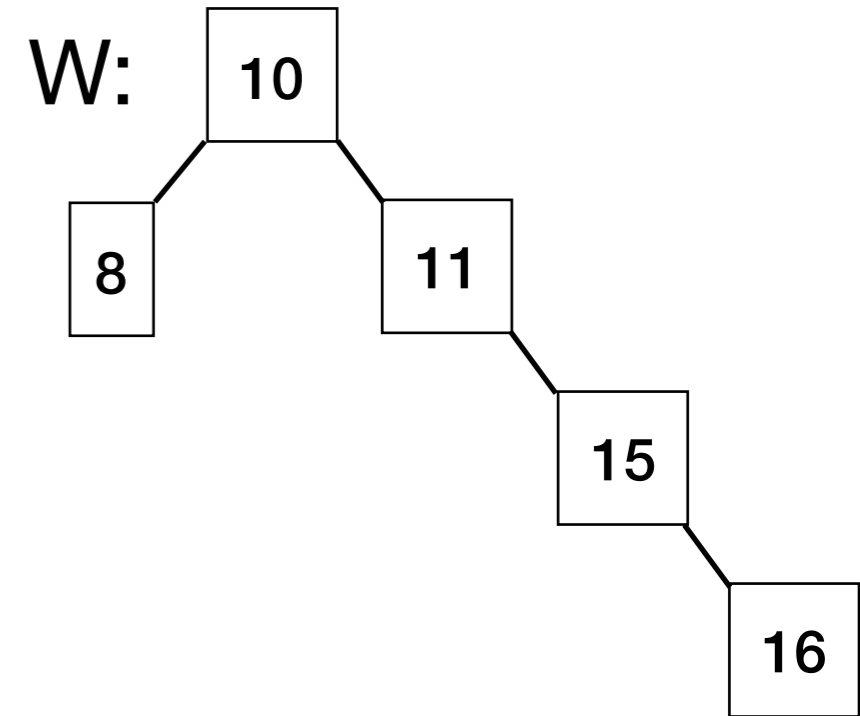
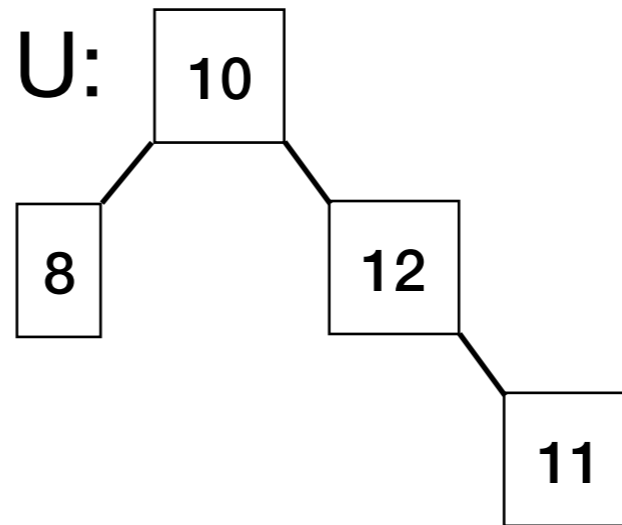
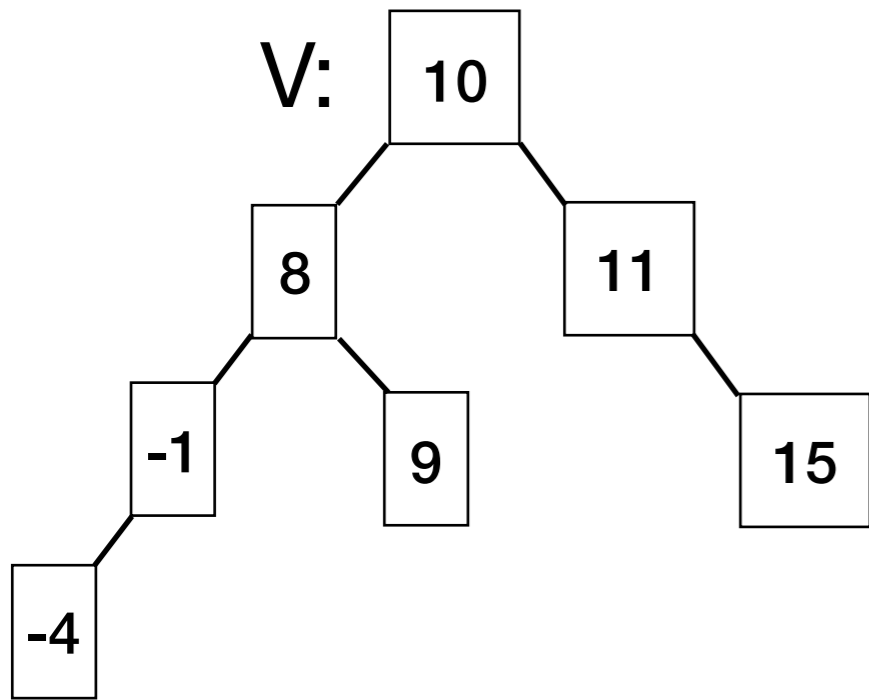
(a) Balance factor: 1



(b) Balance factor: 0

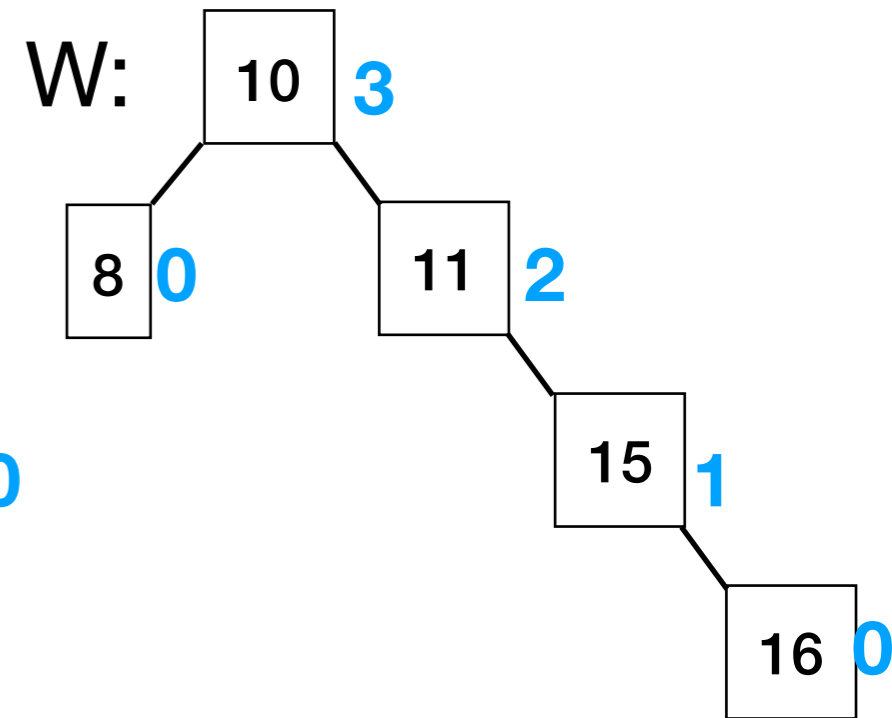
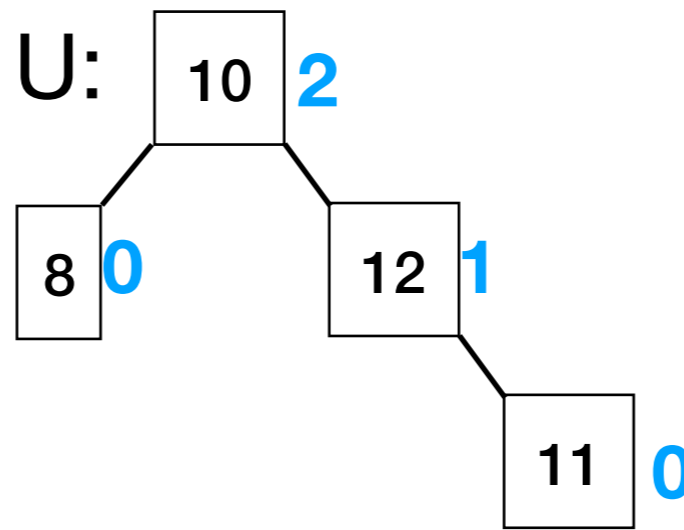
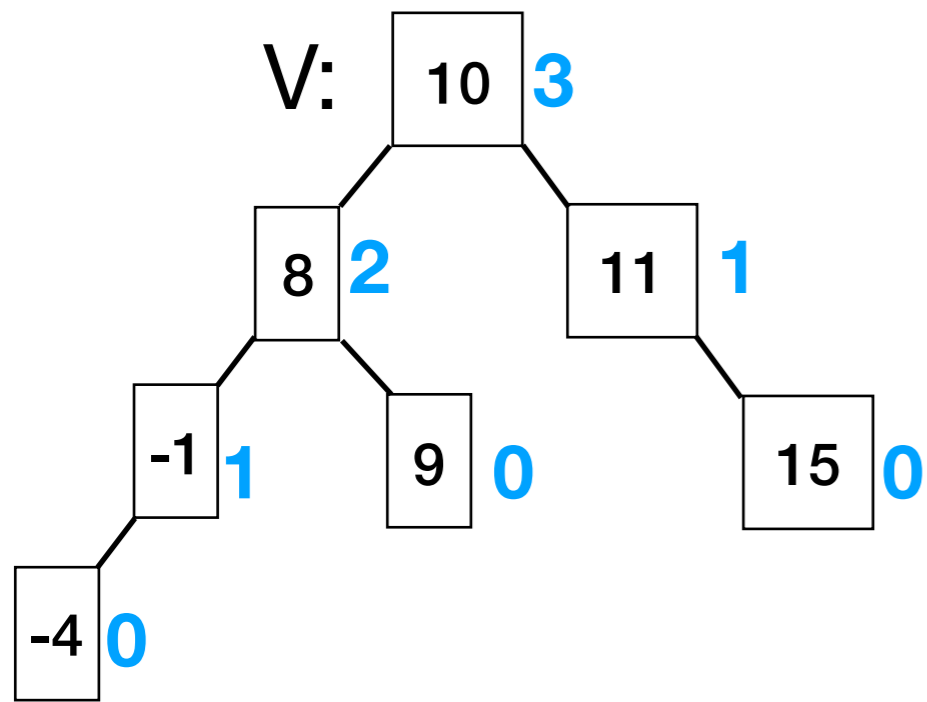


(c) Balance factor: -1



ABCD: Which of these is/are **not** AVL trees?

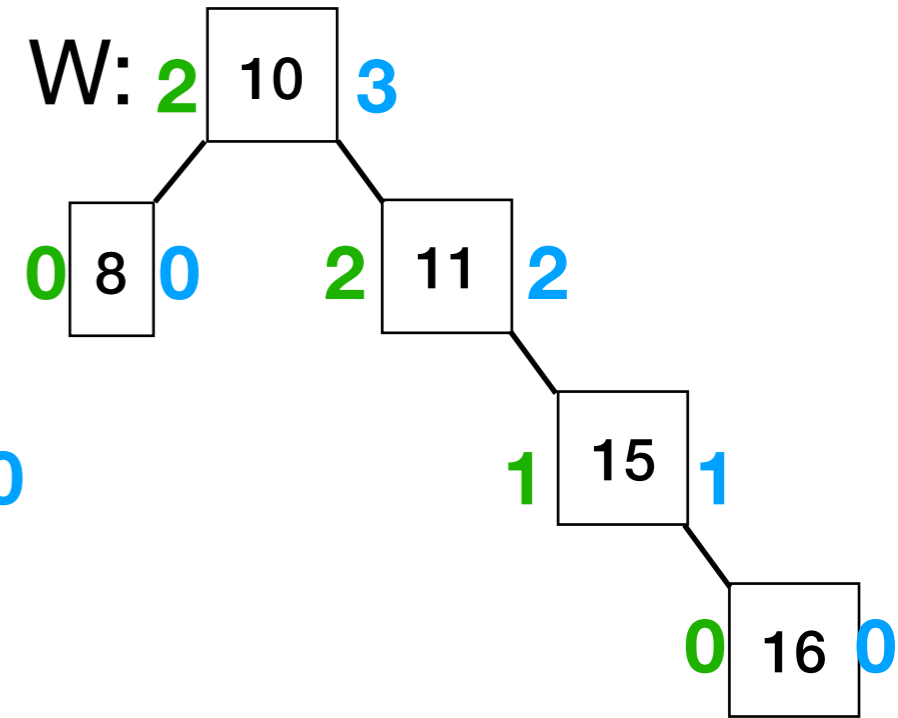
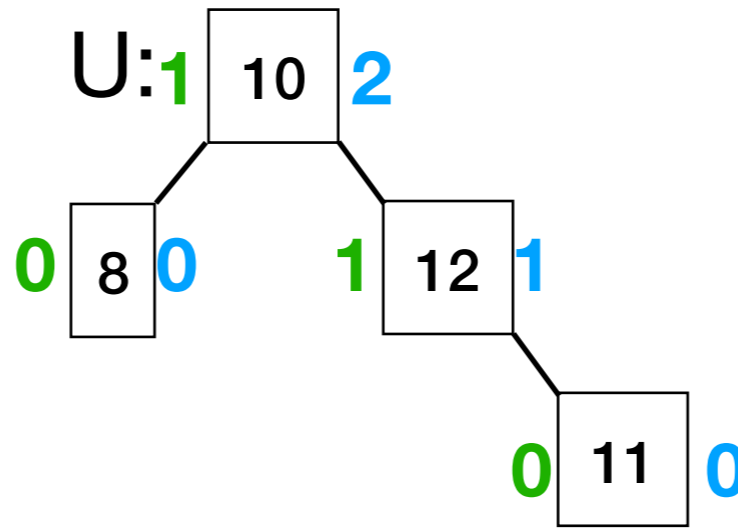
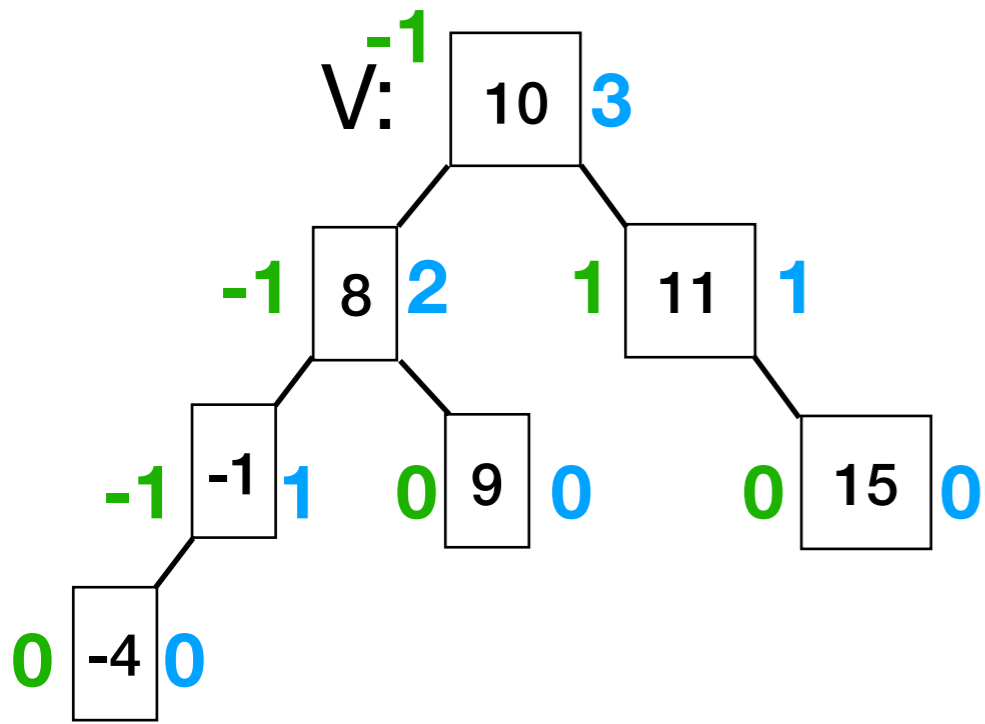
- A. U
- B. W
- C. V and W
- D. U and W



ABCD: Which of these is/are **not** AVL trees?

- A. U
- B. W
- C. V and W
- D. U and W

Heights in blue.



ABCD: Which of these is/are **not** AVL trees?

- A. U
- B. W
- C. V and W
- D. U and W

Heights in blue.

Balance factors in green.

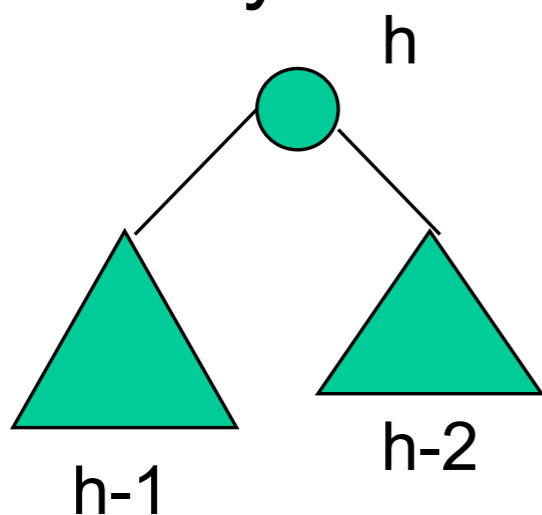
AVL Trees: Insertion

- An AVL tree is a Binary Search Tree in which the following property holds:

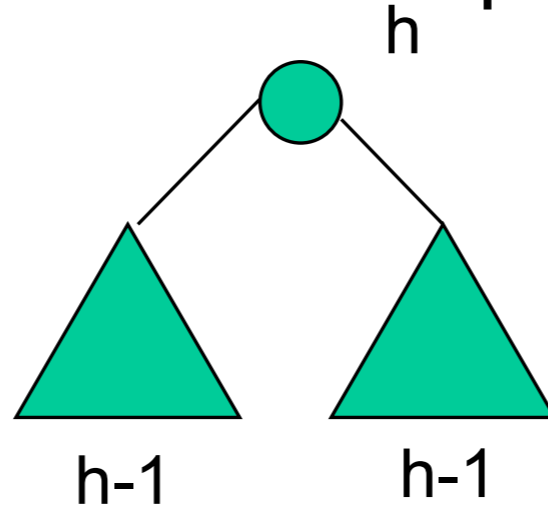
AVL property: $-1 \leq b(n) \leq 1$ for all nodes n .

To insert into an AVL tree:

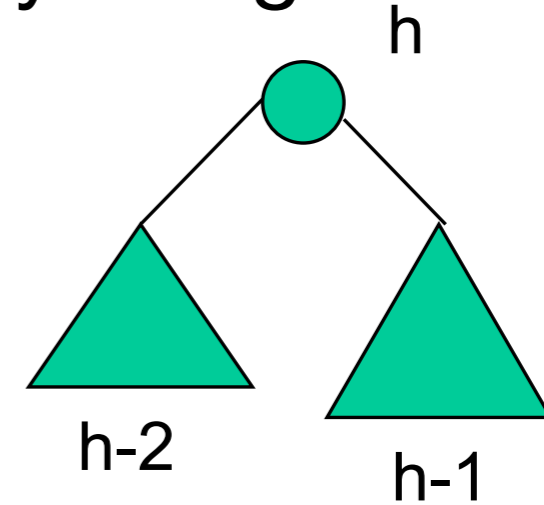
1. Do a normal BST insertion
2. Fix any violations of the AVL property using rotations.



(a) Balance factor: 1



(b) Balance factor: 0



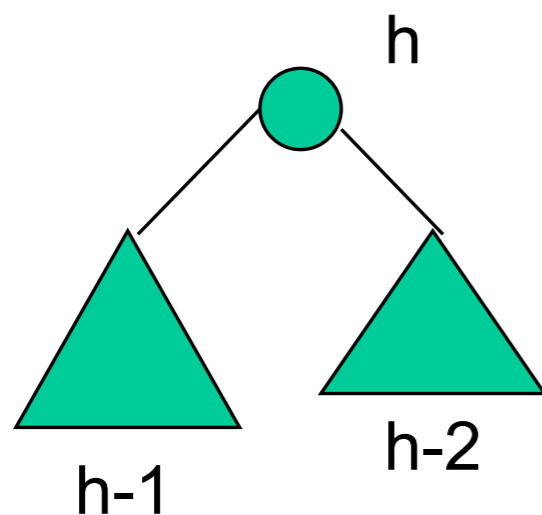
(c) Balance factor: -1

AVL Trees: Insertion

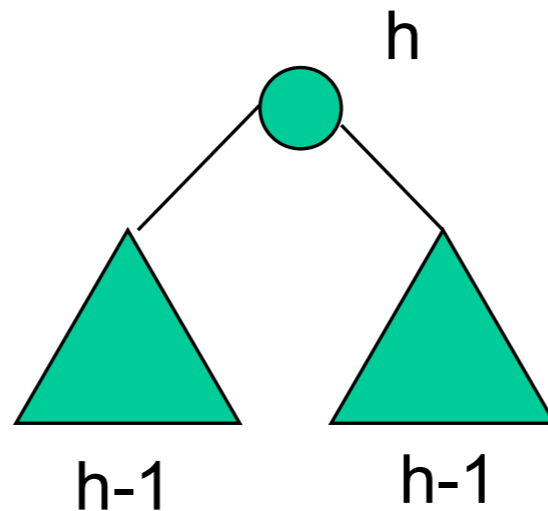
AVL property: $-1 \leq b(n) \leq 1$ for all nodes n .

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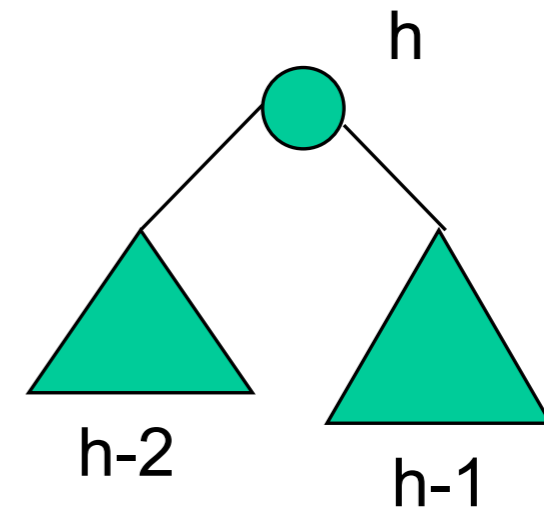
1. Do a normal BST insertion
2. Fix any violations of the AVL property using rotations.



(a) Balance factor: 1



(b) Balance factor: 0



(c) Balance factor: -1

Refresher: BST Insertion

```
/* insert a node with value v into the
 * tree rooted at n. pre: n is not null. */
insert(Node n, int v):
    if n.value == v: return // (duplicate)
    if v < n.value:
        if n has left:
            insert(n.left, v)
        else:
            // attach new node w/ value v to n.left
    else: // v > n.value
        if n has right:
            insert(n.right, v)
        else:
            // attach new node w/ value v to n.right
```

AVL Insertion

```
/* insert a node with value v into the
 * tree rooted at n. pre: n is not null. */
insert(Node n, int v):
    if n.value == v: return // (duplicate)
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        if n has left:
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            // attach new node w/ value v to n.left
    else: // v > n.value
        if n has right:
            insert(n.right, v)
        else:
            // attach new node w/ value v to n.right
    rebalance(n); ←
```