

# CSCI 241

Lecture 8  
Introduction to Trees

# Analyze ALL the Sorts!

	InsertionSort	SelectionSort	MergeSort	QuickSort	RadixSort
Best-case					
Average-case	same as worst	same as worst	same as worst		same as worst
Worst-case					
Stable?					
In-place?					

# Announcements

- Nick's office hours: Tuesdays 1-3pm, CF 167
- Events next week:
  - Monday, Oct. 15 – [Tech Talk: Alaska Airlines](#) – 5 pm in CF 110
  - Tuesday, Oct. 16 – [Peer Lecture Series: CS Success Workshop](#) – 4 pm in CF 420
  - Wednesday, Oct. 17 – [Tech Talk: Integra Group](#) – 5 pm in CF 125
  - Saturday & Sunday, Oct. 20 & 21 – [Fall Game Jam!](#) – 10 am in CF 105
- Regular club meetings:
  - [AI Club](#) - Tuesdays 6pm in PH 228 (talk to Sakari!)
  - [Game Design Club](#) - Mondays 6pm in CF 105 (talk to Kale!)
  - Others - see <https://cse.wvu.edu/cs/cs-clubs>

# Goals:

- Understand the definition of a tree.
- Know the basic terminology associated with trees:
  - Root, child, parent, leaf, height, depth, subtree, descendent, ancestor
- Be able to write a tree class and some simple recursive processing methods.

# Linked List

```
public class ListNode {  
    int value;  
    ListNode next;  
}
```

# Linked List

```
public class List {  
    int value;  
    List next;  
}
```

The node *is the list*.

Next points to the **tail** of the list (also a list!)

# Binary Tree

```
public class Tree {  
    int value;  
    Tree left;  
    Tree right;  
}
```

The node *is the tree*.

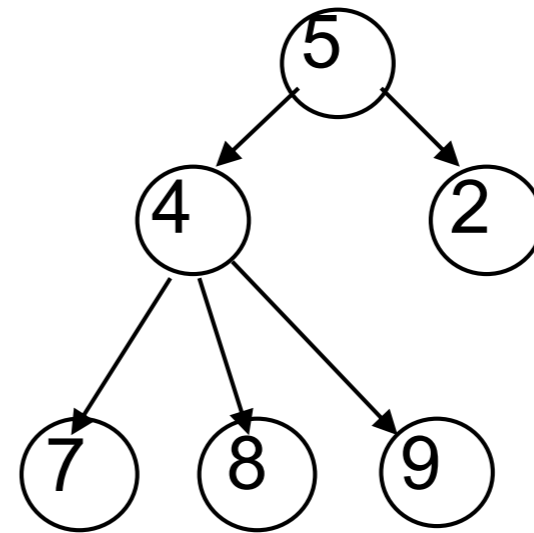
# Tree - Definition

**Tree:** like a linked list, but:

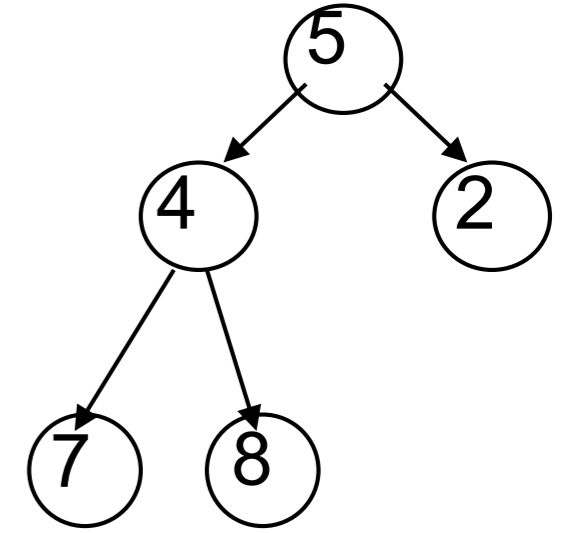
- Each node may have zero or more *successors* (**children**)
- Each node has exactly one *predecessor* (parent) except the *root*, which has none
- All nodes are reachable from *root*

**Binary tree:** A tree, but:

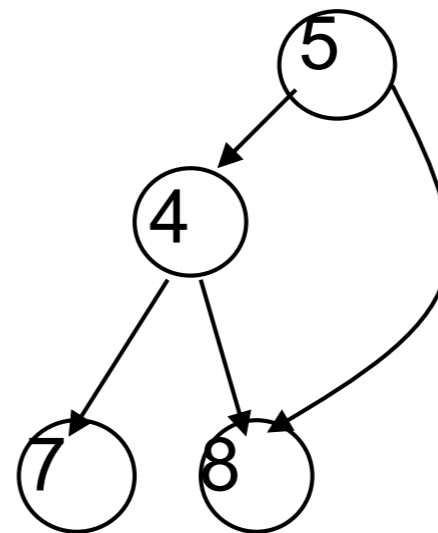
- Each can have at most **two** children (left child, right child)



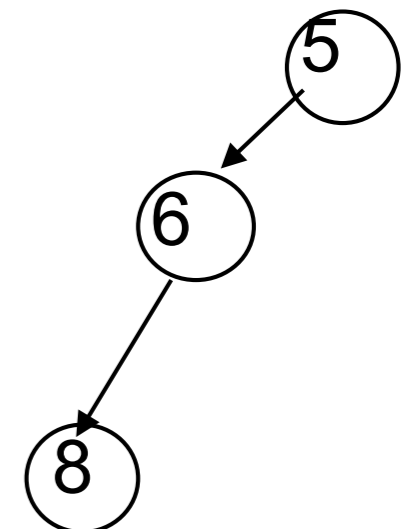
General tree



Binary tree



Not a tree



List-like tree



# Tree Terminology

$M$  is the **root** of this tree

$G$  is the **root** of the **left subtree** of  $M$

$B, H, J, N, S$  are **leaves** (have no children)

$N$  is the **left child** of  $P$

$S$  is the **right child** of  $P$

$P$  is the **parent** of  $N$

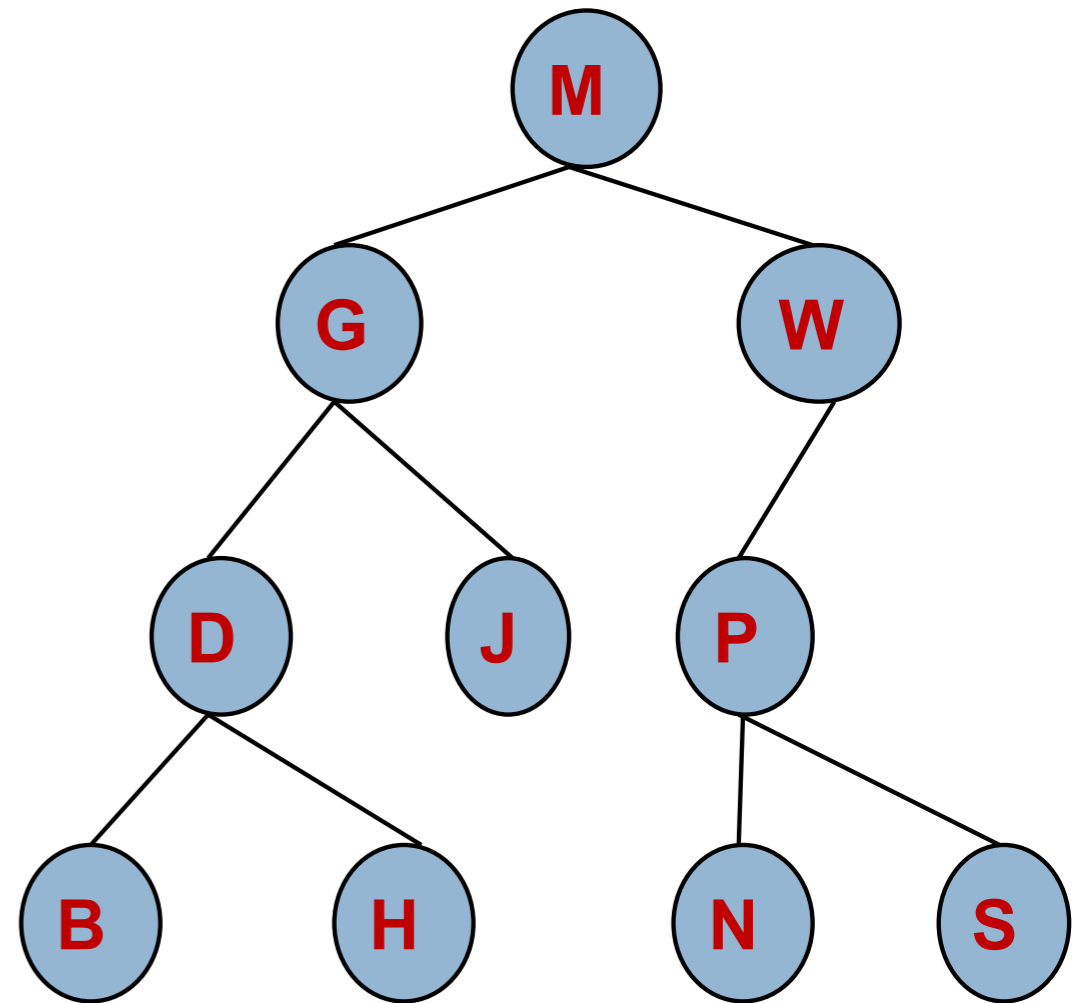
$M$  and  $G$  are **ancestors** of  $D$

$P, N, S$  are **descendants** of  $W$

$J$  is at **depth** 2 (length of path from root)

The subtree rooted at  $W$  has **height** (length of longest path to a leaf) of 2

A collection of several trees is called a \_\_\_\_\_?



```
public class BinaryTreeNode {
    private int value;
    private BinaryTreeNode parent; (null if no left child)
    private BinaryTreeNode left; // left subtree
    private BinaryTreeNode right; // right subtree
                                    (null if no right child)
}
```

```
public class GeneralTreeNode {
    private int value;
    private GeneralTreeNode parent;
    private List<GeneralTreeNode> children;
}
```

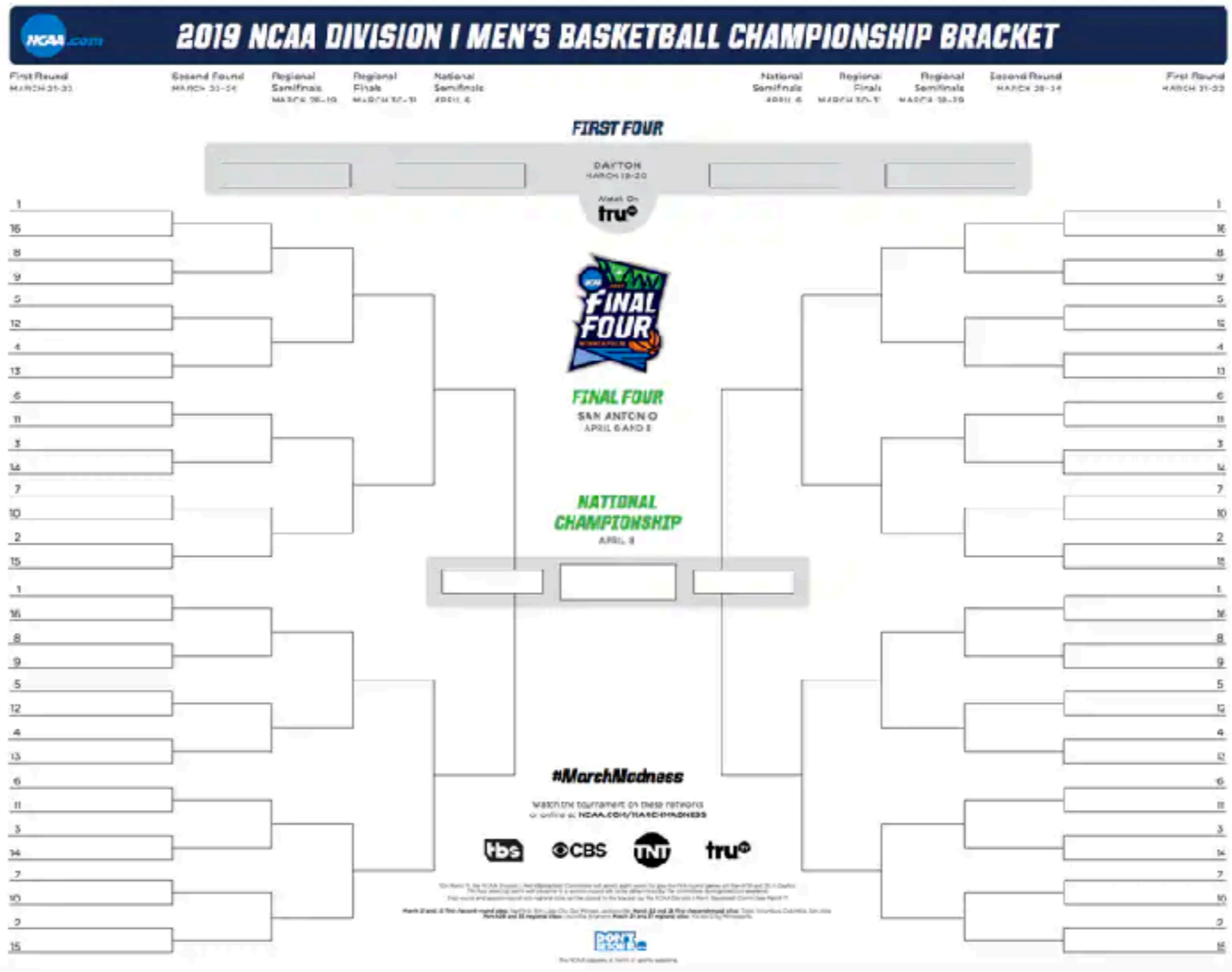
**Why do we need these?**

# Why do we need these?

to represent **hierarchical structure**.

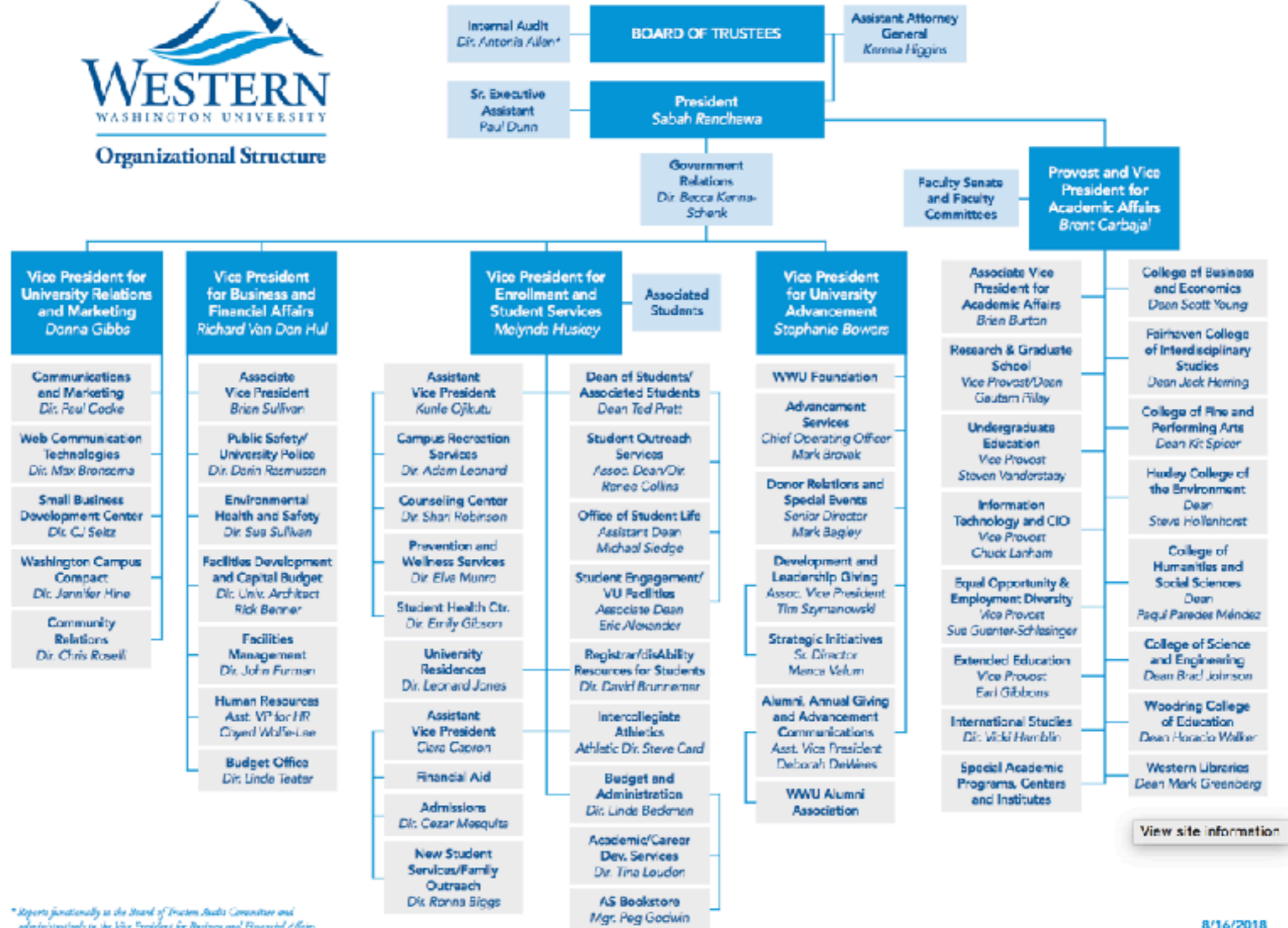
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# Why do we need these?

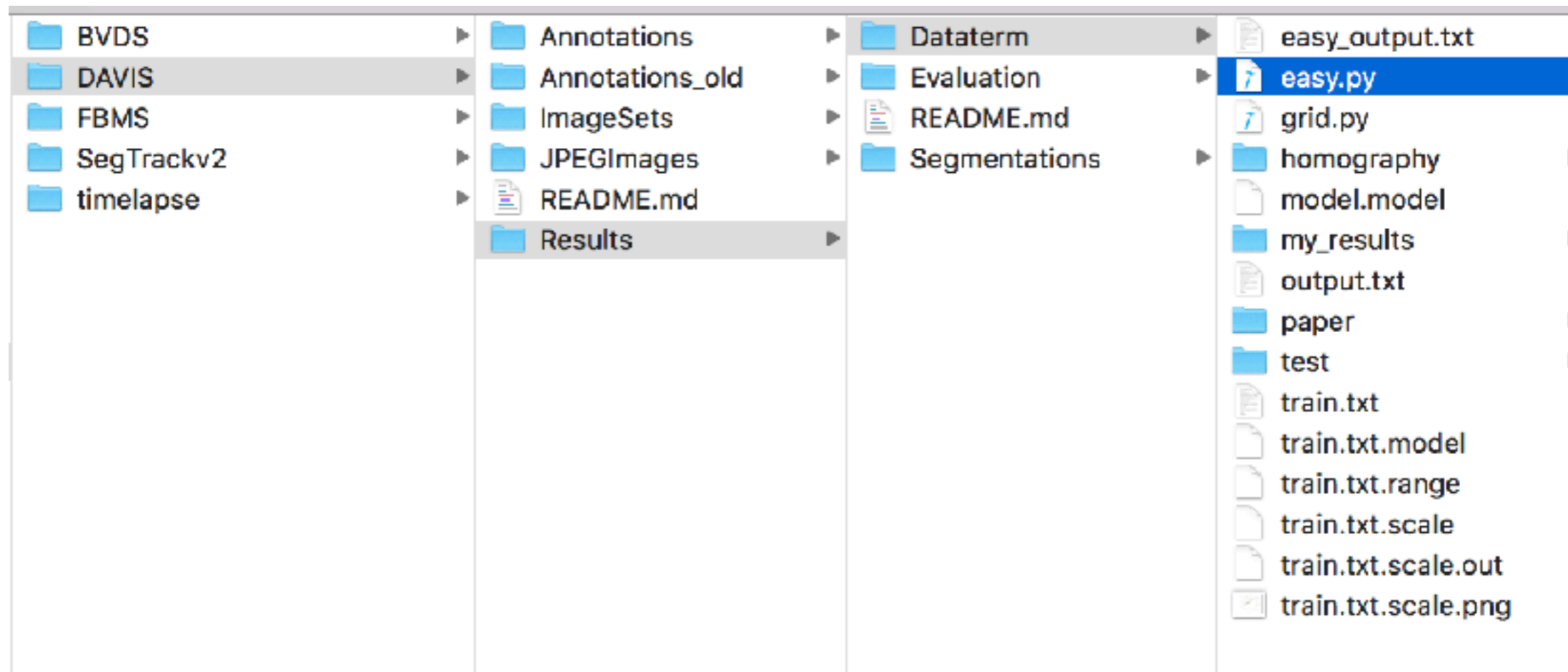
to represent **hierarchical structure.**



\* Reports functionally to the Board of Trustees Audit Committee and administratively to the Vice President for Business and Financial Affairs

# Why do we need these?

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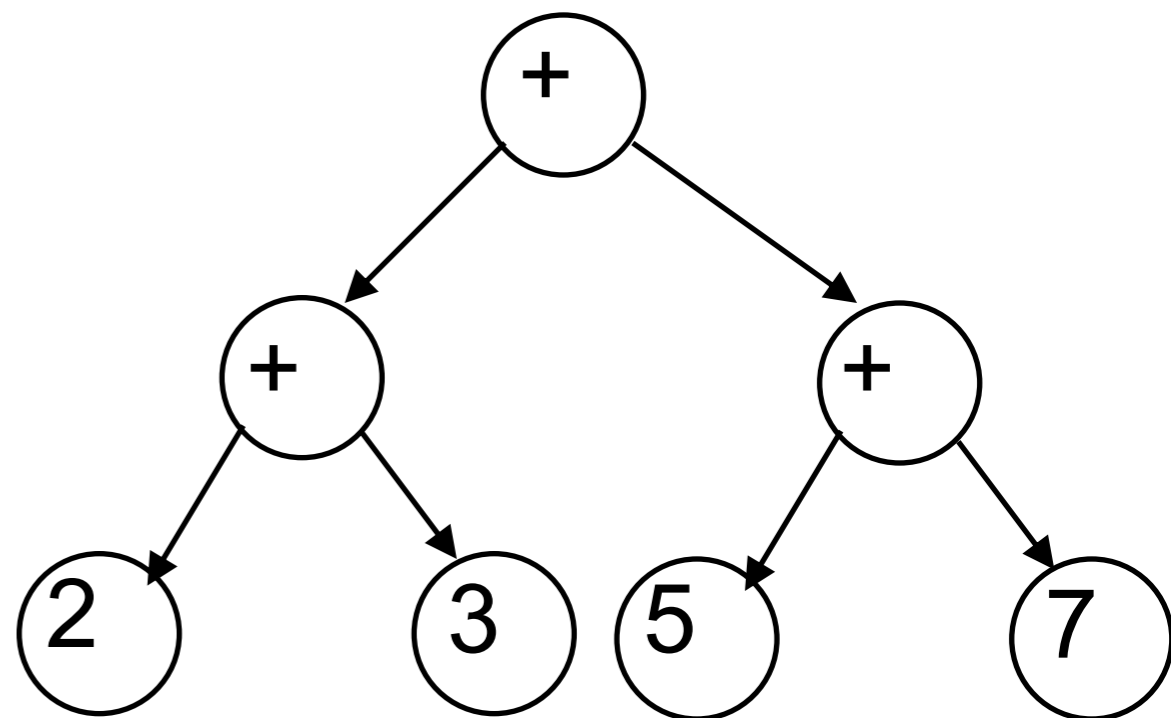
# Why do we need these?

to represent **hierarchical structure**.

## Syntax Trees:

- In textual representation, **parentheses** show hierarchical structure
- In tree representation, hierarchy is explicit in the tree's **structure**

$((2+3) + (5+7))$



Also used for **natural languages** and **programming languages!**



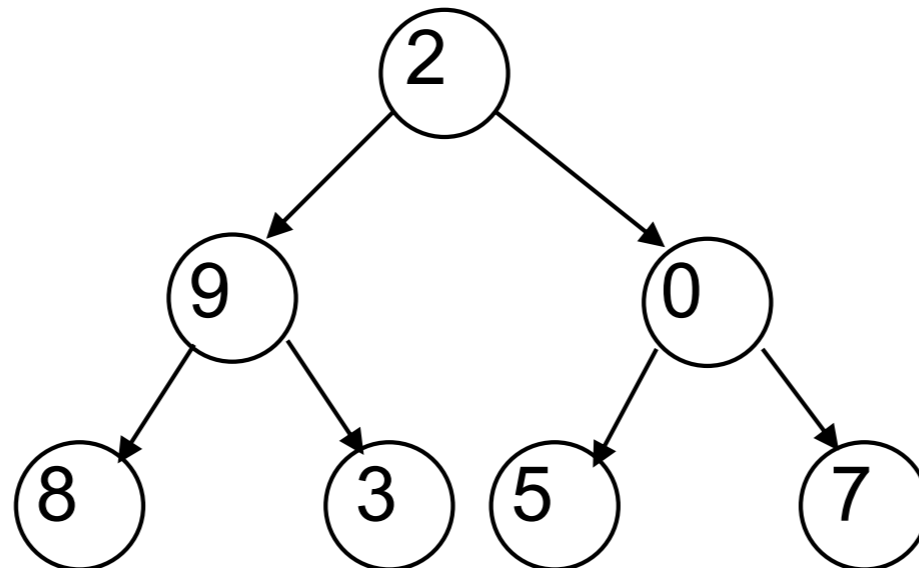
# Why do we need these?

to implement various ADTs **efficiently**.

**TreeSet**, **TreeMap**

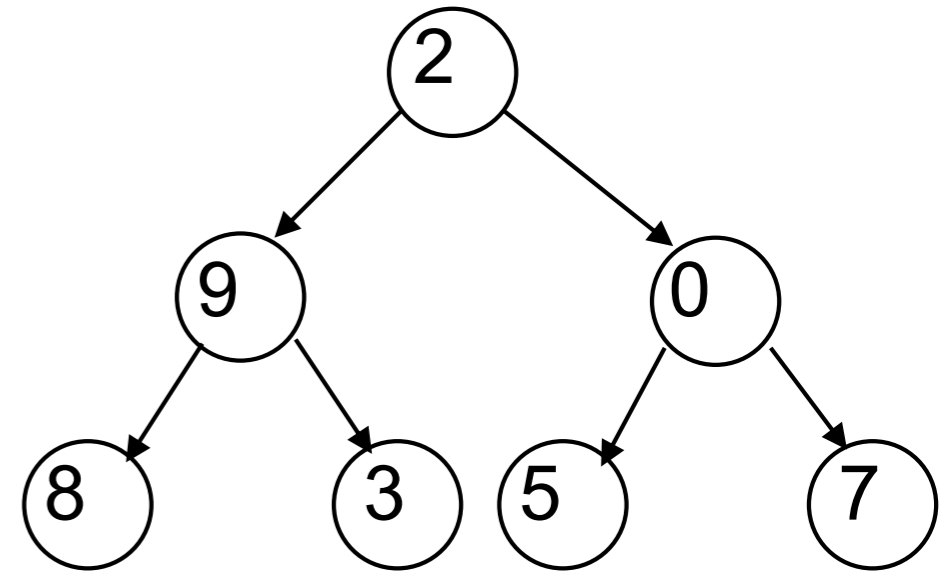
Height of a balanced binary tree is  $O(\log n)$

Consequence: Many operations (find, insert, ...) can be done in  **$O(\log n)$**  in carefully-designed trees.



# Thinking about trees recursively

- **A binary tree is**
  - Empty, or
  - Three things:
    - value
    - a left **binary tree**
    - a right **binary tree**



# Thinking about trees recursively

- **A binary tree is**

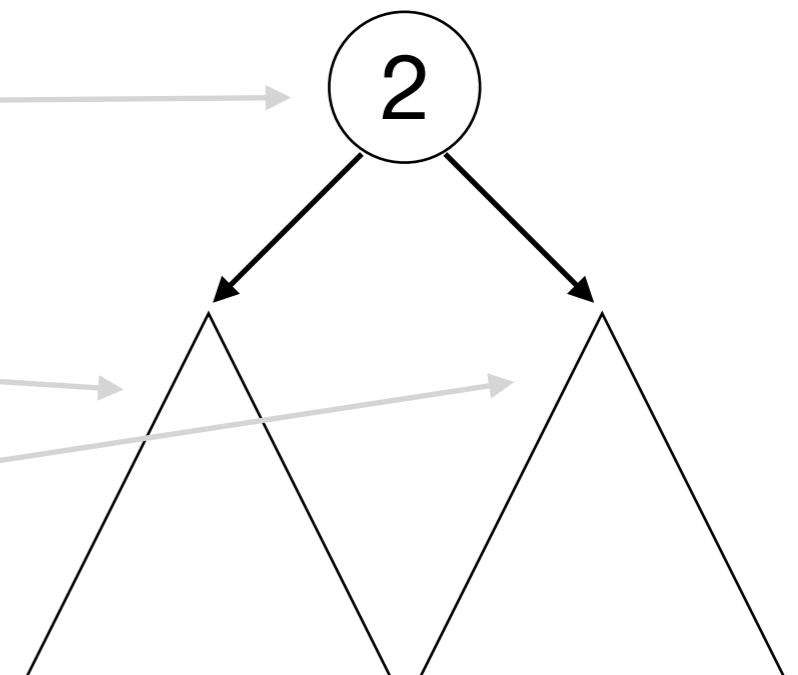
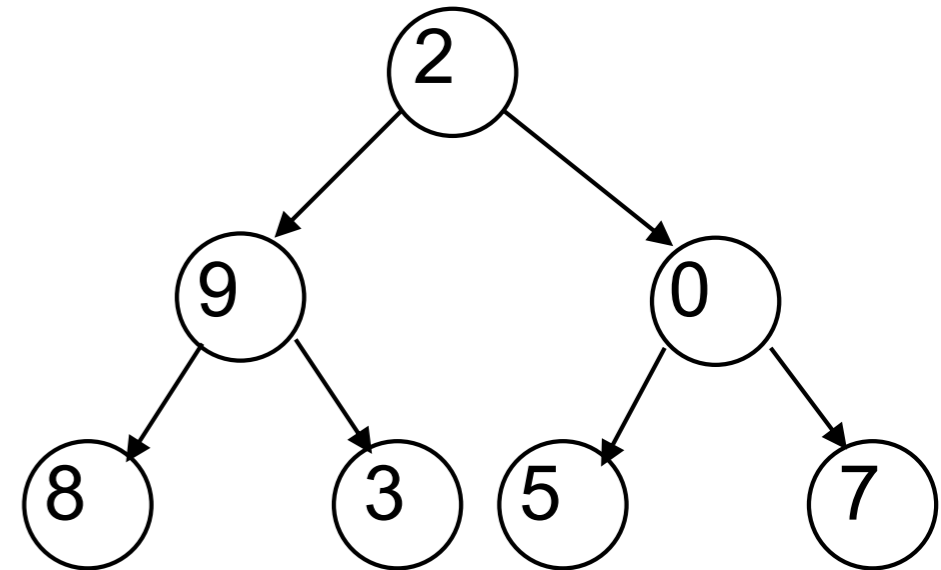
- Empty, or

- Three things:

- value

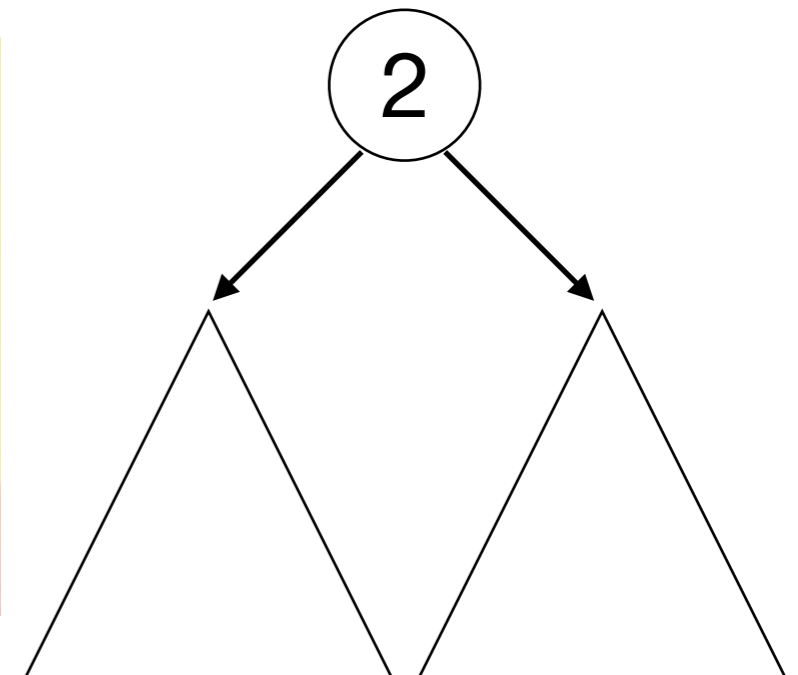
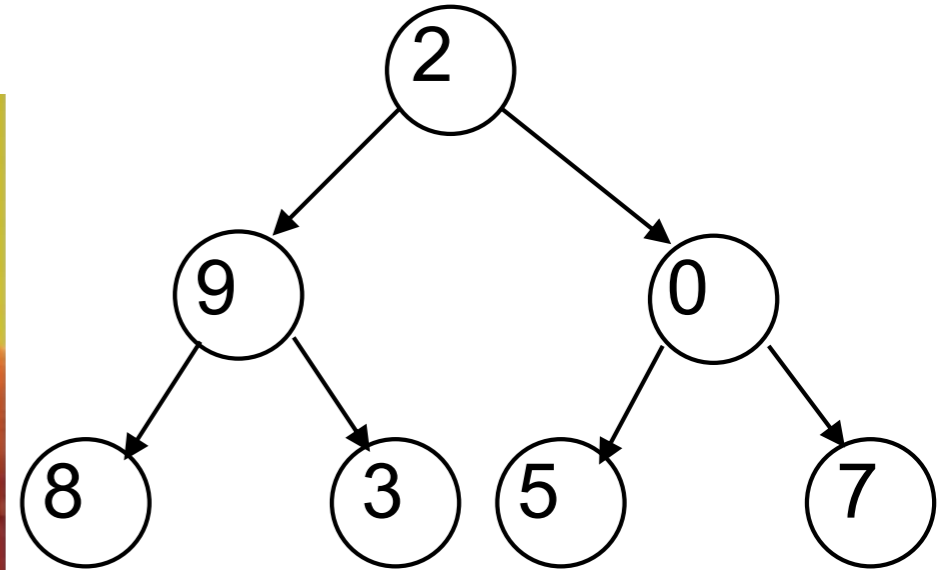
- a left **binary tree**

- a right **binary tree**



# Thinking about trees recursively

- A **binary tree** is
  - Empty, or
  - Three things:
    - value
    - a left **binary tree**
    - a right **binary tree**



# Operations on trees

often follow naturally from the definition of a tree:

- **A binary tree is** Find  $v$  in a binary tree:
  - Empty, or (base case - not found!)
  - Three things:
    - value (base case - is this  $v$ ?)
    - a left **binary tree** (recursive call - is  $v$  in left?)
    - a right **binary tree** (recursive call - is  $v$  in right?)

# Operations on trees

often follow naturally from the definition of a tree:

- **A binary tree is**

- Empty, or

- Three things:

- value

- a left **binary tree**

- a right **binary tree**

Find  $v$  in a binary tree:

```
boolean findVal(Tree t, int v):
```

(base case - not found!)

```
if t == null:
```

```
    return false
```

(base case - is this  $v$ ?)

```
if t.value == v: return true
```

(recursive call - is  $v$  in left?)

```
return findVal(t.left)
```

```
    || findVal(t.right)
```

(recursive call - is  $v$  in right?)

# Tree Traversals

Print (or otherwise process) every node in a tree:

- **A binary tree is**

- Empty, or

- Three things:

- value

- a left **binary tree**

- a right **binary tree**

Print all nodes in a binary tree:

```
boolean printTree(Tree t):
```

(base case - nothing to print)

```
if t == null:
```

```
    return
```

(print this node's value)

```
System.out.println(t.value)
```

(recursive call - print left subtree)

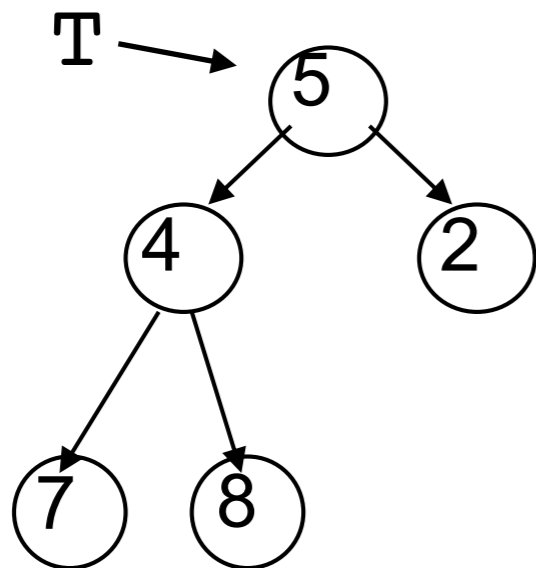
```
printTree(t.left)
```

(recursive call - print right subtree)

```
printTree(t.right)
```

# Tree Traversals

Print (or otherwise process) every node in a tree:



Print all nodes in a binary tree:

```
boolean printTree(Tree t):
```

(base case - nothing to print)

```
if t == null:
```

```
    return
```

(print this node's value)

```
System.out.println(t.value)
```

(recursive call - print left subtree)

```
printTree(t.left)
```

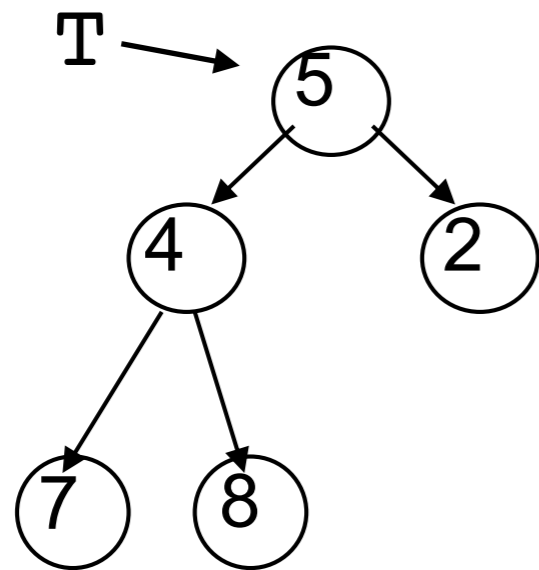
(recursive call - print right subtree)

```
printTree(t.right)
```



# Tree Traversals

Print (or otherwise process) every node in a tree:



**ABCD:** T is a reference to the node with value 5. What is printed by the call `printTree(T)`?

- A. 5 4 2 7 8
- B. 7 4 8 5 2
- C. 7 8 4 2 5
- D. 5 4 7 8 2

Print all nodes in a binary tree:

```
boolean printTree(Tree t):
```

(base case - nothing to print)

```
if t == null:
```

```
    return
```

(print this node's value)

```
System.out.println(t.value)
```

(recursive call - print left subtree)

```
printTree(t.left)
```

(recursive call - print right subtree)

```
printTree(t.right)
```

# Tree Traversals

“Walking” over the whole tree is called a **tree traversal**. This is done often enough that there are standard names. Previous example was a **pre-order traversal**:

1. **Process root**
2. Process left subtree
3. Process right subtree

## Other common traversals:

### in-order traversal:

1. Process left subtree
2. **Process root**
3. Process right subtree

### post-order traversal:

1. Process left subtree
2. Process right subtree
3. **Process root**

# Why do we need these?

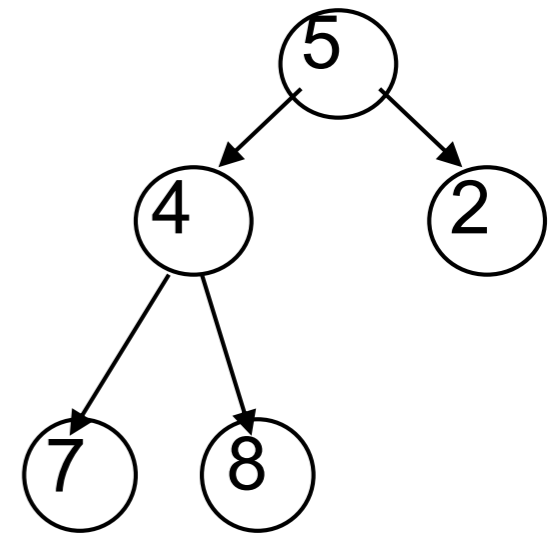
to represent **hierarchical structure**.

Quadrees in graphics and simulation:

<https://www.youtube.com/watch?v=fuexOsLOf10>

# Practice Exercise

- Write the values printed by a:
  - pre-order
  - in-order
  - post-order



traversal of this tree.

# Terminology - Self-Quiz

root

subtree

leaf

child

parent

ancestor

descendant

depth

height

