

CSCI 241

Lecture 3:
Loop Invariants
Insertion and Selection Sort

Roadmap

- Tools for reasoning about algorithms
- Sorting algorithms (A1)
 - Insertion
 - Merge
 - Quick
 - Radix
- Onward to more data structures

Today

- Quiz recap
- Tools for reasoning about algorithms
- Sorting algorithms (A1)
 - Insertion
 - Selection
 - Merge
 - Quick
 - Radix
- Onward to more data structures

```

public class LinkedList {
    public Node head;

    /** if n is in the list, remove all elements
     * *following* it and return true.
     * otherwise, return false leaving the list unchanged
     * precondition: n is not null */
    public boolean truncateAfter(Node n) {
        Node current = head;
        while (current != null) {
            if (current == n) {
                // INSERT CODE HERE

                return true;
            }
            current = current.next;
        }
        return false;
    }
}

```

```

public class Node {
    Node next;
    int value;
}

```

```
public class Factorial {  
    /** return n factorial. precondition: n > 0 */
```

```
public static int fact2(int n) {  
    return n * fact2(n-1);  
}
```

```
public static int fact3(int n) {  
    if (n > 0) {  
        return n * fact3(n-1);  
    }  
    return 1;  
}
```

```
public static int fact4(int n) {  
    if (n > 1) {  
        return n * fact4(n);  
    }  
}
```

```
}
```

```
public class Factorial {
    /** return n factorial. precondition: n > 0 */

    public static int fact2(int n) {
        return n * fact2(n-1);
    }
}
```

```
public static int fact3(int n) {
    if (n > 0) {
        return n * fact3(n-1);
    }
    return 1;
}
```

```
public static int fact4(int n) {
    if (n > 1) {
        return n * fact4(n);
    }
}
}
```

```
public class Factorial {
    /** return n factorial. precondition: n > 0 */

    public static int fact2(int n) {
        return n * fact2(n-1);
    }

    public static int fact3(int n) {
        if (n > 0) {
            return n * fact3(n-1);
        }
        return 1;
    }

    public static int fact4(int n) {
        if (n > 1) {
            return n * fact4(n);
        }
    }
}
```

Goals:

- Be able to analyze an algorithm's correctness using a given loop invariant.
- Be able to describe in words, implement using a loop invariant, and analyze the runtime of:
 - Selection sort
 - Insertion sort

Tools for Reasoning about Algorithms

Precondition, Postcondition

```
/** return the max value in A
 * precondition: A is nonempty
 * postcondition: max value of A is returned */
public int findMax(int[] A) {
    int max = A[0];
    // invariant: max is the largest value in A[0..i]
    for (int i = 1; i < A.length; i++) {
        if (A[i] > max) {
            max = A[i];
        }
    }
    return max;
}
```

Interface, not implementation

The **precondition** is true **before** method execution.
The **postcondition** is true **after** method execution.

Loop Invariant

```
/** return the max value in A
 * precondition: A is nonempty
 * postcondition: max value of A is returned */
public int findMax(int[] A) {
    int currentMax = A[0];
    // invariant: currentMax is the max of A[0..i]
    for (int i = 1; i < A.length; i++) {
        if (A[i] > max) {
            max = A[i];
        }
    }
    return max;
}
```

currentMax is the largest value in:

A[0..1]

A[0..i]

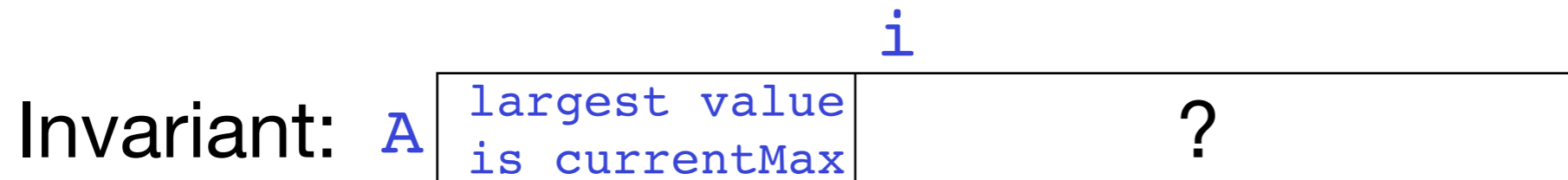
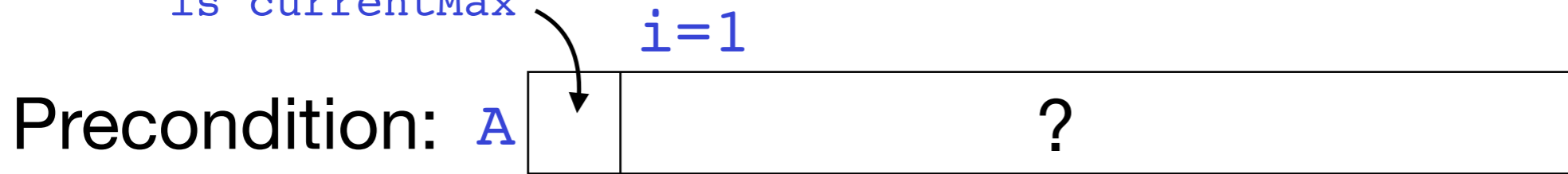
A[0..a.length]

A **loop invariant** is true **before**, **during**, and **after** the loop.

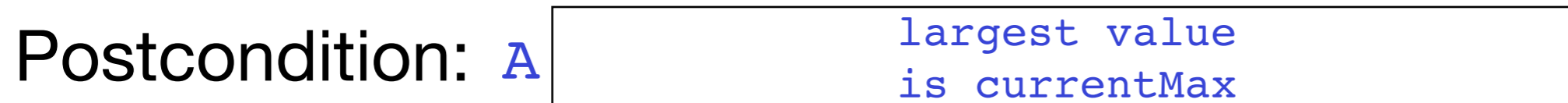
(at the *end* of each iteration)

Loop Invariant

largest value
is currentMax



$i=a.length$



$A[0..1]$



$A[0..i]$



$A[0..a.length]$

The **loop invariant** is true **before**, **during**, and **after** the loop.

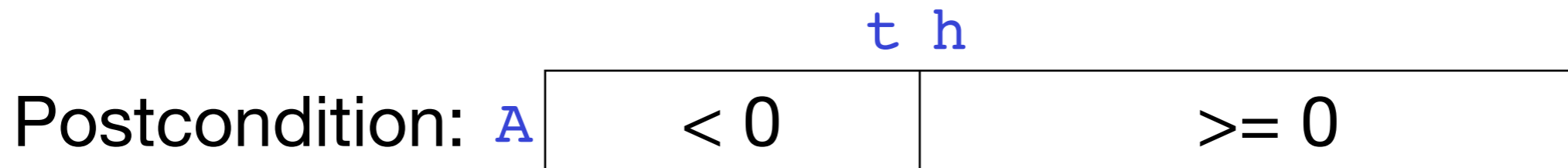
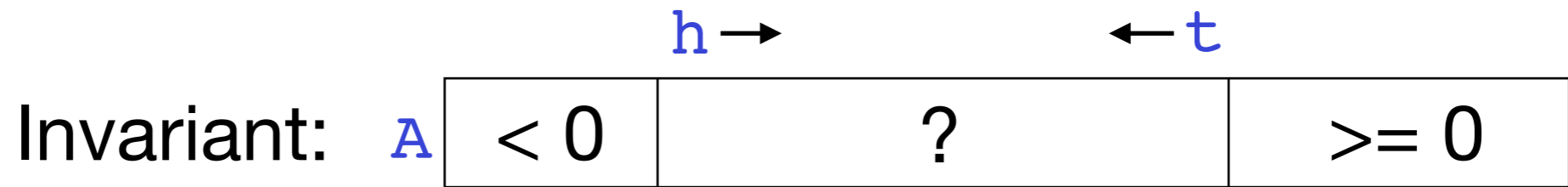
Another Example

```
/** rearrange A so all negative values are to  
 * the left of all non-negative values */  
public void separateSign(int[] A);
```

Postcondition: A

< 0	>= 0
-----	------

```
/** rearrange A so all negative values are to
 * the left of all non-negative values */
public void separateSign(int[] A) {
```



Four concerns:

1. **Initialization:** Make the invariant true at the start.
2. **Termination:** Make the loop end when the postcondition is true.
3. **Progress:** Make progress towards the postcondition.
4. **Maintenance:** Make the invariant true after each iteration.

Insertion Sort

Insert $A[i]$ into the sorted sublist $A[0..i-1]$.

Selection Sort

Find the smallest element in $A[i..n]$ and place it at $A[i]$.

<https://visualgo.net/bn/sorting>

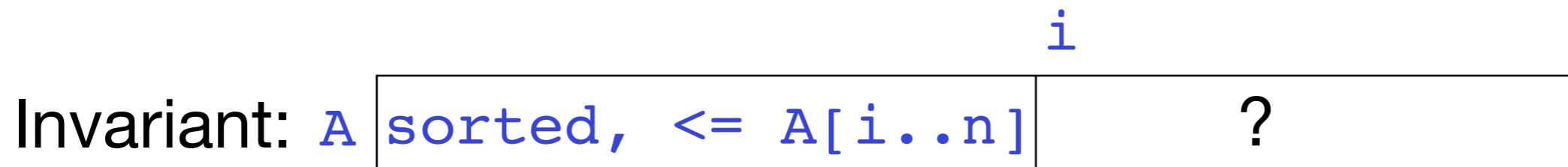
Insertion Sort

Insert $A[i]$ into the sorted sublist $A[0..i-1]$.



Selection Sort

Find the smallest element in $A[i..n]$ and place it at $A[i]$.



<https://visualgo.net/bn/sorting>


```
insertionSort(A):
```

```
    i = 0;
```

```
    while i < A.length:
```

```
        // push A[i] to its sorted position
```

```
        // increment i
```

```
selectionSort(A):
```

```
    i = 0;
```

```
    while i < A.length:
```

```
        // find min of A[i..A.length]
```

```
        // swap it with A[i]
```

```
        // increment i
```

Developing InsertionSort

Precondition:

A

?

Invariant:

A

sorted

i

?

Postcondition:

A

sorted

Four tasks:

1. **Initialization:** Make the invariant true at the start.
2. **Termination:** Make the loop end when the postcondition is true.
3. **Progress:** Make progress towards the postcondition.
4. **Maintenance:** Make the invariant true after each iteration.

Developing SelectionSort

Precondition: A

?

Invariant: A

sorted, $\leq A[i..n]$

i

?

Postcondition: A

sorted

Four concerns:

1. **Initialization:** Make the invariant true at the start.
2. **Termination:** Make the loop end when the postcondition is true.
3. **Progress:** Make progress towards the postcondition.
4. **Maintenance:** Make the invariant true after each iteration.

```

insertionSort(A):
    i = 0;
    while i < A.length:
        j = i;
        while j > 0 and A[j] > A[j-1]:
            swap(A[j], A[j-1])
            j--
        i++

```

ABCD: What's the best and worst-case asymptotic runtime complexity of insertionSort?

	Best	Worst
A	$O(n)$	$O(n)$
B	$O(n^2)$	$O(n)$
C	$O(n)$	$O(n^2)$
D	$O(n^2)$	$O(n^2)$

Why is this best-case runtime interesting?

```
insertionSort1(A):
```

```
    i = 0;
```

```
    while i < A.length:
```

```
        j = i;
```

```
        while j > 0 and A[j] < A[j-1]:
```

```
            swap(A[j], A[j-1])
```

```
            j--
```

```
        i++
```

```
insertionSort2(A):
```

```
    i = 0;
```

```
    while i < A.length:
```

```
        j = i;
```

```
        tmp = A[i];
```

```
        while j > 0 and tmp < A[j-1]:
```

```
            A[j] = A[j-1]
```

```
            j--
```

```
        i++
```

ABCD: What's the best and worst-case asymptotic runtime complexity of insertionSort2?

	Best	Worst
A	$O(n)$	$O(n)$
B	$O(n^2)$	$O(n)$
C	$O(n)$	$O(n^2)$
D	$O(n^2)$	$O(n^2)$

Practice problems

Write a precondition, postcondition, and loop invariant for the *inner* loop of InsertionSort.

Develop SelectionSort using the precondition, loop invariant, and postcondition by completing the four tasks:

Initialization

Termination

Progress

Maintenance