Talked about Threads, Processes, and Memory
Now for "slower memory", e.g. Disks (Chapter 11 Mass Storage Structures)
"Disks" come in a variety of flavors ...
- RAM disk -- dedicate some ram to look like disk
- SSDs (Solid State Disks)
- Hard Disk -- Many Flavors
- Optical Disks
- Network "Disks"
- Magnetic Tape?
- Thumb drives
- Communications with these devices over a "BUS" (From Chapter 12)
  - Parallel or Serial
  - Advanced Technology Attachment (ATA, parallel)
  - serial ATA (SATA, serial)
  - universal serial bus (USB)
  - fibre channel (FC)
  - various flavors of "SCSI" (Small Computer System Interface)
  - others ....
Typical disks:

- Sector -- a single unit of read/write
- Head -- a device to physical read/write on the disk
- Platter -- a side of a disk on which a head reads/writes
- Arm -- device on which heads are mounted, moves in and out
- Cylinder -- all sectors addressable without moving the arm
- Rotational speed in RPM (rotations per minute)
  - 5400, 7200, 10000, 15000 RPM
- Interesting times:
  - Transfer rate
  - Seek time
  - Rotational latency

Typical operations:

- Seek to a cylinder
- Read/Write a sector, select platter and sector on the platter
- More recent (LBA): Read/Write logical disk sector, no seek involved
Mass Storage

Solid-State Disks

☐ nonvolatile memory used like a hard drive
  ☐ flash-memory (typically NAND semiconductors)
  ☐ Advantages?
    ☐ no moving parts, faster transfer, no seek time, less power
  ☐ Disadvantages?
    ☐ reads (fast) vs writes (slower), standard bus tech limits speed
    ☐ max of 100,000 writes (erase, rewrite), lifespan measured in "Drive Writes Per Day"
    ☐ (May use "wear leveling" algorithms, often implemented by the NVM controller.)
  ☐ LBA addressing is used, sometimes OS is told tracks/heads/... but they don’t exist
  ☐ SSD is starting to make rotating disks obsolete
    ☐ 500GB ssd about $61 (10/2020) (down from 2019, was $75)
    ☐ 2TB rotating about $55 (10/2020) (up from 2019, was $50)
       ☐ 2TB ssd about $200 (10/2019) (down from 2019, was $230)

Tape -- still used in some places

☐ Sequential structure, no random access
  ☐ Transfer speeds similar to disk when ready
Disk Structure for most modern disks:
- addressed as a large one-dimensional array of logical blocks
- logical block size some power of 2, 512 usually the smallest
- bad block mapping makes it hard to map logical block to disk geometry
- recent disks -- use same linear size per sector
  - longer tracks have more sectors
  - drive speed changes as head moves in/out

Disk Attachment -- Where is the disk
- Host-Attached storage
  - "same box"
  - High-end, Fibre channel (FC)
    - multiple disks, multiple hosts
- Network-Attached storage
  - NFS, CIFS, Andrew -- network based file systems (later)
  - iSCSI -- SCSI over IP
- Cloud Storage
  - Storage on someone else’s computer
  - API based, WAN based access
- Storage Networks -- private networks not connected to internet
Host Attached Storage vs Network Attached Storage

- Network issues -- storage on network causes network traffic
- Storage Area Network -- e.g. storage devices on one NIC, LAN on another

Disk Scheduling

- Idea that you have a "queue" of disk requests
- How to best schedule them
  - Light load ... no issue
  - Heavy load ... how to best schedule them
    - FCFS scheduling
    - Shortest seek time scheduling
      - may cause starvation
    - Scan algorithm
      - AKA elevator algorithm
      - Circular scan
    - Linux "deadline" scheduler (RedHat)
    - NVM scheduling, typically FCFS
      - priority of reads over writes?
      - writes slow, issues of wear, ...
Other topics

Things to read about

- disk formatting -- partitions, volumes
- bad block management
- Swap space management
- RAID (Redundant arrays of independent disks)
  - making larger virtual disks by striping (RAID 0)
    - Performance gains by parallelism
    - No redundant bits
  - making error correction/recovery by redundant disks
    - RAID 1: mirrored disks
    - RAID 2: Memory-style error-correcting codes (ECC)
    - RAID 3, 4, 5, 6: other techniques ...
- Stable-Storage -- Information is never lost
  - How to implement it?
    - multiple storage devices
    - NVRAM as a cache
OS is a hardware manager ... talked about CPU, Memory, Disks ...

- Other I/O Devices
  - transmission device (network, bluetooth,...)
  - human-interface devices (screen, keyboard mouse, audio, joystick)
  - specialized: sensor and control, ... (large variety)

- Memory mapped I/O
  - Address range communicates to devices, not real memory
  - Device Control register
  - Device Data Register
  - Device Memory -- could be large

- Techniques for I/O
  - Polling -- (assignment 1)
  - Interrupt driven -
    - Start operation, return to other stuff
  - Interrupt from I/O device
  - Interrupt processing needs to be fast
  - DMA and interrupts
Application I/O Interface

- Need an API for standard treatment of I/O devices
- Low level -- Device driver
  - Interface between Kernel and device driver
  - complete to deal with all devices
- Higher level -- user view may look like a "file"
  - UNIX - device file, (/dev/...)
  - Windows -- a device object ... that can be opened by file name
- Device characteristics
  - character-stream vs block
  - sequential vs random access
  - synchronous vs asynchronous
  - sharable vs dedicated
  - speed of operation
  - read/write properties
  - no direct user interaction ... e.g. clocks and timers
- Unix: Block and Character Devices
  - All devices look like a character device, some also look like a block device
  - Interface is slightly different between the two
- Other devices: clocks, network: Different interfaces

Ignore the rest of chapter 12, may come back later
File Systems Interface (Chapter 13)

- File System -- an abstraction on top of storage
- Typical Services
  - File abstraction
  - File manipulation
  - File protection
- Most visible service of OS
- Large code base in most OSes

File abstraction
- Bag of bits?
- Known content? (e.g. is .txt for OS or users?)
  - By the OS?
    - Executable files
  - By user land Tools?
    - Required
File System Basics

- Standard attributes
  - Name: (symbolic, human readable)
  - Identifier: unique tag
  - Type: system tag
  - Location: where it is located on the storage
  - Size: both logical and physical size (if different)
  - Protection: who has what kind of access
  - Time, date, user identification, ...

File Operations

- Creation: Adding information
- Writing: adding information, file position pointer
- Reading: file position pointer also
- Deleting: removing information
- Truncating a file: removing information

May be many other file management routines
- renaming, moving, status, ...
Management of files in the kernel

- Open syscall: looking up information ... look up file only once
- Kernel keeps an "Open File Table" in the kernel
- Open syscall:
  - lookup file in file system (could be expensive)
  - "cache" information in the open file table
  - return a "handle", some data to uniquely represent file
- Close syscall:
  - done using the file, allow file to reclaim space
- Open and Close with shared files
  - multiple applications may open file at the same time
  - in systems with fork(), both processes have access to files
  - Typically ... two levels of tables in this case
    - Kernel wide "open file table"
    - Per process "local file table" that points to open file table
- Kernel global open file table
  - File pointer -- offset into file
  - File-open count -- how many local file entries point here
  - Information for file location on disk
  - Access rights
- Local table: Open flags, pointer to global open file table
Locks and File types

Locks -- shared or exclusive
□ shared read locks
□ exclusive locks
□ mandatory or advisory
□ deadlock issues here

File Types
□ Kinds of data in files
□ executable, text, scripts, DataBase, ....
□ How does OS know what is in the file?
□ file name ... extension (DOS, Windows)
□ .cpp -- file type?
□ C pre-processor input?
□ .app ?
□ OS X, extension on a directory!
□ extra information?
□ Mac: creator -- program that created a file
File Types (page 2)

☐ know how to rebuild executable files? (TOPS 20)
  ☐ Used time information with source to executable
  ☐ Source changed, recompile before running

☐ UNIX?
  ☐ "magic" numbers to start off files
  ☐ file(1) command
File Structure

- Executable ... OS needs to know structure to load file
  - Blitz: OpenFile.LoadExecutable

- Other files?
  - VMS -- knew structure of system files

- Problem?
  - what if your app doesn’t want to use a known structure

- Text vs Data?

- Bag of bits?

- Mac -- Resource and Data "fork"

- Windows -- Multiple "streams" per file

- Internal structure
  - Any kind of packing?
  - Standard encoding?
    - Line in a text file? NL, CR/NL, CR
    - MPE/3000: text file, 80 character lines, all chars present
Access Methods

User level access to the file:
- Sequential
  - "tape model"
  - Sequential access
  - Possibly do "skip +/-n records"
- Rewind
- Go to end
  - (Tape model, multiple files per tape, double EOF => EOT)

- Direct (relative access)
  - Each read/write includes "record" number
  - Each number is a "relative record" number to start of file

- Should an OS provide both?
  - How about sequential access using direct files?
    - like UNIX: keep a file pointer
  - How about direct access using sequential files?
    - very bad!
Other Access Methods?

- Hash table?
  - e.g. Key/Data pairs as basic storage element
  - Also can be stored by trees
- Index file -- keep keys, pointer to data
- IBM ISAM -- indexed sequential-access method
  - two level of indexes to access file
General Disk Structure

File system may depend on storage
- RAM disk -- short life, temp file systems, simple structures
- Collection of disks -- long life, reliable, error protection, hot swapping
- Large disk, subdisks (minidisks, partitions, slices)
  - Allows multiple kinds of file systems on one disk
- Special kinds of file systems?
  - procfs -- a file system interface to "process manager"
  - ZFS -- a "pool" based "general file system"
  - coda, smb, afs, nfs ... -- network file systems
- Volume -- contains a FS.
  - May be anywhere from part of a disk to multiple disks
Directory overview

- Directory Operations
  - lookup (search)
  - add (create)
  - delete
  - list
  - rename
  - traverse the file system
Directory Structures

Single level directory
- Blitz "Stub file system"

Two level directory
- user/file -- top level contains no files
- Or volume:/user/file

Tree structured directories
- current directory, absolute path, relative path

Acyclic Graph structured
- Directory have just "links" to files or directories
- single file can appear in many directories

General Graph structured
- Acyclic?
- Livermore Timesharing System ... full graph
  - traversal algorithms had to detect cycles

Data stored in Directory Entry
- Full information: e.g. DOS
- Pointer to full information: e.g. UNIX UFS
Volume access

Each file system is placed on a "volume"

Multiple volumes to access, How?

- DOS/Windows (in USER space)
  - volume ID
  - path within that volume
  - User needs to see the volume

- UNIX -- File System "mount"
  - Associates a directory on one file system with the root of another
  - System mounts one file system as "Root"
  - Other file systems are mounted on directories of Root
  - User does not need to see mounts
  - User does not need to know file system types
  - Automounting ...
    - to the desktop (Mac)

- Windows?
  - internally does mounts
  - exposes volume via special "mounts"
  - now allows full mounts
File Sharing

On the same OS with multiple users

- need protection and sharing to be considered
- what kinds of sharing
  - read only sharing?
  - read/write sharing?

Remote file systems

- NFS, DFS, SMB, FTP -- different kinds of files
  - (Some systems can "mount" remote files via ftp.)
- sshfs -- an integrated solution for ssh access to files
- Lots of issues in remote file systems -- not much here yet
- client-server fs peer-to-peer
- authentication systems ... distributed naming services ...
- larger number of failure modes
File consistency

How are files shared ... how do reads and writes interact

- Immutable-Shared-Files semantics
  - Once shared, a file can never change

- Session Semantics
  - File gets a "snapshot" at open
  - Changes are not committed until close
  - Changes are not visible unless opened after a close

- UNIX Semantics
  - Writes are visible immediately to any process with an open file
  - Allow processes to interfere with each other.

- Network file systems have done all 3.
  - NFS -- UNIX
  - AFS, Coda -- mostly session semantics
    - (process on the same machine get UNIX semantics)
  - SPRITE (Berkeley, very old) -- read only shared
Protection

reliability -- safe from physical damage
protection -- safe from improper access
Protection may depend of use of file system
- Operations to control: read, write, execute, append, delete, list, change attributes ..
  - Possibly others ... rename, copy, create
- Special directories ...
  - take and give directories at LLNL

Approaches to access control
- Access Control Lists
  - each file has a list of users and allowed operations
  - not on the list? no access
- Drawback?
  - Long lists
Domain based access:
- Owner, Group, Universe
- Each file has protection for each domain
- Access checks user’s domain membership
- Drawback?
  - Hard to select a single user

Typical implementations
- Primary protection by domain
- Secondary protection by ACLs

Examples:
- UNIX: primary protections: read, write, execute
- NT: full control, modify, read&execute, read, write, ...
  - ACL "who" can be a domain or a user
- DOS: nothing!
- Variety of ways to set these:
  - NT: typically a GUI
  - Solaris: has both UNIX and ACL
    - getfacl(1) and setfacl(1)

Read 13.5 Memory-Mapped files ... we talked about them earlier
Typically file systems are stored on disks of some kind ... They provide:
- rewrite: read data, modify, write back to same location
- random access to any block of data ... may take time
- blocks are used for I/O efficiency

Basic File Systems -- Typical hardware and software components
- Disk
- Device Driver -- knows how to control disk
- Basic File System -- uses Device Driver to operate, manages buffers, caches
- File-organization module -- knows about file structure
- Logical file system -- manages meta-data information
  - meta-data -- data about the file, size, date, ...

Management of open files ...
- Idea of a Virtual File system ...
- One interface to ALL file systems implemented by OS
  - UNIX V-node
    - All file systems implement same API for OS to use
    - Core OS knows nothing about actual FS detail
- Best if implemented as a layers of "independent" subsystems
FUSE -- more recent Abstraction...
- Implementation of a file system in user space
- OS passes API calls to user space
- User space program (daemon) implements FS

On Disk Structures Vs In Memory Structures
- On Disk:
  - Total information to access all data
- In Memory:
  - Caches of On Disk information
  - Dynamic information:
    - Mount information
    - Open files and file pointers
      - per-process information (file handle, file descriptor)
- Issue:
  - Keeping data in memory in sync with disk
  - partial writes to disk in case of OS failure
Typical Disk Structures:

- Boot control "block" -- information needed by ROM/OS for boot
- Volume control "block" -- core information on FS
  - UFS: superblock, NTFS: master file table
- Directory Formats
- FS block management structures
- File/Directory block management
Directory Implementation

Directory:
- Keeps names of files with method to lookup meta-data
- Simple Method: linear
  - Fixed or variable sized entries
  - Entry data depends on kind of FS
  - Search time $O(n)$, $n$ number of entries
  - Insert/Delete time?
- Hash table:
  - $O(1)$ search time, insert, delete time
  - Collision techniques?
  - Base hash table size
  - Dynamic issues hash tables
- Some kind of tree storage:
  - Trees in a linear file?
  - File system block as a node?
Allocation methods

Allocation of data blocks (sectors) for files

- Simple: Contiguous Allocation
  - Define a linear ordering of sectors
  - File starts at LBA (logical block address) X
  - Data contained in next Y blocks
- Issues?
  - Random access -- easy
  - Sequential access -- easy
  - Dynamic file size -- hard
  - Creating a new file, unknown space needs
    - Start in largest block
  - Extending a file -- hard
  - Ends up with external fragmentation
  - May need a de-fragmentation function
- Live or offline?
- Used by RT-11, PDP-11 computers
Linked Allocation
- directory/meta-data has first block address
- each block has a "next block" address in the block

Issues?
- creating -- easy
- writing/extending -- easy
- sequential access -- easy, may take longer than contiguous
- random access -- hard
- ends up with internal fragmentation
- dynamic file size -- easy
- data in each sector is less than sector size

Reliability?
- data corrupted (link) => lose the remainder of file

Doubly linked list?
- Store filename, block number?
Allocation methods (page 3)

- **FAT -- File allocation table (MSDOS, OS-2)**
  - array of block numbers, one for each data block on FS
  - links are in the FAT, no loss of data on disk
  - not allocated: 0 entry or on a free list
  - Disk reads for FAT and file
Indexed allocation
□ Block of "pointers to data blocks"
□ Each file has its own index block
□ Directory has pointer to index block
□ Issues?
  □ Create, read, write, append, random access .... easy
    □ Run out of space in index block?
□ Small files ... lots of wasted space in index block
□ Small index blocks ... small files
□ linked scheme, last entry in index block is to next index block
□ multi-level index scheme, top level points to index blocks ...
□ UNIX UFS combined method
  □ small index block, one level regular index block, 2 & 3 level ...

FS Performance ... a major component of "system feels fast"
□ FAT/NTFS systems -- De-fragmentation -> get files closer to contiguous
□ Berkeley’s changes to UFS for FFS
  □ Allocate file in the same cylinder, not just contiguous
  □ Other disk related tweaks .... of which many are not valid any more
Free-Space Management

Free disk space management needs to be done

- Keep track of unallocated blocks
- May use unallocated blocks to help keep track

- Bit Vectors
  - one bit per FS block
  - 0 allocated, 1 free
  - Advantage
    - compact
    - ffs (find first set) instructions
  - Disadvantage
    - large bit maps (e.g. 1TB file system)
    - ffs instructions need all bits in memory

- Linked List
  - Either in the Disk Blocks or the FAT
  - Advantage -- relatively easy
  - Disadvantage -- May be hard to allocate from same cylinder ...
Counting (aka run length encoding)
- Free blocks usually come in groups
- Linked list has first block, number of blocks free
  - Advantages
    - An empty disk has one entry in the list.
  - Disadvantages
    - Turns into simple linked list after much use

Space Maps
- Sun’s ZFS -- designed for a huge number of files
- Can include multiple file systems
- Meta Data I/O is of importance
- Divides space into meta-slabs each with a spacemap
- One spacemap easily fits into memory ... read, modify, write
- ZFS also depends on transaction processing and log file systems
  - more later on log file systems

TRIMing unused blocks
- NVM flash-based, writing is very slow
- Tell device a block is no longer in a file so it can be erased
- Management of free "lists" when rewrite is expensive
Efficiency and Performance

Disk is the major bottleneck in OSes.
- Name lookups can be expensive
- Space allocation can be costly
- Size of pointers to files => space used to store them
  - 16, 32, 64 bit pointers
  - ZFS: 128 bit pointers
- Reading and writing can cause system to slow down
  - E.g. write a block, now need it again
  - (page out, page fault is an example)

Buffer cache
- Cache of Disk blocks Read/Written
- Page cache and FS cache VS Unified buffer cache
  - LRU replacement algorithm in cache
- Synchronous vs Asynchronous writes
- Read Ahead for buffer management of read files
File System Maintenance

File de-fragmentation

- Why needed?
- Which FSes need this?

File system consistency checker

- diskchk in DOS
- fsck in UNIX

- Make sure all structures are in tact.
  - Free inodes and Used inodes add up to total
  - Free blocks and Used Blocks add up to total
  - File meta-data matches reality (e.g. nlinks)
  - All files (inodes) are reachable in directory tree

- (ToyFs needs a fsck program! or a check option to the toyfs program)
Log-Structured File Systems

- DB style transactions as applied to file systems
- Tries to make sure that we never need to repair much
- Basic Idea
  - Write to the "log" what will be done (e.g. metadata)
  - Do what you said
  - Write to log you have done it.
  - Log can be a circular buffer of appropriate size
  - At "recovery time" can see that a log entry was not finished
    - Abort or replay entry
  - Log writes are sequential and thus very fast

- Used in many file systems now, NTFS, LFS (BSD), ext3fs, FFS (BSD)

Other types of things have been used to improve speed and reliability
- ZFS -- snapshot, never overwrites blocks, no FSCK ...

Backups -- another way to preserve your FS data
- Full backups vs Incremental backups

Read 14.8 (WAFL)
File System Internals

- Kinds of file systems
  - general-purpose -- files, directories -- on long term storage
  - tmpfs -- a file system in main memory
  - objfs -- a "virtual" file system, access to kernel symbols
  - ctfs -- a virtual file system, "contract information"
  - lofs -- a "loop back" file system
  - procs -- a virtual file system with system information, process information
  - ufs, ffs, extXfs, zfs -- general purpose file systems

- File-System Mounting
  - Blitz -- toyfs Init() -- read the first sector, get ready to use
  - General term for that is mounting
  - Mount point ... place to access the file system
  - DOS/Windows:
    - drive letter:\path\to\file
  - UNIX:
    - mount on a directory (usually empty, hides directory contents)